SPHERICAL CODES AND DESIGNS

1. INTRODUCTION

A finite non-empty set X of unit vectors in Euclidean space \mathbb{R}^d has several characteristics, such as the dimension d(X) of the space spanned by X, its cardinality n = |X|, its degree s(X) and its strength t(X).

The degree s(X) is the number of values assumed by the inner product between distinct vectors in X; that is,

$$s(X) = |A(X)|, \quad A(X) = \{\langle \xi, \eta \rangle; \xi \neq \eta \in X\}.$$

We shall consider sets X having the property that A(X) is contained in a prescribed subset A of the interval [-1, 1[. Such sets are called *spherical* A-codes. We are interested in upper bounds for n = |X|, and in the structure of spherical A-codes which are extremal with respect to such bounds. For $A = [-1, \beta]$, this problem is equivalent to the classical problem of non-overlapping spherical caps of angular radius $\frac{1}{2} \arccos \beta$; for $\beta \le 0$ exact formulae have been obtained by Rankin [21]. In Section 4 of the present paper we derive bounds for the cardinality of spherical A-codes in terms of the Gegenbauer coefficients of polynomials compatible with A. Apart from these we find, for $s = |A| < \infty$,

$$n \leq \binom{d+s-1}{d-1} + \binom{d+s-2}{d-1}.$$

Several examples are mentioned of sets X for which any of these bounds is attained. As an application, a non-existence theorem for graphs is proved.

The notion (and the name) of a spherical t-design is explained in Section 5. It serves to measure certain regularity properties of sets X on the unit sphere Ω_d . A spherical 1-design has its centre of mass in the centre of the sphere. A spherical 2-design is what Schläfli called a eutactic star, essentially the projection into \mathbb{R}^d of n orthogonal vectors, cf. [6]. A spherical t-design X is defined by requiring that, for $k = 0, 1, \ldots, t$, the kth moments of X are constants with respect to orthogonal transformations of \mathbb{R}^d or, equivalently, that $\sum_{x \in X} W(\xi) = 0$ for all homogeneous harmonic polynomials $W(\xi)$ in \mathbb{R}^d of degree 1, 2, ..., t. The strength t(X) is the maximum value of t for which X is a t-design. We explain how spherical t-designs on the sphere, with the orthogonal group, correspond to classical t-designs on the discrete sphere, with the symmetric group. The cardinality of a spherical t-design X is bounded from below, and we find

$$n \ge {d+e-1 \choose d-1} + {d+e-2 \choose d-1}, \quad n \ge 2{d+e-1 \choose d-1},$$

for t = 2e and t = 2e + 1, respectively. The spherical *t*-design is called *tight* if any one of these bounds is attained.

Combining the notions introduced above, we consider in Section 6 spherical (d, n, s, t)-configurations X. These are sets X of cardinality n on the unit sphere Ω_d , which are spherical t-designs and spherical A-codes with |A| = s; in other words, the strength t(X) is at least t and the degree s(X) is at most s. A condition is given for a spherical A-code to be a spherical t-design, in terms of the Gegenbauer coefficients of an annihilator of the set A. This yields $t \leq 2s$, and $t \leq 2s - 1$ if $A \cup \{1\}$ is symmetric with respect to 0, as well as inequalities for |X|. The cases of equality are discussed. In Section 7 it is shown that $t \geq 2s - 2$ implies that X carries an s(X)-class association scheme. It is proved that there exist no tight spherical 6-designs, apart from the regular heptagon in \mathbb{R}^2 .

Sections 8 and 9 contain many examples of spherical (d, n, s, t)-configurations; there exist tight spherical *t*-designs with t = 2, 3, 4, 5, 7, 11, and nontight spherical (2s - 1)-designs. The constructions of these examples use sets of lines with few angles, and association schemes, respectively. In Section 8 the derived of a spherical *t*-design of degree s on Ω_{d+1} is shown to be a spherical (t + 1 - s)-design on Ω_d , and a spherical (t + 2 - s)-design on Ω_d in the antipodal case. Example 9.2 mentions a relation to the Krein condition.

As in [11], matrix techniques and the addition formula for Gegenbauer polynomials are our basic tools. These orthogonal polynomials are reviewed in Section 2. In Section 3 the inequalities of later sections are prepared, in terms of the Gegenbauer coefficients of certain polynomials, and of the characteristic matrices H_k defined on the set X and on an orthogonal basis of harmonic polynomials of degree k.

We use N for the set of the natural numbers including 0, and J for the all-one matrix.

2. GEGENBAUER POLYNOMIALS

We shall need the following family $\{Q_k(x); k \in \mathbb{N}\}$ of polynomials $Q_k(x)$ in one indeterminate x, defined for a fixed $d \ge 2$.

DEFINITION 2.1 The Gegenbauer polynomial $Q_k(x)$ of degree k is defined by

$$\lambda_{k+1}Q_{k+1}(x) = xQ_k(x) - (1 - \lambda_{k-1})Q_{k-1}(x),$$

$$\lambda_k = k/(d+2k-2), \qquad Q_0(x) = 1, \qquad Q_1(x) = dx.$$

The first few polynomials are:

$$2Q_2(x) = (d + 2)(dx^2 - 1),$$

$$6Q_3(x) = d(d + 4)((d + 2)x^3 - 3x),$$

$$24Q_4(x) = d(d + 6)((d + 2)(d + 4)x^4 - 6(d + 2)x^2 + 3),$$

$$120Q_5(x) = d(d + 2)(d + 8)((d + 4)(d + 6)x^5 - 10(d + 4)x^3 + 15x).$$

Remark 2.2. For $d \ge 3$ the present polynomials are related to the usual [1] Gegenbauer polynomials $C_k^m(x)$ by

$$Q_k(x) = \frac{d+2k-2}{d-2} C_k^{(d-2)/2}(x).$$

For $d = 2, k \ge 1$, they are related to the Chebyshev polynomials of the first kind $T_k(x)$ by

$$Q_k(x) = kC_k^0(x) = 2T_k(x).$$

DEFINITION 2.3.

$$R_{k}(x) := \sum_{i=0}^{k} Q_{i}(x),$$
$$C_{k}(x) := \sum_{i=0}^{\lfloor k/2 \rfloor} Q_{k-2i}(x).$$

Apart from a constant, $R_k(x)$ is the Jacobi polynomial $P_k^{(\mu+1,\mu)}(x)$, with $\mu = \frac{1}{2}(d-3)$, whereas $C_k(x)$ is the usual Gegenbauer polynomial $C_k^{d/2}(x)$. From the definitions the following theorem is easily proved.

THEOREM 2.4.

$$Q_{k}(1) = \binom{d+k-1}{d-1} - \binom{d+k-3}{d-1},$$

$$R_{k}(1) = \binom{d+k-1}{d-1} + \binom{d+k-2}{d-1},$$

$$C_{k}(1) = \binom{d+k-1}{d-1}, \quad \text{for } k \ge 1.$$

The Gegenbauer polynomials are orthogonal polynomials, that is,

$$\int_{-1}^{1} Q_k(x) Q_l(x) (1-x^2)^{(d-3)/2} dx = a_d Q_k(1) \delta_{k,l},$$

where a_d is some positive constant, and $\delta_{k,l}$ is the Kronecker symbol. To any polynomials F(x), $G(x) \in \mathbb{R}[x]$ we associate their Gegenbauer expansions

$$F(x) = \sum_{k=0}^{\infty} f_k Q_k(x), \qquad G(x) = \sum_{k=0}^{\infty} g_k Q_k(x),$$

for well-defined Gegenbauer coefficients f_k and g_k . We shall need the following lemmas, which readily follow from the definitions and the well-known properties of Gegenbauer polynomials, cf. [1], [2].

LEMMA 2.5. Let $Q_i(x)Q_j(x) = \sum_{k=0}^{i+j} q_k(i, j)Q_k(x)$. Then

$$q_0(i,j) = Q_i(1)\delta_{i,j} \quad and \quad q_k(i,j) \ge 0$$

for all i, j, k, with $q_k(i, j) > 0$ if and only if $|i - j| \le k \le i + j$ and $k \equiv i + j \pmod{2}$.

LEMMA 2.6. Let
$$G(x) = Q_l(x)F(x)/Q_l(1)$$
 for some $l \in \mathbb{N}$. Then
 $g_0 = f_l, \qquad (\forall_{k \in \mathbb{N}}(f_k \ge 0)) \Rightarrow (\forall_{k \in \mathbb{N}}(g_k \ge 0)).$

LEMMA 2.7. Let $G(x) = x^i F(x)$ for some $l \in \mathbb{N}$. Then, for each $k \in \mathbb{N}$, the number g_k is a convex linear combination, with strictly positive coefficients, of the numbers f_{k+l-2i} , for $i = 0, 1, ..., \min(l, \lfloor \frac{1}{2}(k+l) \rfloor)$.

Proof. By induction with respect to *l*. For l = 1 we have

$$g_k = \lambda_k f_{k-1} + (1 - \lambda_k) f_{k+1}.$$

3. HARMONIC POLYNOMIALS

Let Ω_d , with measure ω_d , denote the unit sphere in the Euclidean space \mathbb{R}^d of dimension d, endowed with the inner product \langle , \rangle . For any $k \ge 0$, let $\operatorname{Hom}(k) = \operatorname{Hom}_d(k)$ denote the linear space of all functions $V: \Omega_d \to \mathbb{R}$ which are represented by polynomials $V(\zeta) = V(\zeta_1, \ldots, \zeta_d)$, homogeneous of total degree k in the d variables ζ_i . Let $\operatorname{Harm}(k)$ denote the subspace of $\operatorname{Hom}(k)$ consisting of all functions represented by harmonic polynomials of degree k. Then $\operatorname{Harm}(k)$ is invariant under the orthogonal group O(d) of \mathbb{R}^d . Any function $V \in \operatorname{Hom}(k)$ can be uniquely written as

$$V(\zeta) = \sum_{i=0}^{lk/2} \langle \zeta, \zeta \rangle^i W_{k-2i}(\zeta), \qquad W_l \in \operatorname{Harm}(l).$$

Therefore, the following direct sum decompositions hold, cf. [2], [15].

THEOREM 3.1.

$$\operatorname{Hom}(k) = \sum_{i=0}^{\lfloor k/2 \rfloor} \operatorname{Harm}(k - 2i),$$

$$\operatorname{Hom}(k) \oplus \operatorname{Hom}(k - 1) = \sum_{i=0}^{k} \operatorname{Harm}(i).$$

The linear space of the second line consists of all functions on Ω_d repre-

sented by (not necessarily homogeneous) polynomials of total degree $\leq k$ in *d* variables. For the dimensions we have, cf. [15] and Theorem 2.4,

THEOREM 3.2.

dim Hom
$$(k) = C_k(1)$$
, dim Harm $(k) = Q_k(1)$,
dim Hom $(k) \oplus$ Hom $(k - 1) = R_k(1)$.

The addition formula relates the Gegenbauer polynomial $Q_k(x)$, and any orthogonal basis $\{W_{k,i}; i = 1, 2, ..., Q_k(1)\}$ of Harm(k), with norm $W_{k,i} = \omega_d^{1/2}$, as follows, cf. [1], [2], [15].

THEOREM 3.3.

$$\sum_{i=1}^{Q_k(1)} W_{k,i}(\xi) W_{k,i}(\eta) = Q_k(\langle \xi, \eta \rangle); \qquad \xi, \eta \in \Omega_d.$$

DEFINITION 3.4. For any finite non-empty set $X \subset \Omega_d$ of size *n*, for any orthogonal basis $\{W_{k,i}\}$ of Harm(k), with norm $W_{k,i} = \omega_d^{1/2}$, and for any fixed numbering of these, the $n \times Q_k(1)$ matrix

$$H_k := [W_{k,i}(\xi)], \quad \xi \in X, \, i \in \{1, 2, \dots, Q_k(1)\},$$

is called the kth characteristic matrix. Thus, H_0 is the all-one vector of size n.

DEFINITION 3.5. For any $X \subset \Omega_d$ of size *n*, and for any $\alpha \in \mathbb{R}$, $-1 \leq \alpha \leq 1$, the $n \times n$ distance matrix D_{α} is defined by its elements $D_{\alpha}(\xi, \eta) = 1$ for $\langle \xi, \eta \rangle = \alpha$, and $D_{\alpha}(\xi, \eta) = 0$ otherwise, for $\xi, \eta \in X$. The sum of the elements of D_{α} is denoted by d_{α} .

THEOREM 3.6. Let $X \subset \Omega_d$, and let A' be a finite set containing all inner products of the vectors of X. Then

$$H_k H_k^T = \sum_{\alpha \in A'} Q_k(\alpha) D_{\alpha},$$

where the $Q_k(x)$ are the Gegenbauer polynomials, H_k the characteristic matrices, and D_{α} the distance matrices.

Proof. The addition formula 3.3 and Definition 3.4 yield

$$H_k H_k^T = [Q_k(\langle \xi, \eta \rangle)]_{\xi, \eta \in X}.$$

Now apply Definition 3.5.

COROLLARY 3.7.

$$\|H_k^T H_0\|^2 = \sum_{\alpha \in A'} Q_k(\alpha) d_\alpha$$

Proof. In the formula of Theorem 3.6, take the sum of the elements of the matrices.

COROLLARY 3.8. For any polynomial F(x), with Gegenbauer coefficients f_0, f_1, \ldots , the following holds:

$$f_0 n^2 + \sum_{k=1}^{\infty} f_k \| H_k^T H_0 \|^2 = \sum_{\alpha \in A'} F(\alpha) d_{\alpha}.$$

Proof. Use $F(x) = \sum f_k Q_k(x)$ with Theorem 3.6 so as to obtain

$$\sum_{k=0}^{\infty} f_k H_k H_k^T = \sum_{\alpha \in A'} F(\alpha) D_{\alpha},$$

and take the sum of the elements of the matrices.

LEMMA 3.9.

$$||H_i^T H_j - n\Delta_{i,j}||^2 = \sum_{k=1}^{i+j} q_k(i,j) ||H_k^T H_0||^2,$$

where $q_k(i, j)$ is as in Lemma 2.5, and $\Delta_{i,j}$ denotes the appropriate zero matrix for $i \neq j$ and unit matrix for i = j.

Proof. We refer to [11], Lemma 4.5.

4. SPHERICAL CODES

DEFINITION 4.1. Let A be a subset of the interval [-1, 1[. A spherical A-code, for short an A-code, is a non-empty subset X of the unit sphere in \mathbb{R}^d , satisfying $\langle \xi, \eta \rangle \in A$, for all $\xi \neq \eta \in X$.

Thus, an A-code is a set of unit vectors with angles from the prescribed set arc cos A, or a set of points on Ω_d with distances from the prescribed set $(2(1 - A))^{1/2}$. We shall use the notation $A' := A \cup \{1\}$.

DEFINITION 4.2. A polynomial $F(x) \in \mathbb{R}[x]$ is compatible with the set A if $\forall_{\alpha \in A} (F(\alpha) \leq 0)$.

THEOREM 4.3. Let F(x), with Gegenbauer coefficients $f_0 > 0$ and $f_k \ge 0$, for all k, be compatible with the set A. Then the cardinality n of any A-code X satisfies

$$n \leq F(1)/f_0.$$

Equality holds if and only if, for all $\xi \neq \eta \in X$, and for all $k \ge 1$,

$$F(\langle \xi, \eta \rangle) = 0, \qquad f_k H_k^T H_0 = 0.$$

Proof. This is an immediate consequence of Corollary 3.8, since $d_1 = n$.

EXAMPLE 4.4. For given β , with $-1 \leq \beta < 0$, let A be any subset of the nterval $[-1, \beta]$. The polynomial $F(x) = x - \beta$ is compatible with A, and $f_0 = -\beta > 0$, $f_1 = 1/d > 0$. Hence Theorem 4.3 applies, yielding $n \leq 1 - 1/\beta$. An A-code of given dimension $r \leq d$ achieves this bound if and only if it is an r-dimensional regular simplex, with $\beta = -1/r$.

EXAMPLE 4.5. For given α and β , with

 $-1 \leq \alpha \leq \beta < 1$, $\alpha + \beta \leq 0$, $\alpha\beta > -1/d$,

let A be any subset of $[\alpha, \beta]$. Then the polynomial $F(x) = (x - \alpha)(x - \beta)$ has the Gegenbauer coefficients

$$f_0 = \alpha \beta + 1/d, \quad f_1 = -(\alpha + \beta)/d, \quad f_2 = 2/d(d + 2),$$

which are non-negative. Application of Theorem 4.3 yields

 $n \leq d(1-\alpha)(1-\beta)/(1+d\alpha\beta)$

for any A-code X. In addition, equality is only possible if A contains α and β , and if X is an { α , β }-code.

In the special case $\beta = -\alpha$, the lines spanned by the vectors of an $\{\alpha, \beta\}$ -code constitute a set of equiangular lines in \mathbb{R}^d . The sets achieving the above bound are equivalent to the regular 2-graphs, cf. [24], [26]. Also, the case of $\{\alpha, \beta\}$ -codes with $\alpha + \beta < 0$ leads to equiangular lines, now in \mathbb{R}^{d+1} . Indeed, cf. [19], since $\alpha + \beta < 0$ we may determine R and φ such that

$$1 - \alpha = R^{2}(1 - \cos \varphi), \quad 1 - \beta = R^{2}(1 + \cos \varphi), \quad R > 1, \\ 0 < \varphi < \pi.$$

Now define $Y \subseteq \mathbb{R}^{d+1}$ as follows

$$Y = \{R^{-1}((R^2 - 1)^{1/2}, \xi); \xi \in X\}.$$

Clearly, Y is a $\{\pm \cos \varphi\}$ -code on Ω_{d+1} . Hence it carries a set of equiangular lines with the angle φ .

EXAMPLE 4.6. For a given β , with $0 \le \beta < d^{-1/2}$, let A be any subset of $[-1, \beta]$. Define

$$\alpha := -(1+\beta)/(1+d\beta),$$

so $-1 \leq \alpha < 0$. The polynomial

$$F(x) := (x - \alpha)^2 (x - \beta)$$

is compatible with A, and has non-negative Gegenbauer coefficients with $f_0 > 0$. Application of Theorem 4.3 yields

$$n \leq d(1 - \beta)(2 + (d + 1)\beta)/(1 - d\beta^2)$$

for any A-code X, and equality is only possible if X is an $\{\alpha, \beta\}$ -code. It is interesting to observe that the bounds of Examples 4.5 and 4.6 coincide for this particular α .

EXAMPLE 4.7. For given α , β , γ , with $-1 \leq \alpha \leq \beta \leq \gamma < 1$, let A be any subset of $[-1, \alpha] \cup [\beta, \gamma]$. The polynomial $F(x) = (x - \alpha)(x - \beta)(x - \gamma)$ is compatible with A. It has non-negative Gegenbauer coefficients, with $f_0 > 0$, if

 $\begin{array}{l} \alpha+\beta+\gamma\leqslant 0, \qquad \alpha\beta+\beta\gamma+\gamma\alpha\geqslant -3/(d+2),\\ \alpha+\beta+\gamma<-d\alpha\beta\gamma. \end{array}$

Then Theorem 4.3 yields

$$n \leq -d(1-\alpha)(1-\beta)(1-\gamma)/(\alpha+\beta+\gamma+d\alpha\beta\gamma)$$

for any A-code X, and equality is only possible if X is an $\{\alpha, \beta, \gamma\}$ -code. In Example 9.3 we shall give a construction with

$$d = 23$$
, $n = 2048$, $\alpha = -9/23$, $\beta = -1/23$, $\gamma = 7/23$.

We conclude this section by giving yet another bound for the cardinality of an A-code X. This so-called absolute bound only depends on the cardinality of A, not on its specific elements.

THEOREM 4.8. For given $s = |A| < \infty$, the cardinality n of any A-code X satisfies $n \leq R_s(1)$.

Proof. Cf. [16]. For A we define the annihilator polynomial

$$F(x) := \prod_{\alpha \in A} (x - \alpha)/(1 - \alpha).$$

For any $\eta \in X$ we define the function $F_{\eta}: \Omega_d \to \mathbb{R}$ by

$$F_{\eta}(\xi) := F(\langle \xi, \eta \rangle), \qquad \xi \in \Omega_{d}.$$

Thus F_n belongs to the linear space Hom $(s) \oplus$ Hom(s - 1), which has the dimension $R_s(1)$. By definition we have

$$F_{\eta}(\xi) = \delta_{\xi,\eta}, \quad \text{for all} \quad \xi \in X,$$

so that the functions F_n are linearly independent. Hence their number n = |X| cannot exceed the dimension of the linear space, which proves the theorem.

EXAMPLE 4.9. For s = 1 we have $n \le d + 1$, with equality if and only if X is a regular d-simplex, as in Example 4.4.

EXAMPLE 4.10. For s = 2 we have $n \leq \frac{1}{2}d(d + 3)$. Examples meeting the bound exist for

$$d = 2, n = 5;$$
 $d = 6, n = 27;$ $d = 22, n = 275,$

Indeed, the following numbers of equiangular lines exist [17]:

6 in \mathbb{R}^3 ; 28 in \mathbb{R}^7 ; 276 in \mathbb{R}^{23} .

In each case we consider, with respect to a unit vector along any one line, the unit vectors at an obtuse angle along the other lines. These vectors determine a space of one dimension less, and provide the announced examples, cf. Example 4.5.

Theorem 4.8 has an application to graph theory.

THEOREM 4.11. A regular graph on n vertices, whose (0, 1)-adjacency matrix L has the smallest eigenvalue < -1 of multiplicity n - d, satisfies

$$n \leq \frac{1}{2}d(d+1) - 1.$$

Proof. Let k be the valency, and λ the smallest eigenvalue of L. Then

$$G := L - \lambda I - \frac{k - \lambda}{n} J$$

is positive semi-definite of rank $\leq d - 1$. Hence G is the Gram matrix of n vectors in \mathbb{R}^{d-1} of equal length with three distinct inner products. Theorem 4.8 applies with s = 2, hence

$$n \leq \binom{d}{d-2} + \binom{d-1}{d-2} = \frac{1}{2}d(d+1) - 1.$$

EXAMPLE 4.12. There are no strongly regular graphs with

$$n = 28, k = 18, \lambda_1 = 4, \lambda_2 = -2;$$

 $n = 276, k = 165, \lambda_1 = 27, \lambda_2 = -3,$

where, in both cases, k, λ_1 and λ_2 are the eigenvalues of L. This is well known for the 28-graph, but new for the 276-graph (Scott [23] showed the non-existence of rank 3 graphs with these parameters).

5. SPHERICAL DESIGNS

DEFINITION 5.1. A finite non-empty set $X \subseteq \Omega_d$ is a spherical t-design, for short a t-design, for some $t \in \mathbb{N}$ if the following holds for k = 0, 1, ..., t:

$$\forall_{V \in \operatorname{Hom}(k)} \forall_{T \in O(d)} \Big(\sum_{\xi \in X} V(T\xi) = \sum_{\xi \in X} V(\xi) \Big) \cdot$$

Here $T\xi$ denotes the image of $\xi \in \Omega_d$ under the element T of the orthogonal group O(d). Since Hom(k) is spanned by the monomials

$$\xi_1^{k_1}\xi_2^{k_2}\cdots\xi_d^{k_d}, \quad k_i\in\mathbb{N}, \qquad \sum_{i=1}^d k_i=k,$$

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Definition 5.1 amounts to requiring that the kth moments of X are constants with respect to orthogonal transformations, for k = 0, 1, ..., t. Thus, a 1-design is a set $X \subset \Omega_d$ whose centre of mass is the centre of Ω_d , and for a 2-design, in addition, the inertia ellipsoid is a sphere. Another way to express the t-design property requires that, for k = 0, 1, ..., t,

$$\forall_{V \in \operatorname{Hom}(k)} \forall_{T \in O(d)} \left(n^{-1} \sum_{\xi \in X} V(T\xi) = \omega_d^{-1} \int_{\Omega_d} V(\xi) d\omega(\xi) \right),$$

that is, that the kth moments of TX equal the corresponding kth moments of Ω_d , for all $T \in O(d)$. Since for any $V \in \text{Harm}(k)$, with $k \ge 1$, the above integral vanishes, Theorem 3.1 implies the following criterion for t-designs.

THEOREM 5.2. A finite set $X \subset \Omega_d$ is a t-design if and only if

$$\sum_{\xi \in X} W(\xi) = 0 \quad \text{for all} \quad W \in \sum_{k=1}^{t} \operatorname{Harm}(k).$$

THEOREM 5.3. A finite set $X \subset \Omega_d$ is a t-design if and only if its characteristic matrices satisfy any one of the following conditions:

(i) $H_k^T H_0 = 0$ for k = 1, 2, ..., t, or

(ii) $H_k^T H_l = n \Delta_{k,l}$, for $0 \le k + l \le t$.

Proof. The equivalence of Definition 5.1 and (i) follows from Theorem 5.2 and Definition 3.4. The equivalence of (i) and (ii) follows from Lemma 3.9.

Remark 5.4. For $t \ge 2$, let $e := \lfloor t/2 \rfloor$ and $r := e - (-1)^t$. Then

 $H_e^T H_e = nI$ and $H_e^T H_r = 0$

are necessary and sufficient conditions for a t-design. This is a consequence of Theorem 5.3 and Lemmas 2.5 and 3.9.

THEOREM 5.5. For any A-code X, let $A' = A \cup \{1\}$ and let d_{α} denote the sum of the elements of the distance matrix D_{α} . Then

$$\sum_{\alpha\in A'}d_{\alpha}Q_{k}(\alpha)\geq 0,$$

and equality holds for k = 1, 2, ..., t if and only if X is a t-design.

Proof. Apply Corollary 3.7 and Theorem 5.3.

EXAMPLE 5.6. Remark 5.4 says that 2-designs X are characterized by

$$H_0^T H_1 = 0 \quad \text{and} \quad H_1^T H_1 = nI.$$

Now observe that the $n \times d$ characteristic matrix H_1 satisfies

$$H_1H_1^T = d \sum_{\alpha \in A'} \alpha \dot{D}_{\alpha} = d \operatorname{Gram}(X).$$

Hence X is a 1-design if and only if the Gram matrix Gram(X) of its inner products has vanishing row sums. The second condition for X to be a 2-design amounts to the following: Gram(X) has two eigenvalues, namely n/d and 0, the vectors of X span \mathbb{R}^d and may be viewed as the orthogonal projections into \mathbb{R}^d of n orthogonal vectors of length $\sqrt{n/d}$ in n-space, cf. [6], [12]. Examples for such sets X are abundant; for instance, any spanning set of unit vectors along equiangular lines which correspond to a regular 2-graph. However, such a set only yields a 2-design if it is a 1-design; we refer to [24], [25] for many examples of this situation.

EXAMPLE 5.7. A set X is antipodal whenever

$$\forall_{\xi\in X}(-\xi\in X).$$

Obviously, antipodal A-codes provide 1-designs, since

$$A'(X) = -A'(X), \quad d_{\alpha} = d_{-\alpha}, \quad \sum_{\alpha \in A'} d_{\alpha} Q_k(\alpha) = 0 \text{ for odd } k.$$

The antipodal codes on Ω_d are in 1-1 correspondence with the sets of lines through the origin of \mathbb{R}^d , the subject of [11]. An antipodal code on Ω_d is a 3-design if and only if the Gram matrix of a spanning set of vectors of the corresponding set of lines has two eigenvalues. This yields many examples for 3-designs; for instance, the antipodal codes corresponding to regular 2-graphs, cf. [24], [26]. For d = 3 the six vertices of the octahedron, and also the eight vertices of the cube, provide a 3-design.

EXAMPLE 5.8. For a 5-design Theorem 5.5 requires

$$\sum_{\alpha \in A'} \alpha^i d_\alpha = 0 \quad \text{for} \quad i = 1, 3, 5,$$

and

$$\sum_{\alpha \in A'} \alpha^2 d_\alpha = \frac{n^2}{d}, \qquad \sum_{\alpha \in A'} \alpha^4 d_\alpha = \frac{3n^2}{d(d+2)}.$$

For d = 3, the 12 vertices of the icosahedron, and also the 20 vertices of the dodecahedron, provide a 5-design. Further examples are given in subsequent sections.

Remark 5.9. The analogy with the classical t-designs, cf. [14], [5], [27],

is explained as follows. For integers d, v with $1 \le d \le v/2$ we define the 'discrete d-sphere' in \mathbb{R}^v by

$$\Omega := \left\{ x = (x_1, \ldots, x_v) \in \mathbb{R}^v; x_i \in \{0, 1\}, \sum_{i=1}^v x_i = d \right\},\$$

whence $|\Omega| = {\binom{v}{d}}$. We define Hom(t) to be the set of all functions $f: \Omega \to \mathbb{R}$, which are represented by homogeneous polynomials f(x) of degree ≤ 1 in each coordinate x_t , and of total degree t. It turns out that the monomials

$$x_{i_1}x_{i_2}\cdots x_{i_t}$$

form a basis for Hom(t), hence dim Hom(t) = $\binom{v}{t}$. Now a classical t-design $t - (v, d, \lambda)$ is a collection X of d-subsets of a v-set, such that each t-subset is contained in a constant number λ of elements of X. In the setting above, this corresponds to a subset X of Ω subject to the condition

$$\forall_{f \in \operatorname{Hom}(t)} \forall_{T \in \operatorname{Sym}(v)} \Big(\sum_{\xi \in X} f(T\xi) = \sum_{\xi \in X} f(\xi) \Big) \cdot$$

This is equivalent to requiring that the sum over X of any monomial \in Hom(t) is a constant with respect to Sym(v), that is, that the monomial takes the value 1 for a constant number, λ say, of elements of X. Thus Ω , Sym(v), and the classical t-designs, correspond to Ω_d , O(d), and the spherical t-designs, respectively. This correspondence may be pushed still further; for details we refer to [9].

There is no upper bound to the number of points of a t-design, since the union of disjoint t-designs again is a t-design. The following theorem, which in some sense is dual to Theorem 4.3, provides a lower bound.

THEOREM 5.10. Let F(x), with Gegenbauer coefficients $f_0 > 0$ and $f_k \le 0$ for all k > t, satisfy F(1) > 0 and $F(\alpha) \ge 0$ for all $\alpha \in [-1, 1]$. Then the cardinality of any t-design X satisfies

 $n \geq F(1)/f_0.$

Equality holds if and only if, for all $\xi \neq \eta \in X$, and for all k > t,

 $F(\langle \xi, \eta \rangle) = 0, \qquad f_k H_k^T H_0 = 0.$

Proof. This is a consequence of Theorem 5.3 and Corollary 3.8.

THEOREM 5.11. Let X be a (2e)-design. Then

 $n=|X| \ge R_e(1).$

Equality holds if and only if A(X) consists of the zeros of $R_e(x)$.

Proof. Apply Theorem 5.10 for $F(x) = (R_e(x))^2$. It is easily verified, by use of the orthogonality relations for Gegenbauer polynomials and Theorem 2.4, that the bound specializes to

$$n \ge F(1)/f_0 = (R_e(1))^2/R_e(1) = R_e(1),$$

with equality if and only if all elements of A(X) are zeros of F(x). Now it readily follows from Theorem 4.8 that $n = R_e(1)$ implies $|A(X)| \ge e$, so that A(X) must consist of the *e* zeros of F(x).

THEOREM 5.12. Let X be a (2e + 1)-design, Then

$$n=|X|\geq 2C_e(1).$$

Equality holds if and only if A(X) consist of -1 and the zeros of $C_e(X)$. Moreover, in the case of equality X is antipodal.

Proof. Apply Theorem 5.10 for $F(x) = (x + 1)(C_e(x))^2$, then the lower bound specializes to $2C_e(1)$, and the desired result about A(X) is proved by an argument similar to that of Theorem 5.11. In order to prove the last statement, we define Y to be the set of the lines carried by the vectors of X. Clearly $|X| \leq 2|Y|$, with equality if and only if X is antipodal. If $|X| = 2C_e(1)$, then from $C_e(-x) = (-1)^e C_e(x)$ it follows that $(A(X))^2$ has cardinality [e/2] + 1 and contains 0 whenever e is odd. Therefore, the absolute bound for systems of lines [11] yields $|Y| \leq C_e(1)$, whence $|X| = 2|Y| = 2C_e(1)$, so that X is antipodal (and Y meets the absolute bound).

DEFINITION 5.13. A *t*-design is called *tight* if any of the bounds mentioned in Theorems 5.11 and 5.12 is attained.

Clearly, a tight t-design cannot be a (t + 1)-design. We conclude this section by some preliminary examples of tight t-designs.

EXAMPLE 5.14. For d = 2 and any t, a tight t-design is nothing but a regular (t + 1)-gon.

EXAMPLE 5.15. For any d, the d + 1 vertices of a regular simplex in \mathbb{R}^d provide a tight 2-design. The 2d vertices of the cross polytope (the generalization of the octahedron) provide a tight 3-design. Notice that the 2^d vertices of the cube also provide a 3-design (not a 4-design), but not a tight 3-design for $d \ge 3$.

EXAMPLE 5.16. For d = 3 the icosahedron is the only tight 5-design.

6. SPHERICAL (d, n, s, t)-CONFIGURATIONS

DEFINITION 6.1. A (spherical) (d, n, s, t)-configuration is a set $X \subset \Omega_d$ of cardinality *n*, which is a *t*-design and an *A*-code with |A| = s.

Given $X \subseteq \Omega^d$, |X| = n, we denote by s(X) and t(X) the minimum s and the maximum t for which X is a (d, n, s, t)-configuration. Theorem 6.5 will provide a criterion for an A-code to be a t-design, in terms of the Gegenbauer coefficients f_0, f_1, \ldots, f_s of an annihilator F(x) of degree s for the set A.

DEFINITION 6.2. $F(x) \in \mathbb{R}[x]$ is an *annihilator* polynomial for a finite set $A \neq \emptyset$ with $1 \notin A$ if

$$F(1) = 1, \quad \forall_{\alpha \in A}(F(\alpha) = 0).$$

LEMMA 6.3. Let X be an A-code, and let G(x) be an annihilator for A with Gegenbauer coefficients g_0, g_1, \ldots Then

$$n(1 - ng_0) = \sum_{k=1}^{\infty} g_k \| H_0^T H_k \|^2.$$

Proof. Apply Corollary 3.8.

THEOREM 6.4. Let X be an A-code, |X| = n, |A| = s. The Gegenbauer coefficients of an annihilator F(x) of degree s for A satisfy

$$(\forall_{0 \leq i \leq s} (f_i \geq 0)) \Rightarrow (\forall_{0 \leq j \leq s} (f_j \leq 1/n)).$$

If, in addition, $f_j = 1/n$ for some $j \leq s$, then X is an A-code of maximum cardinality.

Proof. For any fixed $j \in \{0, 1, \ldots, s\}$, define

 $G(x) := F(x)Q_j(x)/Q_j(1).$

Clearly, G(x) is an annihilator for A. Lemma 2.6 implies that $g_0 = f_j$ and $g_k \ge 0$ for all k. Hence Lemma 6.3 yields $1 - nf_j = 1 - ng_0 \ge 0$. If equality holds, then the bound of Theorem 4.3 is attained, and X is an A-code of maximum cardinality.

THEOREM 6.5. Let X be an A-code, with |X| = n, |A| = s, and let F(x) be an annihilator of degree s for A with Gegenbauer coefficients f_0, f_1, \ldots, f_s . If X is a t-design with $t \ge s$, then $f_0 = f_1 = \cdots = f_{t-s} = 1/n$. Conversely, if $f_0 = f_1 = \cdots = f_r = 1/n$, and $f_{r+1} > 0, \ldots, f_s > 0$ for some $r \le s$, then X is an (r + s)-design.

Proof. First, suppose X is a t-design with $t \ge s$. For any fixed j = 0, 1, ..., t - s the polynomial

$$G(x) := F(x)Q_j(x)/Q_j(1)$$

is an annihilator for A of degree $j + s \le t$. Hence Theorem 5.3 and Lemma 6.3 yield $0 = 1 - ng_0 = 1 - nf_j$, by use of Lemma 2.6. Conversely, let us consider the annihilator

 $G(x) := x^r F(x)$

for A of degree r + s. Assuming $f_0 = \cdots = f_r = 1/n$ and all $f_i > 0$, we conclude from Lemma 2.7 that $g_0 = 1/n$, $g_k > 0$ for $0 \le k \le r + s$. Lemma 6.3 implies $H_k^T H_0 = 0$ for $1 \le k \le r + s$, whence X is an (r + s)-design.

THEOREM 6.6. Any (d, n, s, t)-configuration X satisfies

 $t \leq 2s$ and $n \leq R_s(1)$.

If t = 2s, or if $n = R_s(1)$, then X is a tight (2s)-design.

Proof. Let F(x) be the annihilator of degree s for A. We first apply Theorem 6.5. If $t \ge s$, then $f_{t-s} \ne 0$, hence $t - s \le s$. This proves $t \le 2s$. In the case of equality Theorem 5.11 implies $n \ge R_s(1)$, whence $n = R_s(1)$ by Theorem 4.8, and X is a tight (2s)-design. For the second part of the theorem we observe that Theorem 3.6 implies

$$\sum_{k=0}^{s} f_k H_k H_k^T = I.$$

Hence the $n \times R_s(1)$ -matrix

$$H := [H_0 \quad H_1 \quad \cdots \quad H_s]$$

has rank *n*, proving once again $n \leq R_s(1)$, cf. Theorem 4.8. Now suppose $n = R_s(1)$, then *H* is non-singular, and all f_k are positive. Therefore Theorem 6.4 implies that all $f_k \leq 1/n$. Hence

$$n\sum_{k=0}^{s} f_k Q_k(1) = n = R_s(1) = \sum_{k=0}^{s} Q_k(1)$$

implies $f_0 = f_1 = \cdots = f_s = 1/n$, $nF(x) = R_s(x)$, and it follows from Theorem 6.5 that X is a (2s)-design. Now the theorem is proved. It is interesting to observe that in the case of equality we have

$$HH^{\mathrm{T}} = H^{\mathrm{T}}H = nI.$$

EXAMPLE 6.7. Tight 4-designs have

$$s(X) = 2$$
, $t(X) = 4$, $n = R_2(1) = \frac{1}{2}d(d+3)$.

Example 4.10 applies, and as a consequence of Theorem 6.6 we have three tight 4-designs, with

$$d = 2, n = 5;$$
 $d = 6, n = 27;$ $d = 22, n = 275.$

THEOREM 6.8. Any (d, n, s, t)-configuration X, which is an A-code with A' = -A', |A| = s, satisfies

$$t \leq 2s-1$$
 and $n \leq 2C_{s-1}(1)$.

If t = 2s - 1, or if $n = 2C_{s-1}(1)$, then X is an antipodal tight (2s - 1)-design. Proof. Applying the absolute bound [11] to the set Y of the lines carried by the vectors of X we obtain

$$n = |X| \leq 2|Y| \leq 2C_{s-1}(1).$$

This also proves $t \le 2s - 1$, since t = 2s is excluded by Theorem 6.6. If t = 2s - 1, then Theorem 5.12 implies $n \ge 2C_{s-1}(1)$, which yields $n = 2C_{s-1}(1)$, and X is antipodal. Now suppose $n = 2C_{s-1}(1)$, then by the above inequality X is antipodal and Y attains the absolute bound for sets of lines. It follows from [11], theorem 6.1, that the annihilator F(x) of degree s for A is given by

$$nF(x) = (1 + x)C_{s-1}(x) = \sum_{k=0}^{s-1} Q_k(x) + \lambda_s Q_s(x),$$

cf. Definitions 2.1 and 2.3. Therefore, Theorem 6.5 implies that X is a (2s - 1)-design. Now the theorem is proved. Sections 8 and 9 contain examples of (2s - 1)-designs which are not antipodal, hence not tight.

7. DISTANCE INVARIANCE AND ASSOCIATION SCHEMES

For any A-code X, the valencies $v_{\alpha}(\xi)$ and the intersection numbers $p_{\alpha,\beta}(\xi, \eta)$ are defined as follows.

DEFINITION 7.1.

$$\begin{aligned} &\forall_{\alpha \in A} \forall_{\xi \in \mathbb{X}} (v_{\alpha}(\xi) := |\{\zeta \in X : \langle \xi, \zeta \rangle = \alpha\}|), \\ &\forall_{\alpha, \beta \in A} \forall_{\xi, \eta \in \mathbb{X}} (p_{\alpha, \beta}(\xi, \eta) := |\{\zeta \in X : \langle \xi, \zeta \rangle = \alpha, \langle \eta, \zeta \rangle = \beta\}|). \end{aligned}$$

DEFINITION 7.2. X is distance invariant if, for all $\alpha \in A'$, the valency $v_{\alpha}(\xi)$ is independent of $\xi \in X$. X carries an s(X)-class association scheme if, for all $\alpha, \beta \in A'$, the intersection number $p_{\alpha,\beta}(\xi, \eta)$ depends only on $\langle \xi, \eta \rangle$.

Thus, the association schemes of Bose and Mesner, cf. [4], [7], [13], are specialized to the present situation. It is interesting to point out that any abstract association scheme can be represented by means of an A-code of a suitable dimension, cf. Section 9. We observe that the triangle inequality on the sphere imposes restrictions on the intersection numbers, namely

$$(p_{\alpha,\beta}(\xi,\eta)\neq 0) \Rightarrow (2(1-\alpha)(1-\beta)(1+\langle\xi,\eta\rangle))$$

$$\geq (1-\alpha-\beta+\langle\xi,\eta\rangle)^2).$$

For any integer $i \ge 0$, let xⁱ have the Gegenbauer expansion

$$x^i = \sum_{k=0}^i f_{i,k} Q_k(x).$$

The 'convolution' of x^i and x^j is defined to be the polynomial

$$F_{i,j}(x) := \sum_{k=0}^{\min(i,j)} f_{i,k} f_{j,k} Q_k(x)$$

LEMMA 7.3. For $0 \le i + j \le t$, and for fixed $\gamma := \langle \xi, \eta \rangle$, the intersection numbers $p_{\alpha,\beta}(\xi, \eta)$ of a (d, n, s, t)-configuration satisfy the linear equation

$$\sum_{\alpha\beta,\epsilon,A} \alpha^i \beta^j p_{\alpha,\beta}(\xi,\eta) = n F_{i,j}(\gamma) - \gamma^j - \gamma^i + \delta_{1,\gamma}$$

Proof. By use of Theorem 5.3, part (ii), the t-design property implies

$$\left(\sum_{k=0}^{i} f_{i,k} H_k H_k^T\right) \left(\sum_{k=0}^{j} f_{j,k} H_k H_k^T\right) = n \sum_{k=0}^{\min(i,j)} f_{i,k} f_{j,k} H_k H_k^T.$$

We rewrite this by use of the addition formula which by Theorem 3.6 reads

 $H_k H_k^T = [Q_k(\langle \xi, \eta \rangle)]_{\xi, \eta \in X}.$

Equate the (ξ, η) -entries on both sides of the formula above, and use the definition of $f_{i,k}$, then

$$\sum_{\alpha,\beta\in A'}\alpha^i\beta^j p_{\alpha,\beta}(\xi,\eta)=nF_{i,j}(\langle\xi,\eta\rangle).$$

This leads to the desired formula since, for $\langle \xi, \eta \rangle = \gamma$,

$$p_{\alpha,1}(\xi,\eta)=p_{1,\alpha}(\xi,\eta)=\delta_{\alpha,\gamma}.$$

THEOREM 7.4. Let X be a (d, n, s, t)-configuration. If $t \ge s - 1$, then X is distance invariant. If $t \ge 2s - 2$, then X carries an s(X)-class association scheme. If $t \ge 2s - 3$, then, for any fixed $\langle \xi, \eta \rangle = \gamma$, the intersection numbers $p_{\alpha,\beta}(\xi, \eta)$ are uniquely determined by $p_{\gamma,\gamma}(\xi, \eta)$.

Proof. Suppose $t \ge s - 1$, and apply Lemma 7.3 for j = 0, $\xi = \eta$, $\gamma = 1$:

$$\sum_{\alpha\in A} \alpha^{i} v_{\alpha}(\xi) = n F_{i,0}(1) - 1; \qquad 0 \leq i \leq s - 1.$$

This linear system of s equations with s unknowns $v_{\alpha}(\xi)$ has a Vandermonde, hence non-singular, matrix. Therefore, the valencies are uniquely determined, and are independent of ξ .

Next suppose $t \ge 2s - 2$. Now Lemma 7.3 yields a linear system of s^2 equations for $0 \le i, j \le s - 1$, with s^2 unknowns $p_{\alpha,\beta}(\xi, \eta)$. The matrix of this system is the direct product of two Vandermonde matrices, hence is non-singular. Therefore, for fixed $\gamma = \langle \xi, \eta \rangle$, the intersection numbers are uniquely determined. The third part of the theorem is proved analogously.

THEOREM 7.5. Any tight t-design carries an s-class association scheme, with s = [t/2].

Proof. Apply Theorems 5.11, 5.12 and 7.4.

Remark 7.6. For $t(X) \ge 2s(X) - 2$, the Bose-Mesner algebra of the association scheme is easily described [7]. It is generated by I and the pairwise orthogonal idempotent matrices

$$J_k := n^{-1}H_kH_k^T, \quad k = 0, 1, \dots, s(X) - 1.$$

This allows us to give explicit formulae for the eigenvalues of the association scheme, cf. Example 8.4 below.

We conclude this section with the following non-existence theorem.

THEOREM 7.7. The only tight 6-design is the regular heptagon in \mathbb{R}^2 .

Proof. For d = 2, the only tight t-designs are the regular (t + 1)-gons. So we have to show that the existence of a 6-design $X \subset \Omega_d$, with $d \ge 3$ and $n = R_3(1)$, leads to a contradiction. By Theorem 7.5 the set X is an A-code of degree s(X) = |A(X)| = 3 which carries a 3-class association scheme. For $d \ge 3$, the eigenvalues of the association scheme are integers, since for $d \ge 3$ the multiplicities $Q_k(1)$ are distinct. This implies that the elements α, β, γ of A(X) are rational. On the other hand, by Theorem 5.11, these α, β, γ are the zeros of the polynomial

$$R_3(x) = \frac{1}{6}d((d+2)(d+4)x^3 + 3(d+2)x^2 - 3(d+2)x - 3),$$

and it is not difficult to show that any rational zero of $R_3(x)$ is the inverse of an integer. We now use an argument devised by van Lint [18] in the theory of perfect codes. By straightforward verification it follows that

$$R_3(-1/(d+2)) > 0, \quad R_3(-1/(d+3)) < 0.$$

Hence $R_3(x)$ has a zero between -1/(d+2) and -1/(d+3), which obviously cannot be the inverse of an integer. This contradiction proves the theorem.

The present theorem, and Examples 4.9, 4.10 and 6.7, suggest the following:

CONJECTURE 7.8. There exist no tight (2e)-designs in Ω_d for $d \ge 3$ and $e \ge 3$.

8. Examples from sets of lines and derived configurations

The unit sphere $\Omega_{d+1} \subset \mathbb{R}^{d+1}$, with vectors $\zeta = (\varepsilon, \eta_1, \eta_2, \ldots, \eta_d) =: (\varepsilon; \eta)$ for short, is partitioned into spheres in parallel spaces of dimension d as follows.

$$\Omega_{d+1} = \bigcup_{-1 \leq \varepsilon \leq 1} \{ (\varepsilon, \xi \sqrt{1-\varepsilon^2}); \xi \in \Omega_d \}.$$

Let $Z \subseteq \Omega_{d+1}$ be any *B*-code containing e := (1, 0, 0, ..., 0) = (1; 0). Define $B^* := B \setminus \{-1\}$ and $s^* := |B^*|$.

DEFINITION 8.1. The derived code Z_{ε} of Z, with respect to e and to any $\varepsilon \in B^*$, is the set

$$Z_{\varepsilon} := \{ \xi \in \Omega_d : (\varepsilon, \, \xi \sqrt{1 - \varepsilon^2}) \in Z \}.$$

Clearly, Z_{ε} is an A-code on Ω_d with

$$A = \{(\beta - \varepsilon^2)/(1 - \varepsilon^2) : \beta \in B^*\}.$$

The following theorem is the spherical analogue of the Assmus-Mattson theorem on designs in codes [3], cf. also [8], theorem 5.3.

THEOREM 8.2. Let $Z \subset \Omega_{d+1}$, containing e, be a t-design and a B-code, with $1 \leq s^* \leq t+1$. Then any non-empty derived $Z_{\varepsilon} \subset \Omega_d$, with respect to e and to $\varepsilon \in B^*$, is a $(t + 1 - s^*)$ -design.

Proof. Define δ by $\delta = 1$ if $(-e) \in \mathbb{Z}$ and $\delta = 0$ if $(-e) \notin \mathbb{Z}$. For any r with $0 \leq r \leq t + 1 - s^*$, any $F_r \in \text{Hom}_d(r)$, and any k with $r \leq k \leq t$, define $G_{r,k} \in \text{Hom}_{d+1}(k)$ by

$$G_{r,k}(\zeta) = G_{r,k}(\varepsilon;\eta) := \varepsilon^{k-r} F_r(\eta) = \varepsilon^{k-r} (1-\varepsilon^2)^{r/2} F_r(\xi),$$

for $\xi \in \Omega_d$. Then

$$\sum_{\zeta\in\mathbb{Z}} G_{r,k}(\zeta) - G_{r,k}(e) - \delta G_{r,k}(-e) = \sum_{\varepsilon\in B^*} \varepsilon^{k-r} (1-\varepsilon^2)^{r/2} \sum_{\zeta\in\mathbb{Z}_{\varepsilon}} F_r(\zeta).$$

Any $T \in O(d)$ induces an orthogonal transformation of Ω_{d+1} fixing e which, applied to Z, leaves the left-hand side invariant, since Z is a *t*-design. Therefore, the right-hand side

$$\sum_{\varepsilon\in B^*} \varepsilon^{k-\tau} (1-\varepsilon^2)^{r/2} \sum_{\xi\in TZ_\varepsilon} F_r(\xi)$$

is independent of $T \in O(d)$. For $k = r, r + 1, ..., r - 1 + s^*$, this yields s^* equations for the s^* unknowns

$$\sum_{\xi\in TZ_{\varepsilon}}F_{\tau}(\xi), \qquad \varepsilon\in B^*.$$

Since the (essentially Vandermonde) determinant

$$\det[e^{k-r}(1-e^2)^{r/2}], \qquad e \in B^*, \ k \in \{r, r+1, \ldots, r-1+s^*\},$$

is non-zero, the unknowns are determined, that is, they are independent of $T \in O(d)$. This holds for any $r = 0, 1, ..., t + 1 - s^*$, and for any $F_r \in Hom_d(r)$. Therefore, any $Z_{\varepsilon} \neq \emptyset$ is a $(t + 1 - s^*)$ -design, and the theorem is proved.

Let Y be any finite non-empty set of lines through the origin of \mathbb{R}^{d+1} , and

let Z be the set of the intersections of the lines with the unit sphere Ω_{d+1} . Then Z is an antipodal B-code for a set B satisfying B' = -B', thus yielding an (d' = d + 1, n', s', t')-configuration. Such Z, and their derived Z_{ε} , are exposed in the following examples.

EXAMPLE 8.3. Let Y be equiangular in \mathbb{R}^{d+1} , so that Z is an antipodal Bcode with $B = \{-1, \varepsilon, -\varepsilon\}, 0 < \varepsilon < 1$. Clearly, Z is a 1-design. Now apply Example 4.7 and Theorem 6.5. If $\varepsilon^2 < 1/(d+1)$, then Z is a 2-design, whence a 3-design, if and only if the bound

$$n' \leq \frac{2(d+1)(1-\varepsilon^2)}{1-\varepsilon^2(d+1)}$$

is attained. This corresponds to the special bound for equiangular lines. Theorem 6.8 implies that Z is a 4-design, whence a tight 5-design, if and only if n' = (d + 1)(d + 2). This corresponds to the absolute bound for equiangular lines. In this case $\varepsilon = (d + 3)^{-1/2}$, and $1/\varepsilon$ is an integer if d > 3. Thus, the icosahedron, and the regular 2 graphs on 28 and on 276 vertices provide tight 5-designs with the parameters

$$(d', n', s', t') = (3, 12, 3, 5), (7, 56, 3, 5), (23, 552, 3, 5).$$

The known regular 2-graphs provide many (d', n', 3, 3)-configurations; there are several infinite series, cf. [24], [26].

The derived configuration Z_{ε} of Z is an A-code, with

$$A = \left\{-\frac{\varepsilon}{1-\varepsilon}, \frac{\varepsilon}{1+\varepsilon}\right\}$$

By Theorem 8.2 we find a 2-design

$$(d, n, s, t) = \left(d, \frac{d}{1 - \varepsilon^2(d+1)}, 2, 2\right)$$

corresponding to each regular 2-graph. This is a tight 4-design if and only if $\varepsilon = (d+3)^{-1/2}$. We only know the existence of the following cases:

$$(d, n, s, t) = (2, 5, 2, 4), (6, 27, 2, 4), (22, 275, 2, 4).$$

Each tight 4-design provides a maximal solution to the problem of $[-1, \beta]$ codes, with $\beta = \epsilon/(1 + \epsilon)$; indeed, the bound of Example 4.6 is achieved. We
point out that any tight 4-design necessarily is the derived of a tight 5-design.

EXAMPLE 8.4. Let Y be a set of lines in \mathbb{R}^{d+1} , each pair of which is either perpendicular or has a given angle arc cos ε . Then $Z \subseteq \Omega_{d+1}$ is an antipodal *B*-code with

$$B = \{-1, 0, \varepsilon, -\varepsilon\}, \qquad 0 < \varepsilon < 1.$$

By application of Theorems 6.4 and 6.5 it can be shown that, in the case $\varepsilon^2 < 3/(d+3)$, Z forms a 4-design, whence a 5-design, if and only if

$$n' = 2 \frac{(d+1)(d+3)(1-\varepsilon^2)}{3-(d+3)\varepsilon^2},$$

that is, if the set of lines meets the special bound. Moreover, Z is a 6-design, whence a tight 7-design, if and only if the set of lines meets the absolute bound, that is,

$$n' = \frac{1}{3}(d+1)(d+2)(d+3).$$

The $n \varepsilon^2(d + 5) = 3$ and $1/\varepsilon \in \mathbb{N} \setminus \{0, 1\}$. Thus, the known sets of lines provide two tight 7-designs with the parameters

$$(d', n', s', t') = (8, 240, 4, 7)$$
 and $(23, 4600, 4, 7)$.

Indeed, the first configuration corresponds to the root system E_8 with $\varepsilon = \frac{1}{2}$ (the Gosset polytope 4_{21} in \mathbb{R}^8 , cf. [6], p. 204); the second configuration corresponds to a subset of the Leech lattice with $\varepsilon = \frac{1}{3}$. Their construction is briefly indicated in Example 8.5.

By Theorem 8.2 the 5-designs Z yield derived 3-designs Z_{ε} , and the 7-designs Z yield derived 5-designs Z_{ε} . Restricting to the case $\varepsilon^2(d + 5) = 3$, $1/\varepsilon \in \mathbb{N} \setminus \{0, 1\}$, of the absolute bound, we find for Z_{ε} :

$$A = \left\{ \frac{\varepsilon}{1+\varepsilon}, \frac{-\varepsilon}{1-\varepsilon}, \frac{-\varepsilon^2}{1-\varepsilon^2} \right\},$$

$$n \prod_{\alpha \in A'} \frac{x-\alpha}{1-\alpha} = Q_0(x) + Q_1(x) + Q_2(x) + \frac{1-\varepsilon^2}{3-\varepsilon^2} Q_3(x),$$

$$(d, n, s, t) = \left(\frac{3-5\varepsilon^2}{\varepsilon^2}, \frac{3-5\varepsilon^2}{2\varepsilon^6}, 3, 5 \right).$$

By Theorem 7.4 these configurations carry 3-class association schemes. For $\varepsilon = \frac{1}{2}$ we have the tight 5-design with parameters (7, 56, 3, 5) which was met in Example 8.3. For $\varepsilon = \frac{1}{3}$ we have the 5-design (22, 891, 3, 5), a first example of a non-tight (2s - 1)-design. The eigenmatrix P and the multiplicity vector μ of its association scheme are as follows:

$$A' = \{-\frac{1}{2}, -\frac{1}{8}, \frac{1}{4}, 1\}$$

$$P = \begin{bmatrix} 42 & 512 & 336 & 1\\ -21 & -64 & 84 & 1\\ 9 & -16 & 6 & 1\\ -3 & 8 & -6 & 1 \end{bmatrix}, \quad \mu = \begin{bmatrix} 1\\ 22\\ 252\\ 616 \end{bmatrix}.$$

EXAMPLE 8.5. Following McKay [20] we consider the lattice generated by the integral linear combinations of the columns of

$$\frac{1}{2} \begin{bmatrix} 4I_k & C_k - I_k \\ O_k & I_k \end{bmatrix},$$

where C_k is a skew Hadamard matrix of order k with the constant diagonal $-I_k$. For k = 4 this is the Gosset lattice 5_{21} in \mathbb{R}^8 ; its 240 vertices of length $\sqrt{2}$ provide the first configuration of Example 8.4. For k = 12 this is the Leech lattice in \mathbb{R}^{24} . Its $2 * \binom{28}{5}$ vectors of length 2 provide an antipodal *B*-code $Z \subset \mathbb{R}^{24}$ with

$$B = \{-1, 0, \pm \frac{1}{2}, \pm \frac{1}{4}\},\$$

and the corresponding set Y of lines meets the absolute bound [11]. From Theorem 6.8 it follows that Z is a tight 11-design with

$$(d', n', s', t') = (24, 196560, 6, 11).$$

The derived configuration of Z with respect to $\varepsilon = \frac{1}{2}$ is the (23, 4600, 4, 7)configuration of Example 8.4. The derived configuration X of Z with respect to $\varepsilon = \frac{1}{4}$ has

$$A(X) = \{-\frac{3}{5}, -\frac{1}{3}, -\frac{1}{15}, \frac{1}{5}, \frac{7}{15}\}, \quad t(X) = 7, \\ (d, n, s, t) = (23, 23 * 2^{11}, 5, 7).$$

Theorem 7.4 implies that X is distance invariant, but it does not guarantee that X carries a 5-class association scheme. However, the third part applies. For instance, for all $\xi, \eta \in X$ with $\langle \xi, \eta \rangle = -\frac{3}{5}$, the intersection numbers $p_{\alpha,\beta}(\xi,\eta)$ only depend on α and β , since the triangle property implies $p_{\gamma,\gamma}(\xi,\eta) = 0$ for these ξ, η .

EXAMPLE 8.6. The regular polytope {3, 3, 5} in \mathbb{R}^4 , cf. [6], p. 153, and [12], is a configuration X with

$$A(X) = \left\{-1, 0, \pm \frac{1}{2}, \frac{\pm 1 \pm \sqrt{5}}{4}\right\}, \quad (d, n, s, t) = (4, 120, 8, 11).$$

It suffices to observe that the annihilator for A(X) of degree 8 has the expansion

$$120F(x) = Q_0(x) + Q_1(x) + Q_2(x) + Q_3(x) + \frac{4}{5}Q_4(x) + \frac{2}{3}Q_5(x) + \frac{3}{7}Q_6(x) + \frac{1}{4}Q_7(x) + \frac{1}{9}Q_8(x).$$

Hence Theorem 6.5 implies that X is a design of strength t(X) = 11.

9. EXAMPLES FROM ASSOCIATION SCHEMES

In order to obtain a further method of construction for (d, n, s, t)-configurations, we consider Bose-Mesner algebras [4], [7]. For any fixed $s \in \mathbb{N}$, let

$$D_0 = I, D_1, \ldots, D_s; \qquad \sum_{i=0}^s D_i = J,$$

be real non-zero symmetric matrices of size n, with entries $\in \{0, 1\}$, which generate an (s + 1)-dimensional linear algebra over \mathbb{R} . This algebra is called the *Bose-Mesner algebra*, for short BM algebra, of the *s*-class association scheme with adjacency matrices D_i . It is well known that any BM algebra \mathcal{A} is commutative, and admits a unique basis of mutually orthogonal idempotents

$$J_0 = n^{-1}J, J_1, \ldots, J_s,$$

cf. [7]. Clearly, any $D \in \mathscr{A}$ is positive semi-definite whenever it has nonnegative components with respect to this basis. Thus for any such D, with unit diagonal, of rank $d \ge 2$, there exists a $d \times n$ matrix C such that

$$C^{T}C = D = I + \alpha_{1}D_{1} + \cdots + \alpha_{s}D_{s},$$

for well-defined $\alpha_i \in \mathbb{R}$. The columns of C represent unit vectors in \mathbb{R}^d with mutual inner products α_i . If $\alpha_i \neq 1$, for all $i \ge 1$, then these columns are distinct, and constitute an A-code $X \subset \Omega_d$ with $A = \{\alpha_1, \ldots, \alpha_s\}$ and |X| = n.

EXAMPLE 9.1. Let \mathscr{A} be a 3-dimensional BM algebra, that is, the adjacency algebra of a strongly regular graph [7]. Let J_1 be one of the non-trivial minimal idempotents of \mathscr{A} , and let $d = \operatorname{rank}(J_1)$ be the corresponding multiplicity. If the given graph is not a ladder graph or its complement, the matrix $D = d^{-1}J_1$ satisfies the above requirements, that is,

$$D = I + \alpha D_1 + \beta (J - I - D_1)$$

for some $\alpha, \beta < 1$, only depending on the spectrum of the graph. Any set $X \subset \Omega_d$ which has D as its Gram matrix is an $\{\alpha, \beta\}$ -code of cardinality n. Using well-known identities concerning the spectrum of a strongly regular graph, one can easily verify, as a consequence of Theorem 4.3, that X is a 2-design with

$$n = d(1 - \alpha)(1 - \beta)/(1 + d\alpha\beta).$$

Therefore, if $\alpha + \beta \leq 0$ holds, X provides a maximal solution to the problem of $\{\alpha, \beta\}$ -codes, cf. Example 4.5. Conversely, let there be given an $\{\alpha, \beta\}$ -code $X \subset \Omega_d$, with $\alpha < \beta < -\alpha$, whose cardinality n achieves the bound of Example 4.5. It turns out that the annihilator of degree s = 2 for A satisfies $f_0 = 1/n, f_1 > 0, f_2 > 0$. Therefore, Theorems 6.5 and 7.4 imply that X is a 2-design, and carries a strongly regular graph. EXAMPLE 9.2. Which of the 2-designs of Example 9.1 are 3-designs? This question has an interesting relation to the *Krein condition*. Let J_0, J_1, \ldots, J_s be the basis of the mutually orthogonal idempotents of a BM algebra \mathcal{A} . The Hadamard product $J_i \circ J_j$, being a principal submatrix of the Kronecker product $J_i \otimes J_j$, has all its eigenvalues in the interval [0, 1], and belongs to \mathcal{A} . Therefore, the coefficients in

$$J_i \circ J_j = \sum_{k=0}^{s} q_{ij}^k J_k$$

satisfy $q_{ij}^k \ge 0$. This is the Krein condition for \mathcal{A} , cf.* [13], [22], [23] and also [7]. Now it turns out that in Example 9.1 the following conditions are equivalent:

$$f_1 \leq \frac{1}{n}, \quad d\alpha\beta + \alpha + \beta + 1 \geq 0, \quad q_{11}^1 \geq 0.$$

Since $f_1 = 1/n$ is a criterion for X to be a 3-design, we have the following elaboration of Example 9.1. A strongly regular graph with $\alpha + \beta < 0$ provides a 3-design if and only if $q_{11}^1 = 0$, in other words, cf. [7], if and only if its 'pseudo-dual' has no triangles. The first and second of the following examples are provided by the Clebsch graph and the Higman-Sims graph, which are 'dual' to their complements. The remaining examples are derived from the McLaughlin graph, cf. Example 6.7.

$$\begin{array}{ll} A = \{-\frac{3}{5}, \frac{1}{5}\}, & (d, n, s, t) = (5, 16, 2, 3), \\ A = \{-\frac{4}{11}, \frac{1}{11}\}, & (d, n, s, t) = (22, 100, 2, 3), \\ A = \{-\frac{1}{3}, \frac{1}{9}\}, & (d, n, s, t) = (21, 112, 2, 3), \\ A = \{-\frac{2}{7}, \frac{1}{7}\}, & (d, n, s, t) = (21, 162, 2, 3). \end{array}$$

Notice that such examples yield 3-designs with s(X) = 2 which are not tight (not antipodal).

EXAMPLE 9.3. Let Γ be the orthogonal complement of the binary Golay code, that is, the unique binary code of length d = 23, size n = 2048, with Hamming distances 8, 12, 16, cf. [10]. Mapping the Hamming cube into the unit sphere in the usual way, we obtain from Γ an A-code $X \subset \Omega_{23}$ with

$$A = \{-\frac{9}{23}, -\frac{1}{23}, \frac{7}{23}\}.$$

The Gegenbauer coefficients of the annihilator of degree 3 for A are easily checked to satisfy

$$0 < f_1 < f_3 < f_2 < f_0 = \frac{1}{n}$$

* The present simple proof of the Krein condition also occurs in N. Biggs, 'Automorphic Graphs and the Krein Condition', Geom. Dedic. 5, 117-127 (1976).

Hence Theorem 6.5 implies that X is a 3-design of strength t(X) = 3, and a maximal code, cf. Example 4.7. Although we know that X carries a 3-class association scheme [7], we cannot deduce this property from Theorem 7.4. In fact, there might exist a maximal A-code (with necessarily the same parameters d_{α} as the Golay code) which does not carry an association scheme. This example shows the difference between the cases $s \ge 3$ and s = 2 (cf. Example 9.1).

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Authors' address:

P. Delsarte, J. M. Goethals, M.B.L.E. Research Laboratory, *Brussels*, Belgium

J. J. Seidel, Technological University, *Eindhoven*, The Netherlands.

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