

High order parametric polynomial approximation of conic sections

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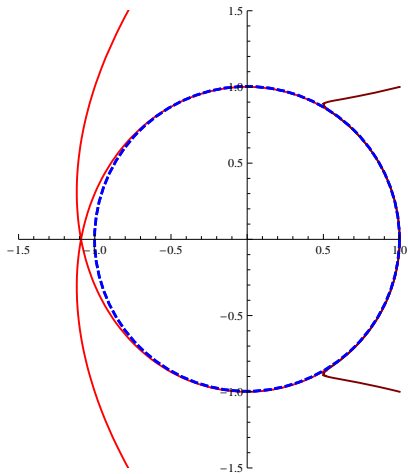


Figure: The unit circle (blue dashed), quintic polynomial approximation given by Lyche and Mørken in 1995 (brown) and the new quintic approximant (red).

Parametric approximation

- Let

$$f(x, y) = 0, \quad (x, y) \in \mathcal{D} \subset \mathbb{R}^2,$$

be a segment of a regular smooth planar curve \mathbf{f} .

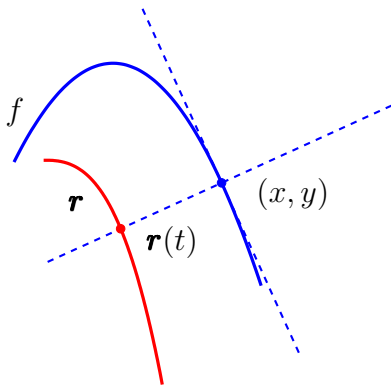
- Suppose

$$\mathbf{r} : I \subset \mathbb{R} \rightarrow \mathbb{R}^2 : t \mapsto \begin{pmatrix} x_r(t) \\ y_r(t) \end{pmatrix}$$

is a parametric approximation of the curve segment \mathbf{f} and

$$f(x_r(t), y_r(t)) = \varepsilon(t), \quad t \in [a, b].$$

- Consider the regular parameterization of \mathbf{r} with respect to the normal of \mathbf{f} , i.e., every point (x, y) on a curve f defines a unique parameter $t := t(x, y)$ on a curve \mathbf{r} :



- This provides an upper bound on Hausdorff and parametric distance.

Theorem

If the curve \mathbf{r} can be regularly reparameterized by the normal to f , and ε is small enough, the distance between curves is bounded by

$$\max_{(x,y) \in \mathcal{D}} \frac{|\varepsilon(t(x,y))|}{\sqrt{f_x^2(x,y) + f_y^2(x,y)}} + \mathcal{O}(\varepsilon(t(x,y))^2).$$

Conic sections

- Only ellipse and hyperbola are interesting to consider (**no polynomial parameterization**).
- A particular conic section is given as

$$f(x, y) := x^2 \pm y^2 - 1 = 0.$$

- The main problem: find two nonconstant polynomials x_n and y_n of degree at most n , such that

$$x_n^2(t) \pm y_n^2(t) = 1 + \varepsilon(t),$$

where ε is a polynomial of degree at most $2n$.

- Assume that at least one point is interpolated, i.e., $\varepsilon(0) = 0$. In order to get ε as small as possible, it is natural to choose $\varepsilon(t) := t^{2n}$.
- If also a tangent direction is prescribed at the interpolation point, we have

$$(x_n(0), y_n(0)) = (1, 0), \quad (x'_n(0), y'_n(0)) = (0, 1).$$

- Thus

$$x_n(t) = 1 + \sum_{\ell=1}^n a_\ell t^\ell,$$
$$y_n(t) = t + \sum_{\ell=2}^n b_\ell t^\ell,$$

and



$$x_n^2(t) \pm y_n^2(t) = 1 + (a_n^2 \pm b_n^2) t^{2n}.$$

- A reparameterization $t \mapsto \sqrt[2n]{|a_n^2 \pm b_n^2|} t$ gives

$$x_n(t) := 1 + \sum_{\ell=1}^n \alpha_\ell t^\ell, \quad y_n(t) := \sum_{\ell=1}^n \beta_\ell t^\ell, \quad \beta_1 > 0,$$

where

$$x_n^2(t) \pm y_n^2(t) = 1 + \text{sign}(a_n^2 \pm b_n^2) t^{2n}.$$

Many acceptable solutions exist.

Table: The number of appropriate solutions in all three cases for $n = 1, 2, \dots, 10$.

n	1	2	3	4	5	6	7	8	9	10
elliptic	1	1	3	6	15	27	63	120	246	495
hyp. $a_n^2 < b_n^2$	1	0	1	2	5	8	19	32	68	120
hyp. $a_n^2 > b_n^2$	0	1	0	2	0	9	0	32	0	125

Solutions

- Solving the equation

$$x_n^2(t) \pm y_n^2(t) = 1 \pm t^{2n}$$

is equivalent to solving

$$x_n^2(t) \pm y_n^2(t) = 1$$

in the factorial ring $\mathbb{R}[t]/t^{2n}$.

- There are additional restrictions, classic algebraic tools can not be applied.

Idea

- Rewrite

$$(x_n(t) + i y_n(t))(x_n(t) - i y_n(t)) = \prod_{k=0}^{2n-1} \left(t - e^{i \frac{2k+1}{2n} \pi} \right),$$

where the right-hand side is the factorization of $1 + t^{2n}$ over \mathbb{C} .

- Thus

$$x_n(t) + i y_n(t) = \gamma \prod_{k=0}^{n-1} \left(t - e^{i \sigma_k \frac{2k+1}{2n} \pi} \right), \quad \gamma \in \mathbb{C}, \quad |\gamma| = 1,$$

where $\sigma_k = \pm 1$.

Best solution

- The best solution should have the **minimum error term** ε .
- It turns out that this happens if

$$\beta_1 = \frac{1}{\sin \frac{\pi}{2n}}.$$

- Surprisingly, any solution for the elliptic case, for which x_n is even and y_n is odd, can be transformed to the hyperbolic solution by the map

$$\begin{aligned}x_n(t) &\mapsto x_n(i t), \\y_n(t) &\mapsto -i y_n(i t).\end{aligned}$$

- In particular, this is true for the best solution too, thus it is **enough to consider the elliptic case only**.

Theorem

Coefficients of the best solution for the elliptic case are

$$\alpha_k = \begin{cases} \sum_{j=0}^{k(n-k)} P(j, k, n-k) \cos\left(\frac{k^2\pi}{2n} + \frac{\pi j}{n}\right); & k \text{ is even,} \\ 0; & k \text{ is odd,} \end{cases}$$

$$\beta_k = \begin{cases} 0; & k \text{ is even,} \\ \sum_{j=0}^{k(n-k)} P(j, k, n-k) \sin\left(\frac{k^2\pi}{2n} + \frac{\pi j}{n}\right); & k \text{ is odd,} \end{cases}$$

where $P(j, k, r)$ denotes the number of integer partitions of $j \in \mathbb{N}$ with $\leq k$ parts, all between 1 and r , where $k, r \in \mathbb{N}$, and $P(0, k, r) := 1$.

Examples

Table: Polynomial approximation of the **unit circle** and maximal normal (radial) error.

n	x_n and y_n	error
3	$x_3(t)=1-2t^2$ $y_3(t)=2t-t^3$	2
4	$x_4(t)=1+(-2-\sqrt{2})t^2+t^4$ $y_4(t)=(\sqrt{2+\sqrt{2}}+\sqrt{2-\sqrt{2}})(t-t^3)$	0.414213
5	$x_5(t)=1+(-3-\sqrt{5})t^2+(1+\sqrt{5})t^4$ $y_5(t)=(1+\sqrt{5})t+(-3-\sqrt{5})t^3+t^5$	0.089987
6	$x_6(t)=1-2(2+\sqrt{3})t^2+2(2+\sqrt{3})t^4-t^6$ $y_6(t)=(\sqrt{2+\sqrt{6}}t-\sqrt{2(3+2\sqrt{3})}t^3+(\sqrt{2+\sqrt{6}})t^5$	0.013886
\vdots	\vdots	\vdots
15	...	$1.07280 \cdot 10^{-15}$

It can be shown that the error is $\mathcal{O}(n^{-2n})$.

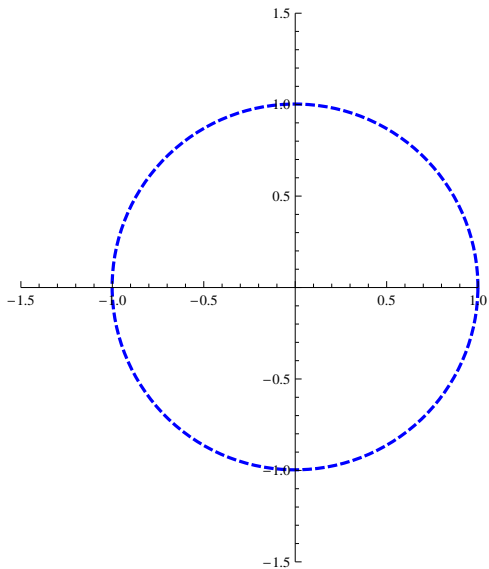


Figure: The unit circle.

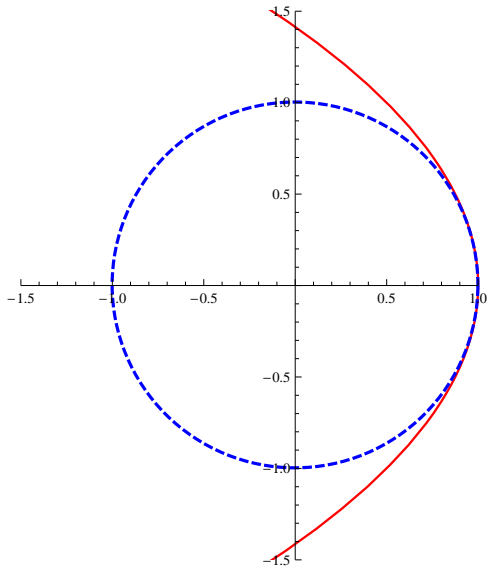


Figure: The unit circle and its polynomial approximant for $n = 2$.

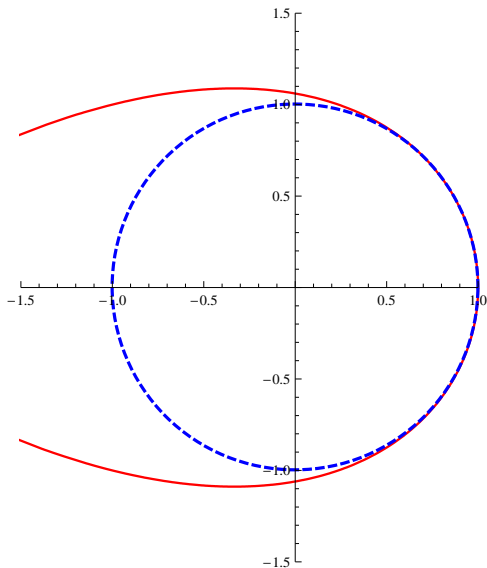


Figure: The unit circle and its polynomial approximant for $n = 3$.

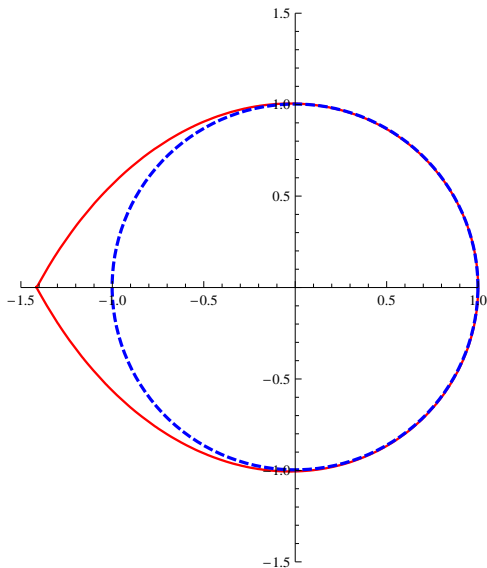


Figure: The unit circle and its polynomial approximant for $n = 4$.

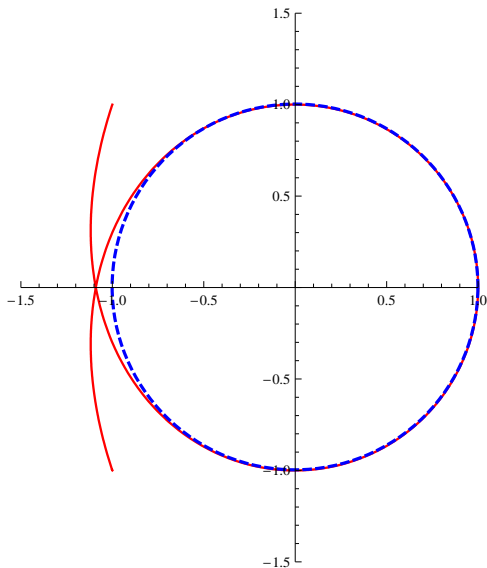


Figure: The unit circle and its polynomial approximant for $n = 5$.

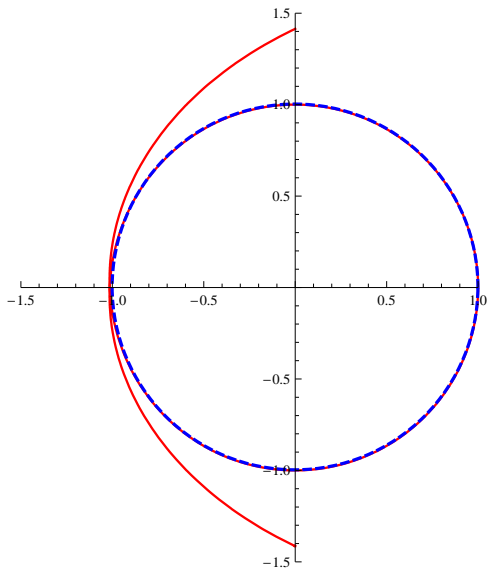


Figure: The unit circle and its polynomial approximant for $n = 6$.

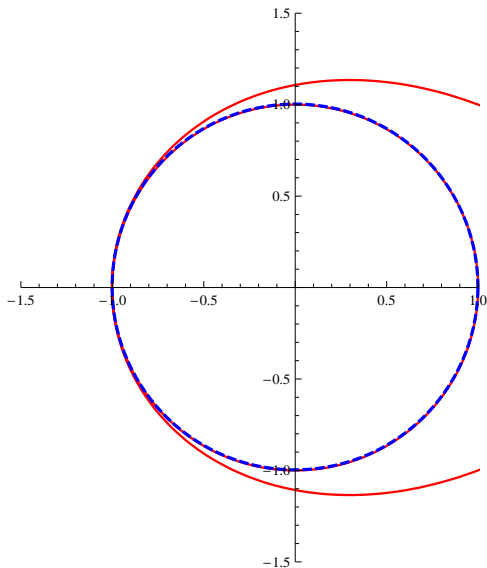
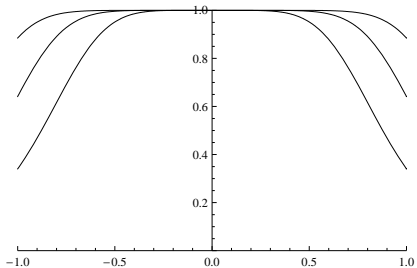
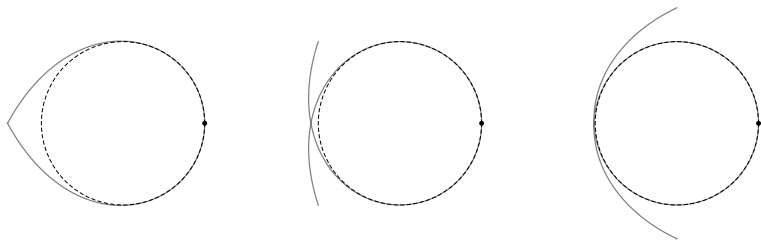


Figure: The unit circle and its polynomial approximant for $n = 7$.

Approximants and curvatures



Cycling

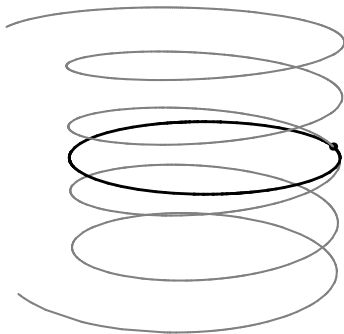


Figure: Unit circle together with the cycles of the approximant for $n = 20$ and $t \in [-1, 1]$.

Höllig-Koch conjecture

Conjecture

A polynomial planar parametric curve of degree n can interpolate $2n$ given points with an approximation order $2n$.

Theorem

Höllig-Koch conjecture holds true for conic sections.

Idea of a proof:

- asymptotic approach,
- a particular nonlinear system has to be studied,
- an existence of a solution guarantees the optimal approximation order,
- solution is obtained by canonical form and optimal solutions for ellipse and hyperbola.

Ellipse

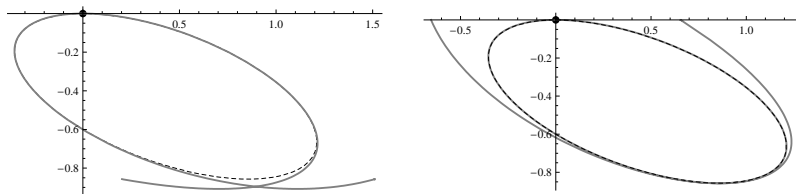


Figure: Approximation of the ellipse $\frac{1}{2}x^2 + xy + \frac{5}{3}y^2 + y = 0$ with the best approximant of degree $n = 5, 7$.

Hyperbola

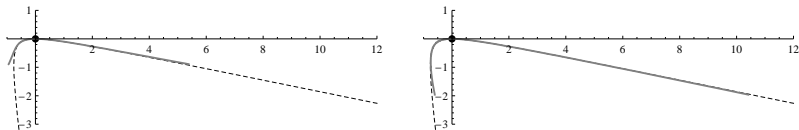
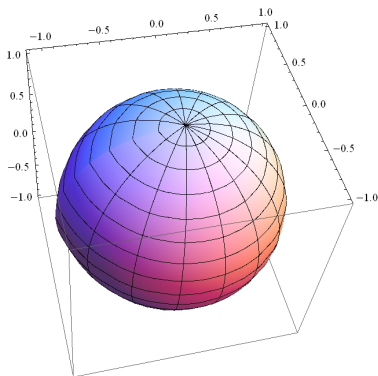


Figure: Approximation of the hyperbola $\frac{1}{5}x^2 + xy + \frac{1}{8}y^2 + y = 0$ with the best approximant of degree $n = 3, 4$.

Sphere approximation



Particular polynomials of degree 5 in u and v yield:

$$x(u, v) = (1 + (-3 - \sqrt{5})u^2 + (1 + \sqrt{5})u^4)(1 + (-3 - \sqrt{5})v^2 + (1 + \sqrt{5})v^4)$$

$$y(u, v) = ((1 + \sqrt{5})u + (-3 - \sqrt{5})u^3 + u^5)(1 + (-3 - \sqrt{5})v^2 + (1 + \sqrt{5})v^4)$$

$$z(u, v) = (1 + \sqrt{5})v + (-3 - \sqrt{5})v^3 + v^5$$