Optimal Polynomial Admissible Meshes on the Closure of $\mathscr{C}^{1,1}$ Bounded Domains

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1 Introducing Polynomial Admissible Meshes

- Defining (Weakly) Admissible Meshes, Optimal Admissible Meshes.
- Main properties and motivations.
- Building Admissible Meshes, the state of the art.
- **2** A new result for $\mathcal{C}^{1,1}$ bounded domains
 - Main Result.
 - Tools: Bernstein Inequality and regularity property of oriented distance function
 - Sketch of the proof.



In 2008 J.P. Calvi and N.Levenberg proposed these well promising definitions .

Admissible Meshes, AM

Let $K \subset \mathbb{R}^d$ (or \mathbb{C}^d) be a compact polynomial determining set. The sequence $\{A_n\}_{\mathbb{N}}$ of finite subsets of *K* is said to be an **Admissible Mesh** for *K* if there exist *C*, *s* > 0 such that

Card
$$A_n = O(n^s)$$

 $\|p\|_K \leq C \|p\|_{A_n} \, \forall p \in \mathscr{P}^n(K).$

Weakly Admissible Meshes, WAM

If instead $C = C_n = O(n^q)$, then we say that A_n is a **Weakly** Admissible Mesh.



By the definitions both AMs and WAMs are determining for $\mathscr{P}^n(K)$, thus we have

Card
$$A_n \ge \dim \mathscr{P}^n(K) = \binom{n+d}{n} = O(n^d)$$

For this reason A. Kroó introduced

Optimal Admissible Mesh

The AM A_n w.r.t. $K \subset \mathbb{R}^d$ is said to be **optimal** if

Card $A_n = O(n^d)$.

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DLS Approximation on a WAM - Calvi Levenberg (2008)

Let $K \subset \mathbb{R}^d$ be compact and polynomial determining, A_n a WAM on it and $f \in \mathscr{C}^0(K)$, then one has

$$\|f - \Lambda_{A_n} f\|_{\mathcal{K}} \leq \left(1 + C_n \left(\|f\|_{\mathcal{K}} (1 + \sqrt{\operatorname{Card}(A_n)})\right)\right) d_n(f, \mathcal{K})$$

Where Λ_{A_n} is the discrete least squares (DLS) operator performed sampling *f* on A_n and $d_n(f, K)$ is the error of best polynomial approximation to *f* on *K*.

Mild regularity of f and $K \Longrightarrow$ convergence of DLS operator.



WAMs work nicely under some fundamental operations.

- Stability under affine mapping, union and tensor product.
- "Weak" stability under polynomial mapping.
- Supersets of WAM are WAMs.
- Good interpolation sets are WAMs.



Discrete Extremal Sets - Bos, De Marchi, Sommariva and Vianello

Starting from a WAM one can extract by standard Numerical Linear Algebra

- AFP Approximate Fekete Points
- ALS Approximate Leja Sequences

such that

Unisolvent sets.

- Slowly increasing Lebesgue constants.
- Same asymptotic (in measure theoretic sense).



Two questions naturally arise..

Question 1

How to build AMs or even WAMs for a given K?

Question 2

How to build Optimal AMs for a given K?

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One should choose/combine different results requiring K to have

particular shape and/or smoothness

- Elich-Zeller: Double degree Chebyshev points for the interval are AM of constant 2.
- **shape:** Use symmetries of *K*, polar coordinates, tensors, quadratic maps..
- **Calvi-Levenberg** used Multivariate Markov Inequality.
- **Kroó:** Star shaped bounded domains with smooth Minkowski functional + Bernstein Inequality .
- Piazzon-Vianello Mapping and Perturbing WAMs and AMs.
- W. Plesniak improve for sub-analytic sets.
- **Kroó** result on Analytic Graph domains.
- Bloom Bos Calvi and Levenberg existence for L-regular sets.

Building AM by Markov Inequality



Markov Inequality (MI)

The set K is preserving a Markov Inequality of constant M_K and exponent r if

$$\||\nabla p|\|_{\mathcal{K}} \le M_{\mathcal{K}} n^{r} \|p\|_{\mathcal{K}}. \quad \forall p \in \mathscr{P}^{n}(\mathbb{R}^{d})$$

MI holds under mild assumptions on K, typically r = 2.

Idea: take any equally spaced grid having step size $O(n^{-r})$.

Calvi Levenberg (2008)

If $K \subset \mathbb{R}^d$ preserves a Markov Inequality of exponent *r*, then it has a AM with $O(n^{rd})$ points.

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Using the classical Bernstein Inequality on segments and star-shaped property it has been proved

Kroó (2011)

Let $K \subset \mathbb{R}^d$ be compact and star-shaped with $\mathscr{C}^{1+\alpha}$ smooth Minkowski functional. Then *K* has an AM with $O(n^{\frac{2d+\alpha-1}{\alpha+1}})$ points.



Perturbation Result Piazzon Vianello

Let $K \subset \mathbb{C}^d$ be a polynomially convex and Markov compact set. If that there exists a sequence of compact sets $\{K_j\}_{\mathbb{N}}$ such that

• there exists $A_{n,j}$ (W)AM for K_j having constant $C_{n,j}$ and

■
$$d_{\mathcal{H}}(K, K_j) \le \epsilon_j$$
 where $\limsup_{i} \epsilon_j C_{n,j} = 0 \forall n$.

Then K has a (W)AM.

Mapping Result Piazzon Vianello

(W)AMs are "weakly" stable under smooth mapping: for any holomorphic map $\varphi : Q \to K$ there exists $j_{\varphi}(n) = O(\log n)$ such that $B_n := \varphi(A_{n:j_{\varphi}(n)})$ is a WAM for *K*.



AM having Card $A_n = O((n \log n)^d)$ as the ones above are termed **nearly optimal.**

- W. Plesniak showed that Piazzon-Vianello results in particular apply to any *compact sub-analytic set* that hence has a nearly optimal AM.
- A. Kroó proved that analytic graph domains have a nearly optimal AM
- Bloom,Bos,Calvi and Levenberg showed (non constructively) that any *L-regular* compact set has.



Question 2

How to build Optimal AMs for a given K?

- Polytopes, Balls have Optimal AMs by 1dim techniques, symmetry or thanks to the particular shape and finite unions.
- The Kroó result applies to **star-shaped** \mathscr{C}^2 smooth sets.
- The Kroó result has been refined: if d = 2, then C² smoothness can be replaced by uniform interior ball condition.

What about sets with a more general shape?



Idea: Smoothness may completely replace Particular Shape.

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Piazzon 2013

Let Ω be a bounded $\mathscr{C}^{1,1}$ domain in \mathbb{R}^d , then there exists an optimal admissible mesh for $K := \overline{\Omega}$.

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Tools I



Bernstein Inequality

Let $p \in \mathscr{P}^n(\mathbb{R})$, then for any $a < b \in \mathbb{R}$ we have

$$\left|\frac{dp}{dx}(x)\right| \le \frac{n}{\sqrt{(x-a)(b-x)}} ||p||_{[a,b]}.$$
 (BI)

and thus
$$\left|\frac{dp}{dx}\left(\frac{a+b}{2}\right)\right| \leq \frac{2\mathbf{n}}{(b-a)} ||p||_{[a,b]}.$$

Markov Tangential Inequality

Let $p \in \mathscr{P}^n(\mathbb{R}^d)$, then for any $x_0 \in \mathbb{R}^d$, r > 0 and $v \in \mathbb{S}^{d-1} \cap \mathcal{T}_x \partial B(x_0, r)$ we have

$$\left|\frac{dp}{dv}(x)\right| \le \frac{n}{r} ||p||_{B(x_0,r]}.$$
 (MTI)

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$\mathscr{C}^{1,1}$ domains

 $\Omega \subset \mathbb{R}^d$ domain whose boundary is locally the graph of a $\mathscr{C}^{1,1}$ function of controlled norm. If Ω is bounded, then all the parameters involved in this definition can be chosen uniformly.



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Geometric Characterization

Bounded $\mathscr{C}^{1,1}$ domains are characterized by the **uniform double sided ball condition**.



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Tools IV



Oriented Distance Function

$$b_{\Omega}(x) := \inf_{y \in \overline{\Omega}} |x - y| - \inf_{z \in \mathbb{C}\Omega} |x - z|$$





Regularity Properties of $b_{\Omega}(\cdot)$

If Ω is a bounded $\mathscr{C}^{1,1}$ domain, then there exists $\overline{\delta}$ such that for any $0 < \delta < \overline{\delta}$ we have

The metric projection on the boundary $x \mapsto \pi_{\partial\Omega}(x)$ is single valued on U_{δ} where U_{δ} is a δ -tubular neighborhood of $\partial\Omega$.

$$\bullet b_{\Omega} \in \mathscr{C}^{1,1}(U_{\delta}).$$

•
$$\nabla b_{\Omega}(x) = rac{x - \pi_{\partial\Omega}(x)}{b_{\Omega}(x)} \neq 0$$
 in U_{δ} .

- $\nabla b_{\Omega}(x)$ defines the outer normal unit vector field w.r.t. Ω .
- Differentiability across the boundary.
- We can take $\overline{\delta}$ as the radius of the ball.
- Level sets of b_{Ω} are $\mathscr{C}^{1,1}$ manifolds.



Bound normal derivatives of polynomials by a **modified Bernstein Inequality** along segments of metric projection. Thanks to boundary regularity.

↓

Find a **norming set** for $\overline{\Omega}$: by union of $m_n = O(\underline{n})$ hypersurfaces which are **level sets** of b_{Ω} and $K_{\delta} := \{x \in \overline{\Omega} : |b_{\Omega}(x)| \ge \delta\}$

$$\|p\|_{K} \leq 2 \max\left\{\|p\|_{K_{\delta}}, \|p\|_{\bigcup_{i=0}^{m_{n}}\Gamma^{i}}\right\}. \quad \forall p \in \mathscr{P}^{n}(\mathbb{R}^{d})$$

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construction





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Bound any directional derivative of polynomials by a **modified Bernstein Inequality** holding in K_{δ} .

Find a **weak norming set** for K_{δ} : w.r.t. $\overline{\Omega}$ by a grid mesh Z_n , of stepsize $O(n^{-1})$ i.e.

$$\|p\|_{K_{\delta}} \leq \|p\|_{Z_{n}} + rac{1}{\lambda} \|p\|_{\overline{\Omega}} \ \lambda > 2. \ \forall p \in \mathscr{P}^{n}(\mathbb{R}^{d})$$

Card
$$Z_n = O(n^d)$$

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Bound tangential derivatives of polynomials on $\Gamma^i s$ by the combination of

- **Regularity of** $\Gamma^i s \Rightarrow$ Ball property.
- Markov Tangential Inequality

↓

Find **weak norming sets** for $\bigcup_{i=0}^{m} \Gamma^{i}$ w.r.t. $\overline{\Omega}$ by the union of geodesic meshes $Y_{n}^{i} \subset \Gamma^{i}$ having *geodesic fill distance* $O(n^{-1})$

$$\|p\|_{\cup_{i=0}^{m_n}\Gamma^i} \leq \|p\|_{Y_n^i} + \frac{1}{\lambda} \|p\|_{\overline{\Omega}} \ \lambda > 2. \ \forall p \in \mathscr{P}^n(\mathbb{R}^d)$$

Card
$$\cup_{i=0}^{m_n} \Gamma^i = m_n O(\mathbf{n}^{d-1}) = O(\mathbf{n}^d).$$

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Finally we set

$$A_n := Z_n \cup \left(\cup_{i=0}^{m_n} \Gamma^i \right)$$

and the inequalities above read as

$$\begin{split} \|p\|_{\overline{\Omega}} &\leq 2\|p\|_{A_n} + \frac{2}{\lambda}\|p\|_{\overline{\Omega}} \quad \forall p \in \mathscr{P}^n(\mathbb{R}^d) \\ & \downarrow \\ \|p\|_{\overline{\Omega}} &\leq \frac{2\lambda}{\lambda - 2}\|p\|_{A_n} \quad \forall p \in \mathscr{P}^n(\mathbb{R}^d) \\ Card(A_n) &= O(n^d). \end{split}$$

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We conclude by recalling the open problem

Conjecture Bloom Bos Calvi and Levenberg

If $K \subset \mathbb{C}^d$ is compact and L-regular then there exists c := c(K) such that any array of degree c(K)n Fekete points forms an Admissible mesh for K, that hence is Optimal.



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Thank You!

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more details....



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Step 1 we build a norming set.

We pick $\delta < r$ and work out a modification of **(BI)** of the form

 $|D_{S}p(x)| \leq n\varphi_{\delta}(b_{\Omega}(x))||p||_{\Omega}$

Where *S* is a segment of metric projection, φ_{δ} arises as follows..

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Proof (2)







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Proof (3)



Then we define a function by integration along segments of metric projection

$$x\mapsto F_{n,\delta}(x):=n\int_0^{-b_\Omega(x)}\varphi_\delta(\xi)d\xi$$



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and consider equally spaced level sets

$$\Gamma^{i}_{n,\delta} := F_{n,\delta}(a^{i}) \ i = 0, 1, \dots m_{n} = O(n),$$

where

$$a^i = 0, \dots, \max_{\Omega} F_{n,\delta}$$

 $1/2 \geq a^{i+1} - a^i.$

Therefore we have

$$\begin{split} \|p\|_{\overline{\Omega}} &\leq & \|p\|_{\cup_i \Gamma_{n,\delta}^i} + 1/2 \|p\|_{\overline{\Omega}} \Rightarrow \\ &\leq & 2 \|p\|_{\cup_i \Gamma_{n,\delta}^i}. \end{split}$$

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For some technical reasons we switch

$$\cup_{i=0}^{m_n} \Gamma_{n,\delta}^i = \bigcup_{i=0}^{\tilde{m}_n} \Gamma_{n,\delta}^i \uplus \bigcup_{i=\tilde{m}_n+1}^{m_n} \Gamma_{n,\delta}^i$$

where the second set is a subset of

$$K_{\delta} := \{x \in \overline{\Omega} : |b_{\Omega}(x)| \ge \delta\}$$

and then we can replace it by K_{δ} itself.

$$||p||_{\overline{\Omega}} \leq 2 \max\{||p||_{\cup_{i}^{\tilde{m}_{n}} \Gamma_{n,\delta}^{i}}, ||p||_{\mathcal{K}_{\delta}}\}.$$

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Step 2: finding a norming mesh Z_n for K_{δ} .

we use (BI) jointly with $B(x, \delta) \subset \overline{\Omega} \ \forall x \in K_{\delta}$ to get

$$|\nabla p(x)| \leq \frac{n}{\delta} ||p||_{\overline{\Omega}}.$$

Thus we can build a suitable Z_n by a grid of step size $\frac{\delta}{4n} = O(n^{-1})$ and hence cardinality

Card
$$Z_n = O(n^d)$$

obtaining

$$||p||_{\mathcal{K}_{\delta}} \leq ||p||_{Z_n} + \frac{1}{4}||p||_{\overline{\Omega}}$$

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Step 3: finding a norming mesh Y_n for $\bigcup_{i=0}^{\tilde{m}_n} \Gamma_{n,\delta}^i$

• Any $\Gamma_{n,\delta}^i$ is a $\mathscr{C}^{1,1}$ manifold.

• we can pick a tangent ball of radius $\delta/2$ lying in Ω .



thus..



We can bound any tangential derivative by **MTI** applied to the ball to get

$$\left|\frac{dp}{dv}(x)\right| \leq \frac{n}{\delta/2} ||p||_{\overline{\Omega}} \, \forall v \in \mathbb{S}^{d-1} \cap \mathcal{T}_x \Gamma^i_{n,\delta}$$

Therefore if we pick a mesh Y_n^i on each $\Gamma_{n,\delta}^i$ having controlled **geodesic fill distance** $h^i = O(n^{-1})$ we get

$$\|p\|_{\cup_{i}^{\widetilde{m}_{n}}\Gamma_{n,\delta}^{i}}\leq \|p\|_{\cup_{i}Y_{n}^{i}}+\frac{1}{4}\|p\|_{\overline{\Omega}}.$$

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The regularity of $\Gamma_{n,\delta}^i$'s ensures that we can produce suitable $Y_{n,\delta}^i$ using $O(n^{d-1})$ points, since we have $m_n = O(n)$ level set $\Gamma_{n,\delta}^i$ we have

Card
$$Y_{n,\delta} := \text{Card} \cup_i Y_{n,\delta}^i = O(n^d)$$

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Step 4: joining all the inequalities.

Finally putting all together we get $||p||_{\overline{\Omega}} \le 2(||p||_{Y_{n,\delta}\cup Z_n} + \frac{1}{4}||p||_{\overline{\Omega}})$ and thus

$$\begin{aligned} \|p\|_{\overline{\Omega}} &\leq 4\left(\|p\|_{Y_{n,\delta}\cup Z_n}\right) \text{ where } \end{aligned} \tag{2} \\ \operatorname{Card}(Y_{n,\delta}\cup Z_n) &= O(n^d) \end{aligned}$$

That is optimal.

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