



## Rational Krylov methods for $f(A)\mathbf{b}$

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## Problem

We consider the vector  $f(A)\mathbf{b}$ , where

- ▶  $A$  is a large  $N$ -by- $N$  matrix,
- ▶  $\mathbf{b}$  is a vector of length  $N$ ,
- ▶  $f$  is a suitable function.

Compute approximation  $\mathbf{f}_m \approx f(A)\mathbf{b}$  from a rational Krylov space.

## What is a rational Krylov space?

Let  $\{\xi_1, \xi_2, \dots\} \subseteq \Xi \subset \bar{\mathbb{C}}$  be a given sequence of **poles**.

Define the polynomials

$$q_m(z) = \prod_{\substack{j=1 \\ \xi_j \neq \infty}}^m (z - \xi_j) \in \mathcal{P}_m.$$

Assume that  $q_m(A)^{-1}$  exists. Then

$$\mathcal{Q}_{m+1}(A, \mathbf{b}) = \mathcal{K}_{m+1}(A, q_m(A)^{-1}\mathbf{b})$$

is the rational Krylov space associated with  $(A, \mathbf{b}, q_m)$ .

## Special cases

- ▶  $\Xi = \{\infty\} \Rightarrow$  polynomial Krylov  $\mathcal{Q}_{m+1} = \mathcal{K}_{m+1}$

[Nauts & Wyatt 83] [van der Vorst 87] [Druskin & Knizhnerman 88]

[Gallopoulos & Saad 92] [Hochbruck & Lubich 97] [Eiermann & Ernst 06]

- ▶  $\Xi = \{\xi\} \Rightarrow$  shift-invert Krylov

[Moret & Novati 04] [van den Eshof & Hochbruck 06]

- ▶  $\Xi = \{0, \infty\} \Rightarrow$  extended Krylov

[Druskin & Knizhnerman 98] [Knizhnerman & Simoncini 08]

- ▶  $\Xi$  arbitrary  $\Rightarrow$  rational Krylov

[Ruhe 84] [Beattie 04] [Beckermann & Reichel 08] [Knizhnerman et al 08]

## Rational Arnoldi algorithm [Ruhe 84/94]

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**Input:**  $A$ ,  $\mathbf{b}$ ,  $\{\xi_1, \xi_2, \dots, \xi_m\}$

$\mathbf{v}_1 := \mathbf{b}/\|\mathbf{b}\|$

**for**  $j = 1, 2, \dots, m$  **do**

$\mathbf{x} := (I - A/\xi_j)^{-1}A\mathbf{v}_j$

$H(1:j, j) := [\mathbf{v}_1, \dots, \mathbf{v}_j]^* \mathbf{x}$

$\mathbf{x} := \mathbf{x} - [\mathbf{v}_1, \dots, \mathbf{v}_j]H(1:j, j)$

$H(j+1, j) := \|\mathbf{x}\|$

$\mathbf{v}_{j+1} := \mathbf{x}/H(j+1, j)$

**end**

---

Yields decomposition  $AV_{m+1}(\underline{I}_m + \underline{H}_m X_m^{-1}) = V_{m+1} \underline{H}_m$ .

## Rational Krylov decompositions

**Theorem (G., 2009):** Let a general decomposition

$$AV_{m+1}\underline{K}_m = V_{m+1}\underline{H}_m$$

be given, where  $V_{m+1}$  has  $m + 1$  linearly independent columns,  $\underline{K}_m \in \mathbb{C}^{(m+1) \times m}$ ,  $\underline{H}_m \in \mathbb{C}^{(m+1) \times m}$ , and  $\underline{H}_m$  is of rank  $m$ .

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1.  $\underline{K}_m$  is of rank  $m$ .
2.  $\text{colspan}(V_{m+1}) = \mathcal{K}_{m+1}(A, \mathbf{q})$  for a vector  $\mathbf{q}$ .
3. For every vector  $\mathbf{b} \in \text{colspan}(V_{m+1})$  there exists a unique polynomial  $q_m$ ,  $\deg(q_m) \leq m$ , such that  $\mathbf{b} = q_m(A)\mathbf{q}$ .  
Hence, if  $q_m(A)$  is invertible,  $\text{colspan}(V_{m+1}) = \mathcal{Q}_{m+1}(A, \mathbf{b})$ .

## Rational Krylov approximations

A special case is the (reduced) decomposition

$$AV_m K_m = V_{m+1} \underline{H}_m.$$

As an approximation to  $f(A)\mathbf{b}$  we consider

$$\mathbf{f}_m := V_m f(H_m K_m^{-1}) V_m^\dagger \mathbf{b}.$$

**Theorem (Interpolation):** There holds

$$\mathbf{f}_m = r_m(A)\mathbf{b} = \frac{p_{m-1}}{q_{m-1}}(A)\mathbf{b},$$

where  $r_m$  Hermite-interpolates  $f$  at  $\Lambda(H_m K_m^{-1})$ .

---



## Example

The iteration

$$\mathbf{v}_1 = \mathbf{b},$$

$$\beta_j \mathbf{v}_{j+1} = (I - A/\xi_j)^{-1} (A - \alpha_j I) \mathbf{v}_j, \quad j = 1, \dots, m,$$

yields a decomposition  $AV_{m+1} \underline{K}_m = V_{m+1} \underline{H}_m$  with

$$V_{m+1} = [\mathbf{v}_1, \dots, \mathbf{v}_{m+1}],$$

$$\underline{K}_m = \begin{bmatrix} 1 & & & & \\ \beta_1/\xi_1 & 1 & & & \\ & \beta_2/\xi_2 & \ddots & & \\ & & \ddots & 1 & \\ \hline & & & & \beta_m/\xi_m \end{bmatrix} \quad \text{and} \quad \underline{H}_m = \begin{bmatrix} \alpha_1 & & & & \\ \beta_1 & \alpha_2 & & & \\ & \beta_2 & \ddots & & \\ & & \ddots & \alpha_m & \\ \hline & & & & \beta_m \end{bmatrix}.$$

## Example

The iteration

$$\mathbf{v}_1 = \mathbf{b},$$

$$\beta_j \mathbf{v}_{j+1} = (I - A/\xi_j)^{-1}(A - \alpha_j I)\mathbf{v}_j, \quad j = 1, \dots, m,$$

can be used for explicit rational interpolation:

By Theorem (Interpolation) we know that

$$\mathbf{f}_m = V_m f(H_m K_m^{-1}) \mathbf{e}_1 = r_m(A) \mathbf{b} = \frac{p_{m-1}}{q_{m-1}}(A) \mathbf{b},$$

where  $r_m$  Hermite-interpolates  $f$  at  $\Lambda(H_m K_m^{-1}) = \{\alpha_1, \dots, \alpha_m\}$ .

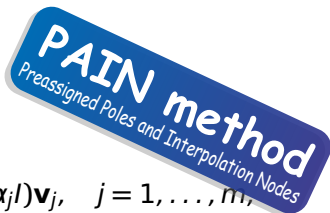
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## Example

The iteration

$$\mathbf{v}_1 = \mathbf{b},$$

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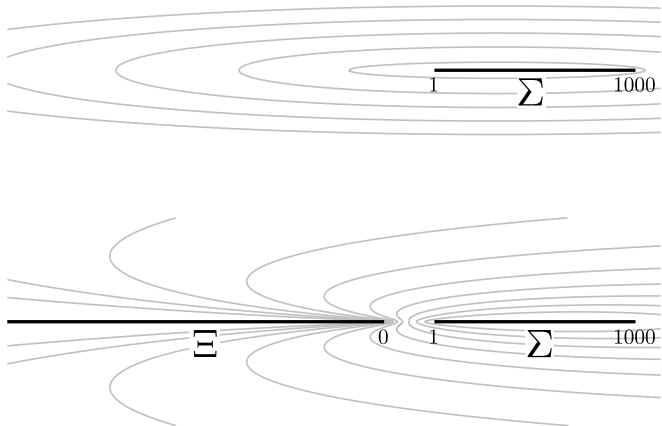
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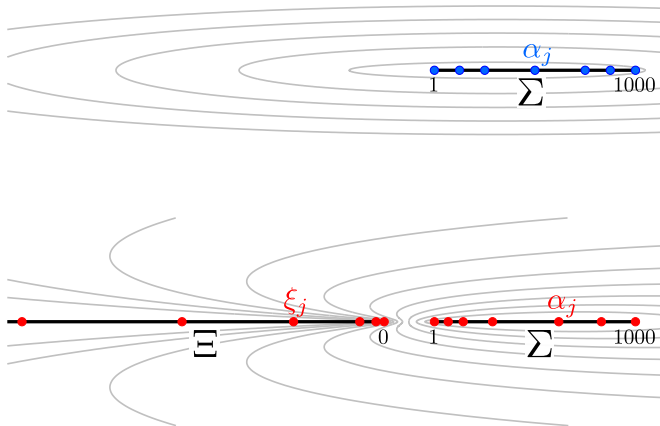
## Remarks

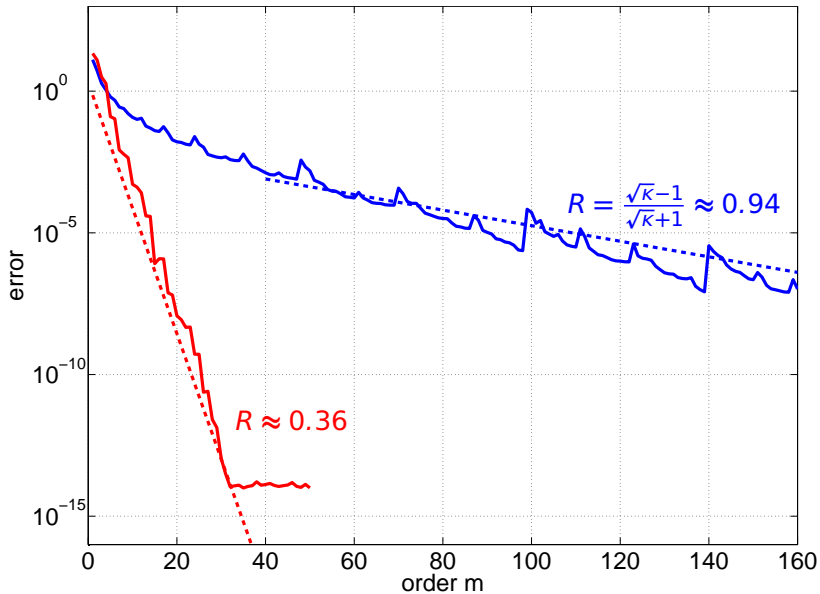
- ▶ 2 vectors storage need, 0 inner-products
  - ▶ If all  $\xi_j = \infty \Rightarrow$  polynomial interpolation at  $\{\alpha_1, \dots, \alpha_m\}$
  - ▶ Polynomial interpolation methods have been considered before [Huisinga et al 99] [Bergamaschi, Caliari & Vianello 04]
    - ▶ For  $\{\alpha_1, \dots, \alpha_m\}$  use Leja points, scaled to a set of unit capacity for stability [Reichel 90].
    - ▶ No such scaling is necessary with the PAIN method: simply choose  $\beta_j$  such that  $\|\mathbf{v}_{j+1}\| = 1, j = 1, \dots, m.$
  - ▶ For rational interpolation use **Leja-Bagby points**.
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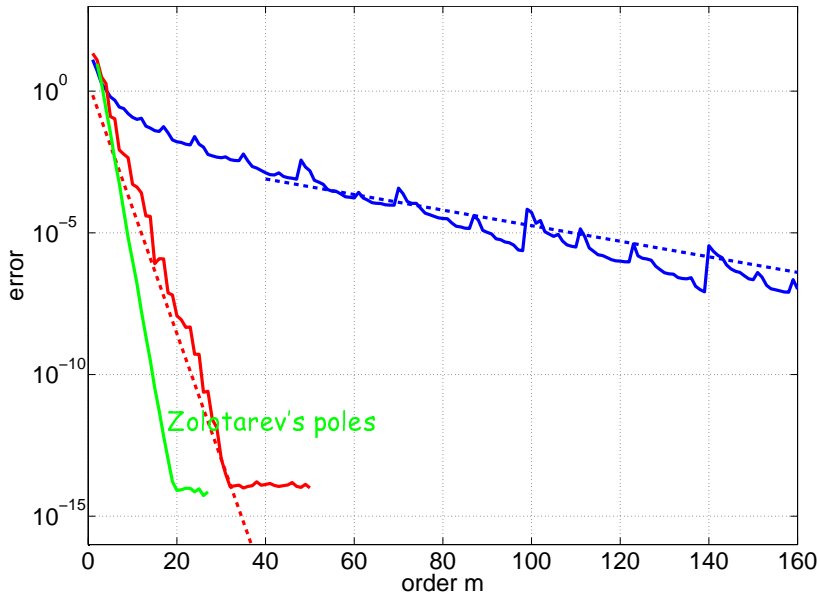
**Compute:**  $f(A)\mathbf{b} = \sqrt{A}\mathbf{b}$ ,  $A = \text{diag}(1, \dots, 1000)$ ,  $\mathbf{b} = [1, \dots, 1]^T$ .



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If good (or best) rational approximation  $r_m^*$  to  $f$  is known explicitly, one can directly evaluate  $r_m^*(A)\mathbf{b} \approx f(A)\mathbf{b}$ .

[Trefethen et al 06] [Frommer et al 06] [Schmelzer et al 07] [Hale et al 08]

However, using the poles  $\xi_j$  of  $r_m^*$  and suitable interpolation nodes  $\alpha_j$  as inputs for PAIN, we can achieve essentially the same accuracy at the same computational cost.

Moreover, the PAIN method is implicitly based on **exact** interpolation of  $f$  and hence robust to perturbations in  $r_m^*$ :

$$\limsup_{m \rightarrow \infty} \|f(A)\mathbf{b} - \mathbf{f}_m\|^{1/m} \leq R < 1$$

if  $\alpha_j, \xi_j$  are equilibrium-distributed on  $\Sigma, \Xi$ .

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## Rayleigh-Ritz extraction

There is a way to automatically choose near-optimal interpolation points  $\{\alpha_1, \dots, \alpha_m\}$  at iteration  $m$ :

1. Compute orthonormal basis  $V_m$  of  $\mathcal{Q}_m = q_{m-1}^{-1}\mathcal{K}_m$ .
2. "Determine" Rayleigh quotient  $A_m = V_m^*AV_m$ .
3. Compute  $\mathbf{f}_m = V_m f(A_m)V_m^* \mathbf{b}$ .

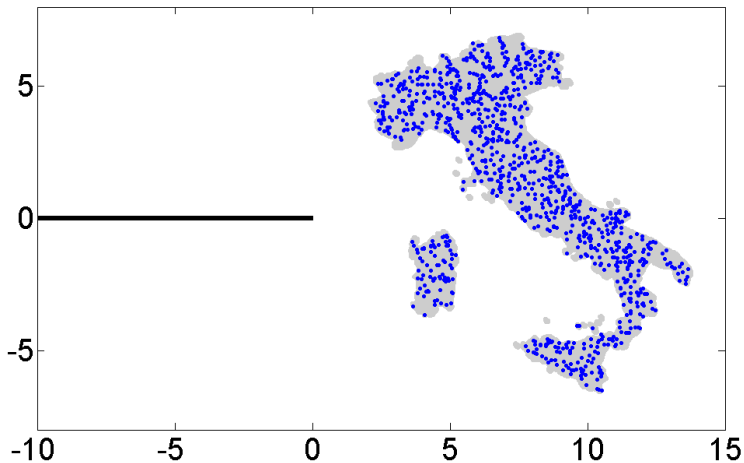
**Theorem:**  $\{\alpha_1, \dots, \alpha_m\} = \Lambda(A_m)$ .

**Theorem:**  $\|f(A)\mathbf{b} - \mathbf{f}_m\| \leq C \min_{p \in \mathcal{P}_{m-1}} \|f - p/q_{m-1}\|_{F(A)}$ .

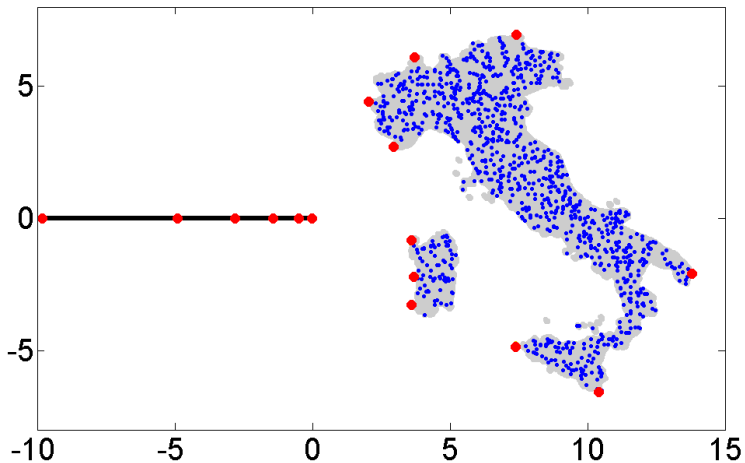
**Price:**  $m$  vectors storage need,  $m^2/2$  inner-products.

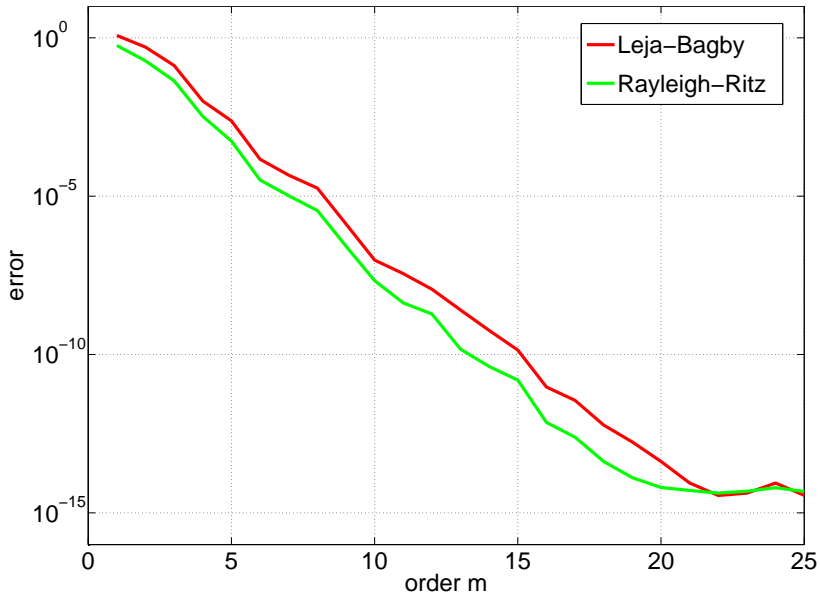
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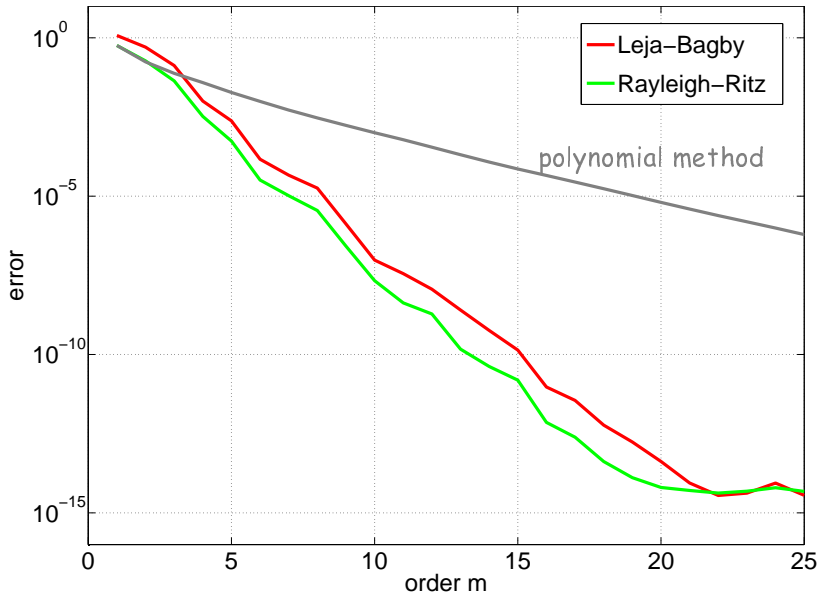
Compute:  $f(A)\mathbf{b} = \log(A)\mathbf{b}$ ,  $A$  normal with 1000 eigenvalues



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## Parameter-dependent problems

In practice, one often is not interested in  $f(A)\mathbf{b}$  but in  $f^\tau(A)\mathbf{b}$ , for  $\tau \in T$  from some **parameter set**  $T$ .

Given a **single** rational Krylov decomposition (as before)

$$AV_m K_m = V_{m+1} \underline{H}_m, \quad \mathcal{R}(V_m) = \mathcal{Q}_m = q_{m-1}^{-1} \mathcal{K}_m,$$

we compute **several** approximations

$$\mathbf{f}_m^\tau = V_m f^\tau(H_m K_m^{-1}) V_m^\dagger \mathbf{b} = r_m^\tau(A) \mathbf{b} = \frac{p_{m-1}^\tau}{q_{m-1}}(A) \mathbf{b},$$

where  $r_m^\tau$  Hermite-interpolates  $f^\tau$  at  $\Lambda(H_m K_m^{-1})$ .

---

## Example: Transfer function

$f^\tau(z) = (z - \tau)^{-1}$ , spectrum  $\Sigma = [0, +\infty)$ , parameters  $T = i[1, c]$ .

Let  $\omega_m(z) = (z - \alpha_1) \cdots (z - \alpha_m)$ , then

$$r_m^\tau(z) = \frac{1 - \frac{q_{m-1}(\tau) \omega_m(z)}{q_{m-1}(z) \omega_m(\tau)}}{z - \tau} = \frac{p_{m-1}^\tau}{q_{m-1}}$$

Hermite-interpolates  $f^\tau$  at  $\{\alpha_1, \dots, \alpha_m\}$ . Hence,  $\mathbf{f}_m^\tau = r_m^\tau(A)\mathbf{b}$ .

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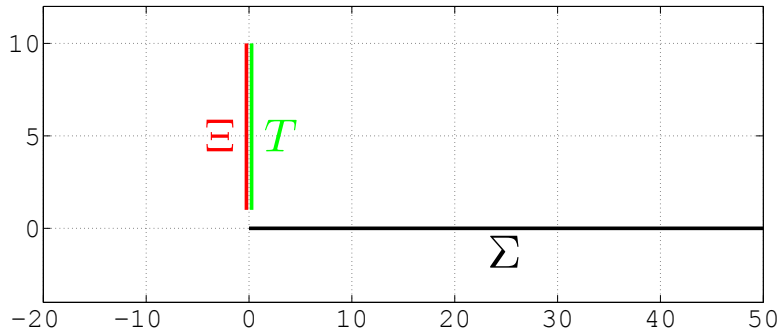
The relative error is

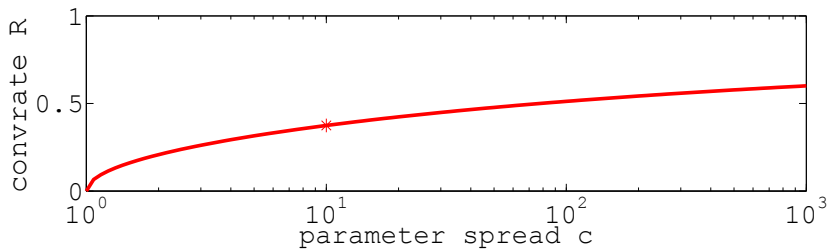
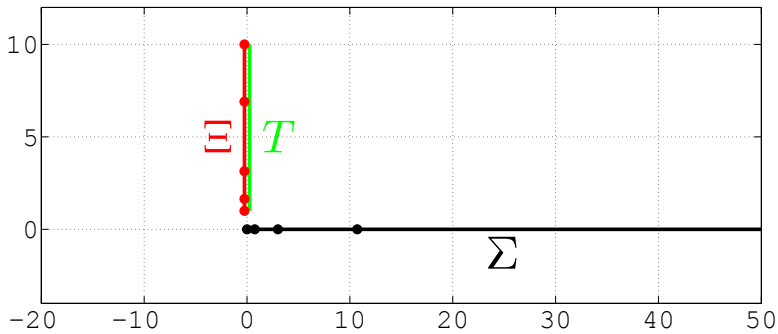
$$[f^\tau(z) - r_m^\tau(z)]/f^\tau(z) = \frac{q_{m-1}(\tau) \omega_m(z)}{q_{m-1}(z) \omega_m(\tau)}, \quad z \in \Sigma, \tau \in T,$$

and if  $\Xi = T$ , its minimization is related to the ADI problem.

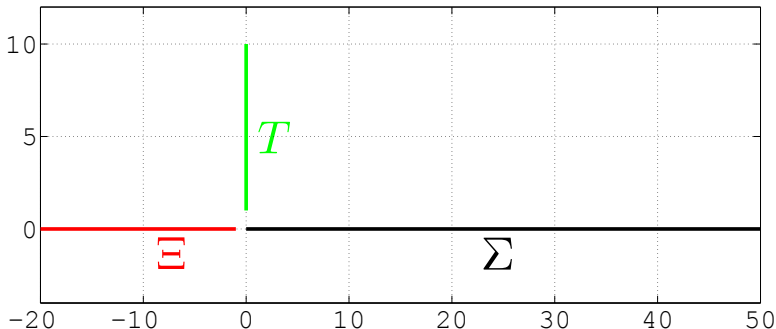
[Knizhnerman, Druskin & Zaslavsky 08]

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i.e.,  $\Xi = [-\infty, 0)$ . How to select  $\{\xi_1, \xi_2, \dots\} \subset \Xi$ ?



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Use standard tool to solve nonstandard approximation problem:

Assume to have a single repeated pole  $\xi$ . Then

$$r_m^\tau(z) = \frac{p_{m-1}^\tau(z)}{q_{m-1}(z)} = \frac{p_{m-1}^\tau(z)}{(z - \xi)^{m-1}} = \hat{p}_{m-1}^\tau(\hat{z}), \quad \hat{z} = (z - \xi)^{-1},$$

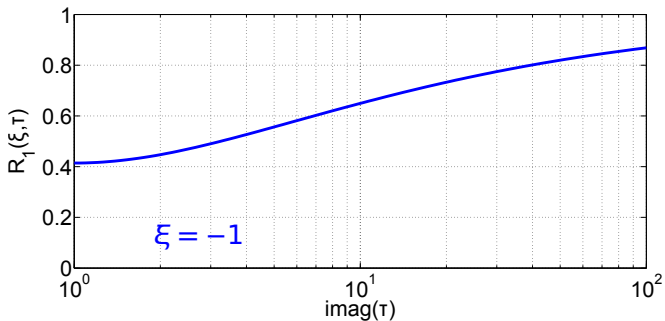
i.e., we have a polynomial problem: among  $p \in \mathcal{P}_{m-1}$

$$\text{minimize } \|f^\tau(\hat{z}^{-1} + \xi) - p(\hat{z})\|_{\hat{\Sigma}}, \quad \hat{\Sigma} = \{\hat{z} : z \in \Sigma\}.$$

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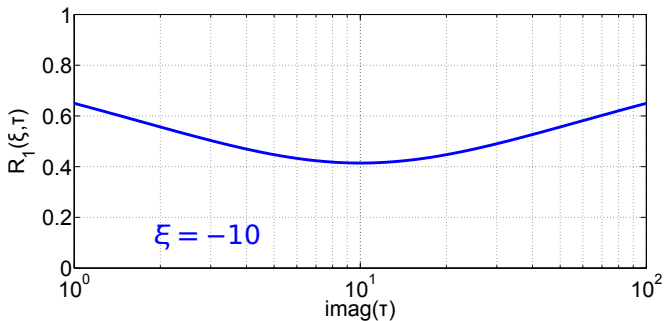
Apply Walsh's theory on polynomial approximation to obtain the asymptotic convergence rate  $R_1(\xi, \tau)$ . For the transfer function:

$$R_1(\xi, \tau) = \left( 1 + \frac{\sqrt{8}d^{3/4} + 4d^{1/2} + \sqrt{8}d^{1/4}}{1+d} \right)^{-1/2}, \quad d = -\tau^2/\xi^2.$$



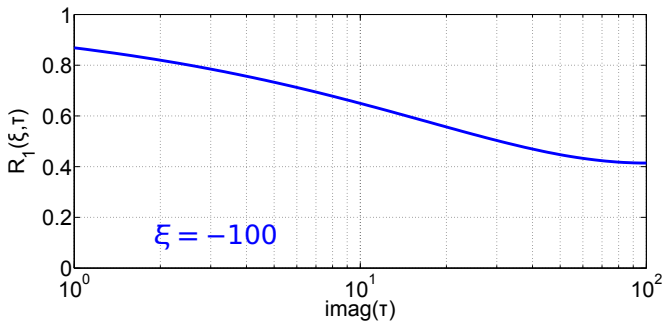
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Consider  $p$  poles  $\{\xi_1, \dots, \xi_p\}$  repeated cyclically.

The product form of the error

$$f^\tau(z) - r_m^\tau(z) = \frac{\frac{q_{m-1}(\tau) \omega_m(z)}{q_{m-1}(z) \omega_m(\tau)}}{z - \tau}$$

allows to conclude that

$$R(\{\xi_1, \dots, \xi_p\}, \tau) = \prod_{j=1}^p R_1(\xi_j, \tau)^{1/p}$$

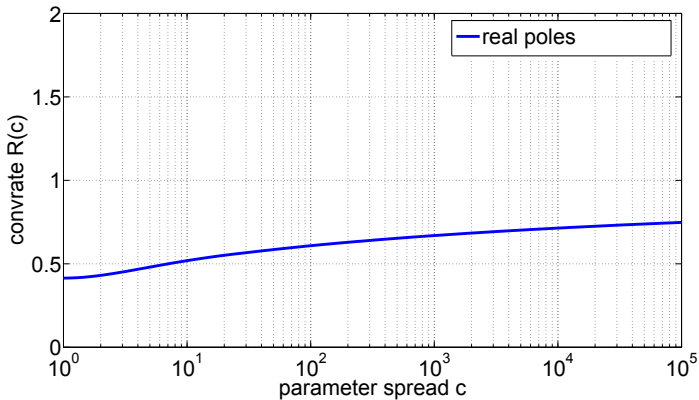
is the asymptotic convergence rate for this pole sequence.

$\Rightarrow$  Find  $\{\xi_1^*, \dots, \xi_m^*\}$  minimizing the worst-case rate

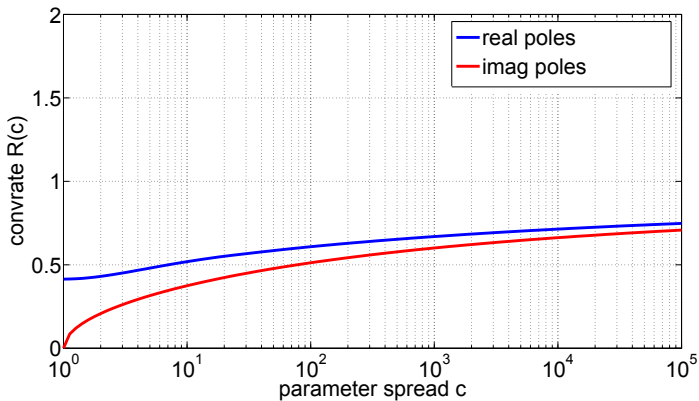
$$\max_{\tau \in T} R(\{\xi_1, \dots, \xi_m\}, \tau).$$

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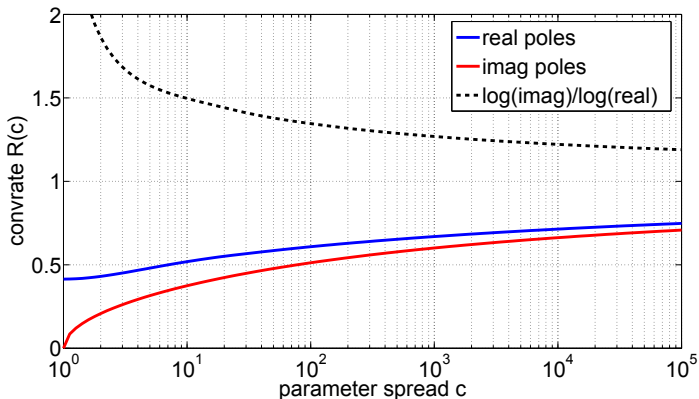
Find optim. poles by nonneg. minimization  $\|\mathbf{e} - M\mathbf{x}\|_\infty$ . Here is the optimal overall-convergence rate on  $T = i[1, c]$  depending on  $c$ .



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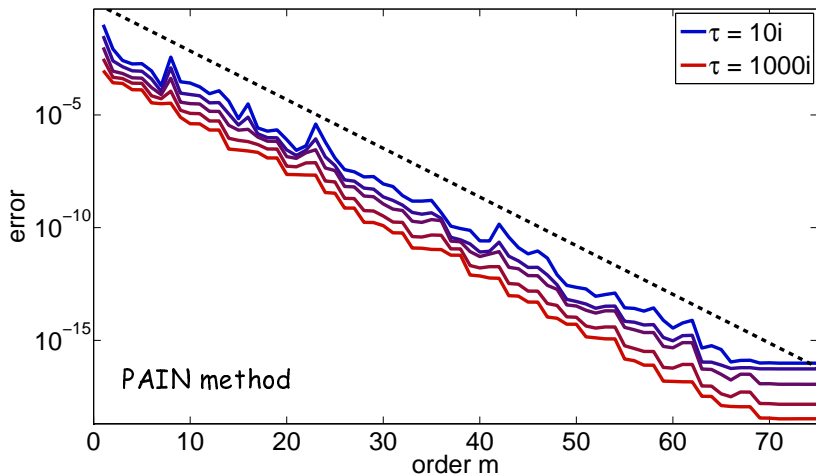
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Ex: If solve of complex system is 1.5 as expensive as real solve, use imaginary poles only if  $c < 10$ !

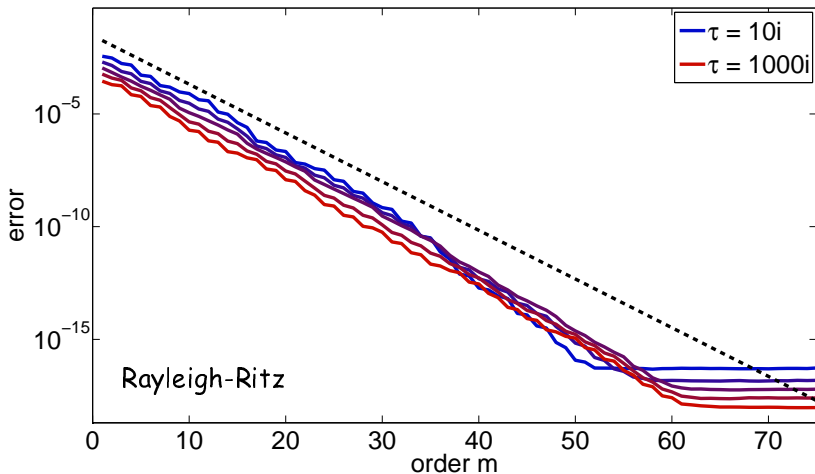
Compute:  $f^\tau(A)\mathbf{b} = (A - \tau I)^{-1}\mathbf{b}$ ,

$A = \text{diag}(0, \dots, 1e4)$ ,  $\tau \in T = i[10, 1000]$ ,  $\mathbf{b} = \text{randn}$ .



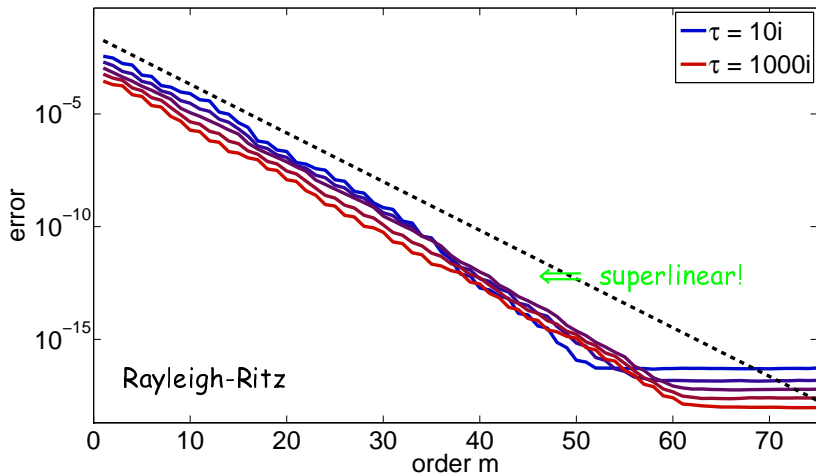
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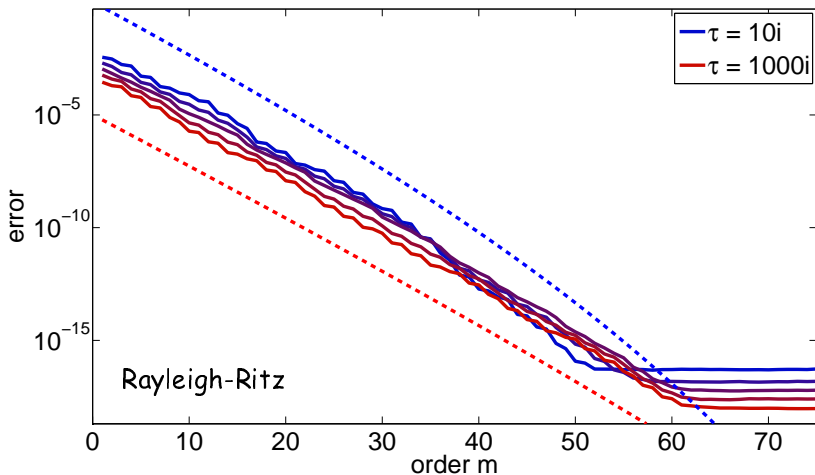
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## Summary

- ▶ Have characterized the general form of a rational Krylov decomposition.
  - ▶ All existing rational Krylov methods fit into this framework.
  - ▶ Propose "PAIN" as an efficient and robust rational Krylov method for problems with known spectral properties.
  - ▶ Have presented simple method for finding constrained pole sequences yielding asymptotically optimal convergence.
  - ▶ This method may be applied for general  $f$  by using Cauchy integral representation.
  - ▶ Can explain superlinear convergence observed with Rayleigh-Ritz extraction for Hermitian problems using weighted potential theory [Beckermann, G. & Vandebril 09].
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# See my poster for "rational Ritz values" and "inexact solves":



## Rational Krylov methods and approximation of $f(A)b$

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Equipe d'Analyse Numérique et d'Optimisation



### Matrix Functions

Given a square matrix  $A$  of size  $N \times N$ , a vector  $b$  of length  $N$  and a scalar function  $f(z)$ ,

$$f(A)b := p(A)b,$$

where  $p \in \mathcal{P}_{N-1}$  is a polynomial of degree  $\leq N-1$  that Hermite-interpolates  $f$  at the eigenvalues of  $A$ . In typical applications the matrix  $A$  is large and sparse.

#### Some Applications

- $f(z) = (z - i\omega)^{-1}$ : model reduction in the frequency domain,
- $f(z) = \exp(-tz)$ : time-integration of linear ODE's, exponential integrators, e.g., in geophysics or chemistry,
- $f(z) = \sqrt{z}$ : simulation of Brownian motion of molecules or sampling from Gaussian Markov random fields,
- $f(z) = \text{sign}(z)$ : simulations in quantum chromodynamics.

### Rational Krylov Spaces

**Definition:** Given a sequence of polynomials

$$q_{m-1}(z) = \prod_{\substack{j=1 \\ \xi_j \neq \infty}}^{m-1} (z - \xi_j), \quad m = 1, 2, \dots,$$

where  $\xi_j \in \mathbb{C} \setminus \Lambda(A)$ . Then the associated rational Krylov spaces of order  $m$  are defined as

$$\mathcal{K}_m(A, b) := q_{m-1}(A)^{-1} \mathcal{X}_m(A, b),$$

where  $\mathcal{X}_m(A, b) = \text{span}\{A^0 b, A^1 b, \dots, A^{m-1} b\}$ .

### Rational Ritz values

... are the eigenvalues of the Rayleigh quotient  $A_m = Q_m^* A Q_m$ , denoted by  $\Theta = \{\theta_1, \dots, \theta_m\}$ .

Let  $A$  be Hermitian. Then the  $\theta_k$ 's lie in the spectral interval of  $A$  and *interlace* the eigenvalues  $\Lambda(A) = \{\lambda_1, \dots, \lambda_N\}$ :

(\*) In any interval  $(\theta_k, \theta_{k+1})$  there is at least one eigenvalue  $\lambda_k$  of  $A$ .

Moreover, the rational Ritz values are zeros of orthogonal rational functions and may be characterized as (see, e.g., [2, 3])

(\*\*) The  $\theta_k$ 's are the zeros of the minimizer of  $\|p(A)q_{m-1}(A)^{-1}b\|$  among all monic  $p \in \mathcal{P}_m^\infty$ .

Logarithmic potential theory can explain the *asymptotic* distribution of the rational Ritz values. Therefore we consider

- a sequence of Hermitian matrices  $\{A_N\}$ , each of size  $N \times N$ , whose eigenvalue counting measures converge to a Borel probability measure  $\sigma$  in the weak-\* sense,
- a sequence of vectors  $\{b_N\}$ , each of length  $N$ ,
- a ray sequence of integers  $\{m_N\}$  such that  $m_N/N \rightarrow t \in (0, 1)$  as  $N \rightarrow +\infty$ ,
- a sequence of polynomials  $\{q_N\}$ , each of degree  $m_N - 1$ , whose zero counting measures converge to a Borel measure  $\nu$ ,  $\|\nu\| = t$ ,
- the sequence  $\{\Theta_N\}$  of rational Ritz values of order  $m_N$ .

**Tools from Potential Theory** Associated with a (signed) Borel measure  $\mu_1$  is the logarithmic potential

$$U^{\mu_1}(z) := \int \frac{1}{|z - w|} d\mu_1(x).$$

### Inexact solves & error estimators

In each iteration of the rational Arnoldi method a linear system of the form  $(A - \xi_j I)x_j = q_j$  is solved. If the residuals are collected in a matrix  $R_m$ , then (1) becomes

$$A Q_{m+1} K_m = Q_{m+1} H_m + R_m. \quad (2)$$

Setting  $E_m := -R_m K_m^\dagger Q_m^*$ , we observe that we have computed an exact Arnoldi decomposition

$$(A + E_m) Q_{m+1} K_m = Q_{m+1} H_m$$

for the matrix  $A + E_m$ . The Rayleigh quotient  $\tilde{A}_m$  computed from the data  $K_m$  and  $H_m$  satisfies

$$\begin{aligned} \tilde{A}_m &= Q_m^* (A + E_m) Q_m \\ &= Q_m^* A Q_m + Q_m^* (-R_m K_m^\dagger Q_m^*) Q_m \\ &= \tilde{A}_m - Q_m^* R_m K_m^\dagger Q_m^*. \end{aligned}$$

Here,  $\tilde{A}_m := Q_m^* A Q_m$  is referred to as the *corrected Rayleigh quotient*, because it is a compression of  $A$  instead of  $A + E_m$ . It can be computed from  $\tilde{A}_m$  without explicit projection, only by additional inner-products  $Q_m^* R_m$ .

We now decompose the error

$$\|f(A)b - f_m\| \leq \underbrace{\|f(A)b - f(A + E_m)b\|}_{\text{sensitivity error}} + \underbrace{\|f(A + E_m)b - f_m\|}_{\text{approximation error}},$$

and estimate

$$\text{sensitivity error} \approx \|f(\tilde{A}_m) Q_m^* b - f(\tilde{A}_m) Q_m^* b\|.$$

It is advisable to terminate the rational Arnoldi method if the ap-