

A. Configuration model

• vertex set $[m]$, $\forall m \in \mathbb{N}$

• We consider a degree sequence $\underline{d} = (d_i)_{i \in [m]}$ with $d_i = \text{degree of vertex } i$, s.t

* $d_i \geq 1 \quad \forall i \in [m]$

* $\sum_{i \in [m]} d_i =: l_m$ even so that $E(G) = \frac{l_m}{2}$

$\forall G$ compatible with \underline{d}

• We would like to construct a RG with uniform probability over $G_m^{\underline{d}} = \{G \in \mathcal{G}_m : d_G(i) = d_i, \forall i \in [m]\}$

① $|G_m^{\underline{d}}| = ?$

② $\mathcal{G}_m^{\underline{d}} \neq \emptyset ?$

\hookrightarrow One can add assumptions on $\underline{d} = (d_i)_{i \in \mathbb{N}}$ so that \underline{d} is graphical $\Leftrightarrow \mathcal{G}_m^{\underline{d}} \neq \emptyset$

One can address problem ① by considering multigraphs over $[m]$:

$$\mathcal{M}_m^{\underline{d}} = \{G \text{ multigraph over } [m] : d_G(i) = d_i \forall i \in [m]\}$$

The measure over $\mathcal{M}_m^{\underline{d}}$ is constructed as follows:

• assign to each vertex i , d_i half-edges

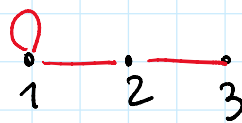
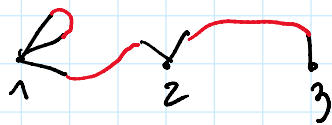
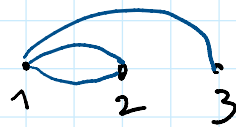
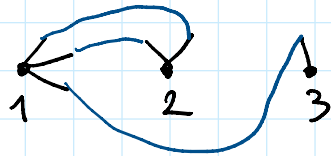
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1. Couple two half-edges to form an edge, iteratively (no matter of the order), until no more half-edges unpaired

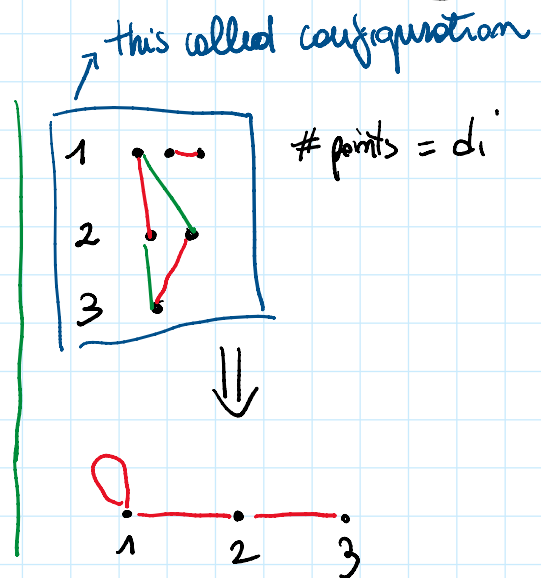
2. every matching is taken with equal probability

Example



This procedure can be visualized as:

- list all the vertices from 1 to n
- For each vertex i , draw d_i points
- Construct a matching on the total number of points = $2M$
- Collaps the d_i points into vertex i



Notice that # matching over $2M$ points is:

$$\frac{\binom{2M}{2} \binom{2M-2}{2} \cdots \binom{2}{2}}{M!} = (2M-1)!!$$

However $\phi: \text{Conf}(\underline{d}) \rightarrow \mathcal{M}_m^{\underline{d}}$ is surjective but not injective

hence the measure over $\mathcal{M}_m^{\underline{d}}$ that we obtain is not uniform. But it can be written explicitly.

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Denoting by $CM(m, \underline{d})$ the corresponding RG, we have

$$P_{m, \underline{d}}(G) := P(CM(m, \underline{d}) = G) = \frac{N(G)}{(2m-1)!!} \quad \forall G \in \mathcal{M}_m^{\underline{d}}$$

$$\begin{aligned} \text{where } N(G) &= \# \{ \text{configurations over } 2m \text{ mapped to } G \} \\ &= \prod_{i \in [m]} d_i! \cdot \frac{1}{2^{\# \text{ loops in } G} \cdot \prod_{e \in E(G)} \text{mult}(e)!} \quad \left(\begin{array}{l} \text{if } G \text{ is simple} \\ \prod d_i! \end{array} \right) \end{aligned}$$

B. Random Graph on $\mathcal{G}_m^{\underline{d}}$

Goal: define a measure over $\mathcal{G}_m^{\underline{d}}$. Starting $(\mathcal{M}_m^{\underline{d}}, P_{m, \underline{d}})$:

1^o METHOD: discard graphs that are not simple

2^o METHOD: erase loops and multiple-edges

→ this changes the degree sequence: $\underline{d} \rightarrow \underline{d}^{\text{er}}$

1. ASSUMPTIONS ON DEGREES:

let $\underline{d} = (d_i)_{i \in [m]}$ s.t. $\forall m \in \mathbb{N} : \sum_{i \in [m]} d_i =: 2m$ even

let $U \sim \text{Uniform}[m]$ and $D_m := d_U$. Then

a. \exists r.v. D s.t. $D_m \xrightarrow{d} D$. Equivalently:

$$P(D_m = k) = \frac{1}{m} \sum_{i \in [m]} \mathbb{1}_{\{d_i = k\}} \xrightarrow{m \rightarrow \infty} P(D = k) =: p_k$$

↳ If $(d_i)_{i \in [m]}$ are i.i.d r.v., this convergence (in probability or a.s.) is verified by LLN, with $D \stackrel{d}{=} d_i$.

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b. $E(D_m) \xrightarrow{m \rightarrow \infty} E(D)$ ($< \infty$ for having sparse RG)
 \hookrightarrow for $(d_i)_{i \in \mathbb{N}}$ iid r.v., this should understood as a convergence in probability

c. $E(D_m^2) \xrightarrow{m \rightarrow \infty} E(D^2)$

Remark: If we take r.v. $(d_i)_{i \in \mathbb{N}}$ iid st

$$P(d_i \geq k) \sim c \cdot k^{-(\tau-1)} \quad (\text{with } \tau > 2)$$

\Rightarrow The resulting graph have typical vertex degree that satisfies the corresponding power law decay (SCALE-FREE PROPERTY)

Consider $CM(m, \underline{d})$, where $\underline{d} = (d_i)_{i \in \mathbb{N}}$ satisfies assumptions a., b., c.

Theorem: $P(CM(m, \underline{d}) \stackrel{e \in G_m^{\underline{d}}}{\text{is simple}}) \xrightarrow{m \rightarrow \infty} e^{-\frac{\nu}{2} - \frac{\nu^2}{4}}$

where $\nu := \frac{E[D(D-1)]}{E(D)}$.

Proof (idea):

1st STEP: $S_m(G) := \# \text{ loops in } G$
 $M_m(G) := \# \text{ multiple edges in } G$ $\forall G \in \mathcal{M}_m^{\underline{d}}$

$\Rightarrow (S_m, M_m) \xrightarrow{d} (S, M)$

where S, M are independent: $S \sim \text{Poi}(\frac{\nu}{2})$, $M \sim \text{Poi}(\frac{\nu^2}{4})$

2nd STEP: $P(CM(m, \underline{d}) \text{ is simple}) = P_{m, \underline{d}}(S_m = 0, M_m = 0)$

Lemma. $P(\text{CM}(m, \underline{d}) \text{ is simple}) = P_{m, \underline{d}}(S_m=0, M_m=0)$
 $\approx P(S=0)P(M=0) = e^{-\frac{v}{2} - \frac{v^2}{4}}$ #

CONSEQUENCES:

1. Size of $G_m^{\underline{d}}$: Under assumptions a., b., c.:

$$|G_m^{\underline{d}}| = \frac{(m-1)!!}{\prod_{i \in [m]} d_i!} \cdot e^{-\frac{v}{2} - \frac{v^2}{4}} (1 + o(1)) \quad \text{as } m \rightarrow \infty$$

Proof: Let $G \in G_m^{\underline{d}} \subset \mathcal{M}_m^{\underline{d}} \Rightarrow P_{m, \underline{d}}(G) = \frac{\prod_{i \in [m]} d_i!}{(m-1)!!}$

$$e^{-\frac{v}{2} - \frac{v^2}{4}} (1 + o(1)) = P(\text{CM}(m, \underline{d}) \text{ is simple}) = \sum_{G \in G_m^{\underline{d}}} P_{m, \underline{d}}(G) = \frac{\prod_{i \in [m]} d_i!}{(m-1)!!} |G_m^{\underline{d}}|$$

Example: Take $d_i = r \quad \forall i \in [m]$ so that $v = \frac{r(r-1)}{2} = r-1$

$$\Rightarrow |G_m^r| = \frac{(rm-1)!!}{(r!)^m} \cdot e^{-\frac{r-1}{2} - \frac{(r-1)^2}{4}} (1 + o(1))$$

2. Under assumptions a., b., c., it holds:

Proposition: Let $E_m \subseteq G_m^{\underline{d}} \subset \mathcal{M}_m^{\underline{d}}$ s.t. $P(\text{CM}(m, \underline{d}) \in E_m) \xrightarrow{m \rightarrow \infty} 1$,

Let $G(m, \underline{d})$ be a uniform RG chosen from $G_m^{\underline{d}}$. Then

$$P(G(m, \underline{d}) \in E_m) \xrightarrow{m \rightarrow \infty} 1$$

Proof: By hyp. $P(\text{CM}(m, \underline{d}) \in E_m^c) \xrightarrow{m \rightarrow \infty} 0$.

$\Rightarrow P(\text{CM}(m, \underline{d}) \in E_m^c) = P(\text{CM}(m, \underline{d}) \in G_m^{\underline{d}} \cap E_m^c) = P(\text{CM}(m, \underline{d}) \in E_m^c) \xrightarrow{m \rightarrow \infty} 0$

$$\begin{aligned}
 \text{Then } \underline{P(G(m, d) \in E_m^c)} &= P(CM(m, d) \in E_m^c \mid CM(m, d) \text{ is simple}) \\
 &= \frac{P(CM(m, d) \in E_m^c \cap G_m^d)}{P(CM(m, d) \text{ is simple})} \leq \frac{P(CM(m, d) \in E_m^c)}{e^{-\frac{k}{2} - \frac{k^2}{4}}} \xrightarrow{m \rightarrow \infty} 0
 \end{aligned}$$

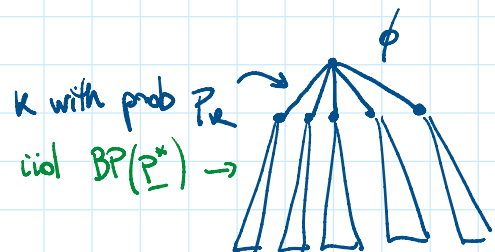
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[C.] Unimodular Branching Process

It is a random tree, similar to Galton-Watson BP(\underline{P}) but with a different branching rule between root and other vertices:

- the root has offspring distribution $\underline{P} = (P_k)_{k \geq 0} \stackrel{d}{\sim} X$
- a vertex \neq root has offspring distr. $\underline{P}^* = (P_k^*)_{k \geq 0}$ s.t.

$$P_k^* := \frac{(k+1) \cdot P_{k+1}}{\sum_{k \geq 0} k \cdot P_k} = \frac{(k+1) \cdot P_{k+1}}{\mathbb{E}(X)}$$



Hence:

$$\text{let } \eta := \sum_{k \geq 0} P_k^* \cdot \eta^k =: \mathcal{G}^*(\eta) \quad (\text{extinction probability of BP}(\underline{P}^*))$$

$$\text{Then } \mathcal{G} = \sum_{k \geq 0} P_k \cdot (1 - \eta^k)$$

• Notice that $\eta = 1 \iff \sum_{k \geq 0} k \cdot P_k^* \leq 1$

$$\iff \sum_{k \geq 0} \frac{k \cdot (k+1) P_k}{\mathbb{E}(X)} = \frac{\mathbb{E}(X(X-1))}{\mathbb{E}(X)} \leq 1$$

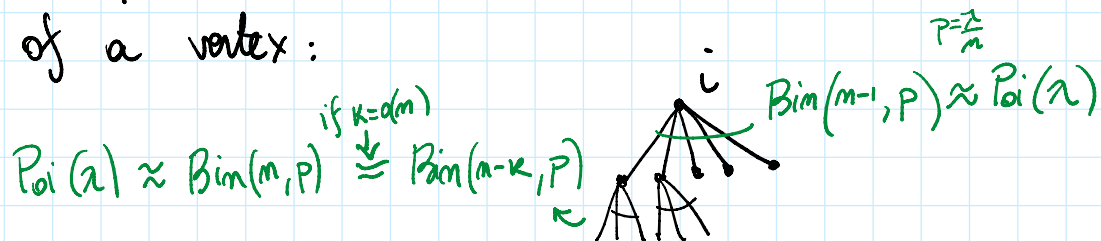
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In conclusion, it holds:

Theorem: if $\nu \leq 1 \Rightarrow \text{Unim BP}(\underline{P})$ dies a.s. ($\mathcal{L} = 0$)
 if $\nu > 1 \Rightarrow \text{Unim BP}(\underline{P})$ survives with positive prob. ($\mathcal{L} > 0$)

Informal explanation of unimodularity

- In ER random graph $G(m, p)$ if we look at the neighborhood of a vertex:

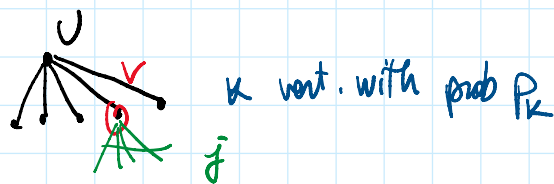


- The same fact holds for inhomogeneous RG
- In CM (m, \underline{d}) the degree sequence is assigned:

Let us pick a vertex $U \sim \text{uniform}[m]$ as a root-vertex

We know:

- $P(d_U = k) \approx P_k = P(D=k)$



- Let V vertex at distance 1 from U

$$P(\# \text{ of } \underset{V}{\text{spring}} = j) = P(d_V = j+1, \text{ and } V \text{ is connected to } U)$$

$$\stackrel{\text{Uniform matching}}{\approx} \frac{(j+1) \# \{v: d_v = j+1\}}{\sum_{l=1}^{\infty} l \cdot \# \{v: d_v = l\}} \cdot \frac{1}{m} \approx \frac{(j+1) \cdot P_{j+1}}{\sum_{l=1}^{\infty} l \cdot P_l} = P_j^*$$