

# Random Graphs and Networks - 6<sup>th</sup> LECTURE

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Recall: Given a BP ( $Poi(\lambda)$ ) (with offspring distr.  $Poi(\lambda)$ ):

1.  $\left\{ \begin{array}{l} \cdot \text{ if } \lambda \leq 1 \Rightarrow \text{BP}(Poi(\lambda)) \text{ dies out a.s. } (\eta = 1) \\ \cdot \text{ if } \lambda > 1 \Rightarrow \text{BP}(Poi(\lambda)) \text{ dies out with prob. } \eta < 1 \\ \text{and } \eta = e^{-\lambda(\eta-1)} \quad \left[ \eta = 1 - e^{-\lambda\eta} \right] \quad \eta = 1 - \eta \end{array} \right.$
2. if  $\lambda > 1$ , then  $\text{BP}(Poi(\lambda)) \setminus \{\text{dying out}\} \stackrel{d}{=} \text{BP}(Poi(\mu_\lambda))$   
 where  $\mu_\lambda$  is conjugate of  $\lambda$ :  $\mu_\lambda e^{-\mu_\lambda} = \lambda e^{-\lambda}$  (duality)
3. Exploration process on BP ( $Poi(\lambda)$ ):  
 Size of total progeny  $T \stackrel{d}{=} \min\{n : S_n = 0\}$

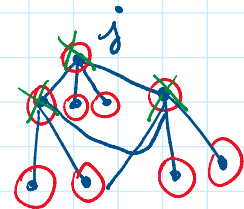
## A. Exploration of the components of $G(n, p)$ [ $p = \frac{\lambda}{n}, \lambda > 0$ ]

• Let  $j \in [n]$  and explore  $\mathcal{C}(j) = \{k \in [n] : j \leftrightarrow k\}$

Procedure - It is a dynamics along which vertices may be live, dead or neutral:

- At step  $k=0$ :  $j$  is live, while all vertices  $i \neq j$  are neutral
- At step  $k \geq 1$ , we choose a live vertex  $i$ , then

- \*  $i$ : live  $\rightarrow$  dead
- \* If  $l \sim i$ :  $\left\{ \begin{array}{l} \text{neutral} \rightarrow \text{live} \\ \text{live} \rightarrow \text{live} \\ \text{dead} \rightarrow \text{dead} \end{array} \right.$
- \* Vertices  $l \neq i$  keep their value



Then set  $S_k = \#$  live vertices at step  $k$ ,  $\forall k \geq 0$

$D_k = \#$  dead vertices at step  $k$

$\hookrightarrow D_k = k$

$N_k = \#$  neutral vertices at step  $k$

with  $S_k + N_k = m - k$

Remark:  $|\mathcal{C}(j)| =: T = \min\{k : S_k = 0\}$

1. Let  $X_i = \#$  vertices becoming live at step  $i$ ,  $\forall i \in \mathbb{N}$ :

$$S_k = S_{k-1} + X_k - 1 = \dots = \sum_{i=1}^k X_i - (k-1)$$

with  $S_0 = 1$

Hence  $(S_k)_{k \in \{0, \dots, T\}}$  is a RW with increments  $(X_{i-1})_{i \in \mathbb{N}}$

1<sup>st</sup> FACT:  $X_i \sim \text{Bin}(N_{i-1}, p)$



2<sup>nd</sup> FACT:  $N_i \sim \text{Bin}(m-1, (1-p)^i)$



$\text{Bin}(m-1, 1 - (1-p)^k)$

$$\begin{aligned} \Rightarrow S_k &= m - k - N_k = \underbrace{(m-1)}_{\text{Bin}(m-1, 1 - (1-p)^k)} - N_k - k + 1 \\ &\stackrel{d}{=} \text{Bin}(m-1, 1 - (1-p)^k) - k + 1 \end{aligned}$$

Heuristics: If  $i$  is "small" (w.r.t.  $m$ ) then:

$$N_i \approx m-1. \text{ Then } X_i \approx \text{Bin}(m-1, p) \underset{p = \frac{\lambda}{m}}{\approx} \text{Poi}(\lambda).$$

But we know that  $\text{BP}(\text{Poi}(\lambda))$  undergoes a phase transition w.r.t. to  $\lambda$   $\left( \begin{array}{l} \text{if } \lambda \leq 1 \Rightarrow \text{a.s. finite} \\ \text{if } \lambda > 1 \Rightarrow \text{infinite with prob } \xi \end{array} \right)$

## B. Phase transition

Recall  $|\mathcal{C}_{\max}| = \text{size of maximal component in } \mathcal{G}(m, p)$ ,

wh.  $\mathbb{P}(|\mathcal{C}_{\max}| > 1) > 0$  iff  $\lambda > 1$  and  $\xi > 0$ .

Recall  $|C_{\max}|$  = size of maximal component in  $G(m, p)$ ,  
 while  $|C_2|$  = size of the second maximal component.

Theorem: let  $p = \frac{\lambda}{m}$ . Then

1. if  $\lambda < 1$ :  $\frac{|C_{\max}|}{\lg m} \xrightarrow{d, P} c_\lambda = (\lambda - 1 - \lg \lambda)^{-2} = \left[ I_\lambda(1) \right]^{-2}$   
 ( $|C_{\max}| \sim c_\lambda \lg m$ )

2. if  $\lambda > 1$ : i.  $\frac{|C_{\max}|}{m} \xrightarrow{d, P} \zeta$  (where  $\zeta$  is sol:  
 $\zeta = 1 - e^{-\lambda \zeta}$ )  
 ( $|C_{\max}| \sim \zeta \cdot m$ )

ii.  $\frac{|C_2|}{m} \xrightarrow{d, P} 0$  (hence  $C_{\max}$  is unique)

Comments:

- $C_{\max}$  is called giant component
- $G(m, p)$  with  $p = \frac{\lambda}{m}$ ,  $\lambda > 1$  is highly connected
- This abrupt change while passing  $\lambda = 1$  is called phase transition and:

- $\lambda < 1$  is called subcritical
- $\lambda = 1$  is called critical
- $\lambda > 1$  is supercritical

- Result 2. ii. can be improved to show that  $\exists c > 0$

$$P(|C_2| \leq c \lg m) \xrightarrow{m \rightarrow \infty} 1$$

C. Tools: Chernoff bounds

Recall the Markov inequality:

if  $X \geq 0$  r.v.,  $a > 0$ :  $P(X \geq a) \leq \frac{E(X)}{a}$

Chernoff bounds are obtained as an application of this bound, that applies to general  $X$  r.v. (with finite average):

$$* a > \mathbb{E}(X) : P(X \geq a) = P(e^{Xt} \geq e^{at}) \stackrel{M.I.}{\leq} e^{-at} \mathbb{E}(e^{Xt})$$

$$* a < \mathbb{E}(X) : P(X \leq a) = P(e^{-Xt} \geq e^{-at}) \stackrel{M.I.}{\leq} e^{at} \mathbb{E}(e^{-Xt})$$

Thus, minimizing over  $t > 0$ , we get

$$P(X \geq a) \leq \inf_{t > 0} \left\{ e^{-at} \mathbb{E}(e^{Xt}) \right\} = e^{-\sup_{t > 0} \{ at - \lg \mathbb{E}(e^{Xt}) \}}$$

• Let  $(X_k)_{k \in \mathbb{N}}$  iid r.v. s.t.  $\mathbb{E}(X_k) = \mu < \infty$

$$\text{Let } S_m = \sum_{k=1}^m X_k \Rightarrow \mathbb{E}(S_m) = m \cdot \mu$$

For longer  $m$  (by LLN + CLT):  $S_m \approx m\mu + O(\sqrt{m})$

$$\text{If } a > \mu : P(S_m \geq am) = P(S_m \geq m\mu + (a-\mu) \cdot m)$$

Theorem [Chernoff bounds]

$$\bullet \text{ If } a > \mu \Rightarrow P(S_m \geq am) \leq e^{-mI(a)}$$

$$\bullet \text{ If } a < \mu \Rightarrow P(S_m \leq am) \leq e^{-mI(a)}$$

where  $I: \mathbb{R} \rightarrow \bar{\mathbb{R}}$  given by  $I(a) = \sup_{t \in \mathbb{R}} \{ at - \lg \mathbb{E}(e^{tX_0}) \}$

Comments:

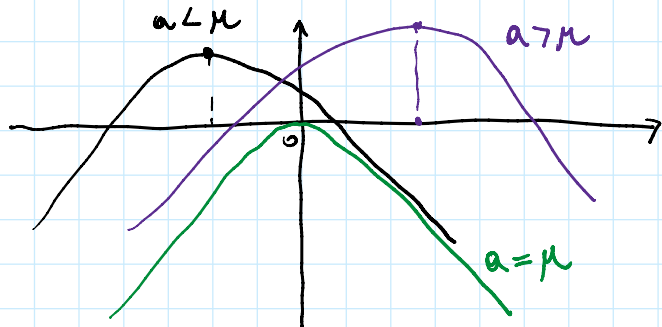
•  $I$  is the Legendre transform of  $\lg \mathbb{E}(e^{tX_1})$

• Let  $g(t) = t \cdot a - \lg \mathbb{E}(e^{tX_1})$ . It turns out that

\*  $g(t)$  is concave

$$* g'(t) \Big|_{t=a} = a - \mathbb{E}(X_1) = \begin{cases} < 0 & \text{if } a < \mu \\ = 0 & \text{if } a = \mu \end{cases}$$

$$* g'(t)|_{t=0} = a - \mathbb{E}(X_1) = \begin{cases} = 0 & \text{if } a = \mu \\ > 0 & \text{if } a > \mu \end{cases}$$



- ↓
- if  $a = \mu$ , the bound on useless
  - if  $a \neq \mu$ , then  $I(a) > 0$ .

Proof: Application of the exponential Markov inequality:  

$$\mathbb{E}[e^{tS_m}] = \prod_{j=1}^m \mathbb{E}[e^{tX_j}] = e^{m \log \mathbb{E}(e^{tX_1})}$$
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Examples:

1. Let  $S_m \sim \text{Poi}(m\lambda)$ .  $S_m \stackrel{d}{=} \sum_{k=1}^m X_k$ ,  $(X_k) \text{ iid } \sim \text{Poi}(\lambda)$

Recall  $\mathbb{E}[e^{tX_1}] = e^{-\lambda(1-e^t)}$ . Hence:

$$I_\lambda(a) = \sup_{t \in \mathbb{R}} \{ta + \lambda(1-e^t)\} = \lambda - a - a \log \frac{\lambda}{a}$$

Notice that  $I_\lambda(1) = \lambda - 1 - \log \lambda$

$$\Rightarrow \text{if } a > \lambda: P(S_m \geq am) \leq e^{-m I_\lambda(a)}$$

2. Let  $S_m \sim \text{Bin}(m, p)$ .  $S_m \stackrel{d}{=} \sum_{k=1}^m X_k$ , with  $(X_k) \text{ iid } \sim \text{Be}(p)$ .

$$\text{But } \mathbb{E}[e^{tX_1}] = pe^t + (1-p) = 1 + p(e^t - 1) \left( \leq e^{-p(1-e^t)} \right)$$

hence

$$I(a) = \sup_{t \in \mathbb{R}} \{ta - \log(1 + p(e^t - 1))\} \geq I_p(a)$$

↳ rate function of  $\text{Poi}(p)$

$$= a \log \frac{p}{a} + (1-a) \log \frac{1-p}{1-a}$$

$$= a \log \frac{p}{a} + (1-a) \log \frac{1-p}{1-a}$$

↳ rate function of  $\text{Poi}(p)$

D. Sub-critical regime ( $p = \frac{\lambda}{m}$ ,  $\lambda < 1$ )

The proof amounts to show:

- Upper bound:  $\mathbb{P}_{m,p}(|\mathcal{C}_{\max}| \geq c \log m) \rightarrow 0$  if  $c > c_\lambda$
- Lower bound:  $\mathbb{P}_{m,p}(|\mathcal{C}_{\max}| \leq c' \log m) \rightarrow 0$  if  $c' < c_\lambda$

Upper bound: let  $c > c_\lambda$

$$\mathbb{P}_{m,p}(|\mathcal{C}_{\max}| \geq c \log m) \leq \sum_{j \in [m]} \mathbb{P}_{m,p}(|\mathcal{C}(j)| > c \log m) = \theta(1)$$

GOAL:  $\mathbb{P}_{m,p}(|\mathcal{C}(j)| > c \log m) = \theta(m^{-1})$

if  $(S_k)_{k=1, \dots, T}$  is the exploration process of  $\mathcal{C}(j)$ :

$$\mathbb{P}_{m,p}(|\mathcal{C}(j)| > c \log m) = \mathbb{P}_{m,p}(T > c \log m)$$

• For given  $k = k(m) = \theta(m)$ , we have

$$\begin{aligned} \mathbb{P}_{m,p}(T > k) &= \mathbb{P}_{m,p}(S_k \geq 1) = \mathbb{P}(\text{Bin}(m-1, 1 - (1-p)^k) \geq k) \\ &\approx \mathbb{P}(\text{Poi}(k\lambda) \geq k+1) \stackrel{\lambda < 1}{\leq} e^{-k \cdot I_\lambda(1)} = m^{-\underbrace{c I_\lambda(1)}_{> 1}} = \theta(m^{-1}) \end{aligned}$$

$k \cdot p = \frac{k\lambda}{m}$

$\hookrightarrow$  Chernoff bounds

Taking  $k = c \log m$  we get

Remark: this argument breaks:

- when  $c \leq c_\lambda$  ( $\Rightarrow \mathbb{P}(T > k) \leq m^{-a}$ ,  $a \leq 1$ )

- when  $C \leq C_{\infty}$  ( $\Rightarrow P(T > k) \leq m^{-a}$ ,  $a \leq 1$ )
- when  $R = \gamma \cdot m$  ( $\Rightarrow 1 - (1-p)^k \sim 1 - e^{-\lambda k}$ )
- when  $\lambda > 1$  (Chernoff bound)