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Tikhonov regularization and matrix function evaluation

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Introduction and Motivations

We consider ill-conditioned linear systems

Ax = b

We mainly focus the attention on **full-rank** problems in which the singular values of *A* decay gradually to zero.

We want to construct an iterative solver able to overcome some of the typical drawback of the classical iterative solvers:

- **Semi-convergence**: the iterates initially approach the solution but quite rapidly diverge
- Strong dependence on the parameter-choice strategy: in order to prevent divergence a reliable stopping criterium has to be used
- **Poor accuracy**: typically holds for Krylov type methods based on the use of $A^T A$ (CGLS)

Reformulation of the problem

The basic idea is to solve the system Ax = b in two steps, first solving in some way the regularized system

$$(A + \lambda I) x_{\lambda} = b$$

and then recovering the solution x from the system

$$\left(A + \lambda I\right)^{-1} Ax = x_{\lambda}$$

that is equivalent to compute

 $x = f(A)x_{\lambda}$

where

$$f(z) = 1 + \lambda z^{-1}$$

By the definition of *f* the attainable accuracy depends on the conditioning of $(A + \lambda I)^{-1} A$. Theoretically the best situation is attained defining λ such that

$$\kappa(A + \lambda I) = \kappa((A + \lambda I)^{-1} A$$

In the SPD case taking

$$\lambda = \sqrt{\lambda_1 \lambda_N} \approx 1/\sqrt{\kappa(A)}$$

where λ_1 and λ_N are respectively the smallest and the largest eigenvalue of *A*, we obtain

$$\kappa(A + \lambda I) = \kappa((A + \lambda I)^{-1} A) = \sqrt{\kappa(A)}$$

The computation of the matrix function

The Arnoldi Algorithm

For the computation of $f(A)x_{\lambda}$ we use the standard Arnoldi method projecting the matrix A onto the Krylov subspaces generated by A and x_{λ}

Extension to Tikhonov regularization

The method can be extended to problems in which the exact right hand side *b* is affected by noise. We assume to work with a perturbed right-hand side $\tilde{b} = b + e_b$. Given λ and *H* we solve the regularized system

$$(A^T A + \lambda H^T H) x_{\lambda} = A^T \hat{t}$$

and then we approximate *x* by computing

$$\overline{x} = \left(A^T A\right)^{-1} \left(A^T A + \lambda H^T H\right) x_{\lambda} = f(Q) x_{\lambda}$$

where $Q = (H^T H)^{-1} (A^T A)$. As before, for the computation of $f(Q)x_{\lambda}$ we use the standard Arnoldi method projecting the matrix Q onto the Krylov subspaces generated by Q and x_{λ} . Now, at each step we have to compute the vectors $w_j = Qv_j, j \ge 1$, with $v_1 = x_{\lambda}/||x_{\lambda}||$, that is, to solve the systems $(H^T H) w_j = (A^T A) v_j$.

Theoretical error analysis

The field of values of A is defined as

$$F(A) := \left\{ \frac{y^H A y}{y^H y}, y \in \mathbb{C}^N \setminus \{0\} \right.$$

THEOREM 1 Assume that $F(A) \subset \mathbb{C}^+$. Then for the error $E_m := f(A)x_{\lambda} - ||x_{\lambda}|| V_m f(H_m)e_1$ we have

$$E_m \| \le K \frac{\lambda}{a^{m+1}} \prod_{i=1}^m h_{i+1,i} \| x_\lambda \|,$$

where a > 0 is the leftmost point of F(A) and $2 \le K \le 11.08$ (Crouzeix 2007; in the symmetric case K = 1).

The rate of convergence is almost independent of the choice of λ , and is closely related with the rate of the decay of $\prod_{i=1}^{m} h_{i+1,i}$.

THEOREM 2 Let σ_j , $j \ge 1$, be the singular values of an operator A.

$$\sum_{j\geq 1}\sigma_j^p<\infty \text{ for a certain }p>0$$

then

If

$$\prod_{i=1}^{m} h_{i+1,i} \le \left(\frac{\eta e p}{m}\right)^{m/p} \quad where \quad \eta \le \frac{1+p}{p} \sum_{j \ge 1} \sigma_j^p$$

- For discrete ill-posed problems the rate of decay of ∏^m_{i=1} h_{i+1,i} is superlinear and depends on the *p*-summability of the singular values of *A*, i.e., on the degree of ill-posedness of the problem.
- Each Arnoldi-based method (CG, FOM, GMRES) shows the same rate of convergence

Error analysis in computer arithmetics

<u>ASP</u>. We assume that x_{λ} , solution of $(A + \lambda I) x_{\lambda} = b$, is approximated by \overline{x}_{λ} with an accuracy depending on the choice of λ and the method used. In this way, the Arnoldi algorithm actually constructs the Krylov subspaces $K_m(A, \overline{x}_{\lambda})$. We obtain

Numerical experiments f_{0} f_{0}

Figure 1 - Stability behaviour of the error for the ASP method (1a) and the ATP method (1b) for different values of λ for the problem BAART 240.



Figure 2 - Maximum attainable accuracy with respect to the choice of lambda and with right-hand side affected by Gaussian noise $\sigma=10^{-3}$ (2b). The dimension of each problem is N=160.

An example of image restoration







Image restoration of

a 100×100 subimage of coins.png from MATLAB's Image Processing Toolbox with $\lambda = 100$.

As filters we consider

$$\begin{split} H_{1,2} &= \left(\begin{array}{cc} I \otimes H_1 \\ H_1 \otimes I \end{array}\right), \text{ where } H_1 = \left(\begin{array}{cc} 1 & -1 \\ \ddots & \ddots \\ & 1 & -1 \end{array}\right) \in \mathbb{R}^{(n-1) \times n}, \\ \\ H_{2,2} &= \left(\begin{array}{ccc} 4 & -1 & -1 \\ -1 & 4 & -1 & -1 \\ \ddots & \ddots & \ddots \\ -1 & -1 & 4 & -1 \\ & -1 & -1 & 4 \end{array}\right) \in \mathbb{R}^{N \times N}. \end{split}$$

Conclusions

• Both methods are stable w.r.t. the choice of the number of iterations, i.e. they do not require a reliable stopping rule

 $K_m(A, x_{\lambda}) = span\{x_{\lambda}, Ax_{\lambda}, ..., A^{m-1}x_{\lambda}\}$

For the construction of the subspaces $K_m(A, x_\lambda)$ the Arnoldi algorithm generates an orthonormal sequence $\{v_j\}_{j\geq 1}$, with $v_1 = x_\lambda / ||x_\lambda||$, such that $K_m(A, x_\lambda) = span \{v_1, v_2, ..., v_m\}$. For every m,

 $AV_m = V_m H_m + h_{m+1,m} v_{m+1} e_m^T.$

where $V_m = [v_1, v_2, ..., v_m]$, H_m is an upper Hessenberg matrix with entries $h_{i,j} = v_i^T A v_j$ and e_j is the *j*-th vector of the canonical basis of \mathbb{R}^m .

The *m*-th Arnoldi approximation to $x = f(A)x_{\lambda}$ is defined as

 $x_m = \|x_\lambda\| V_m f(H_m) e_1$

At each step of the Arnoldi algorithm we have to compute the vector $w_j = Av_j$. The method theoretically converges in a finite number of steps. For the computation of $f(H_m)$ we employ the Schur-Parlett algorithm (Golub and Van Loan 1983). $\|E_m\| \leq \|f(A)\overline{x}_{\lambda} - \|\overline{x}_{\lambda}\|V_m f(H_m)e_1\| + \|f(A)(x_{\lambda} - \overline{x}_{\lambda})\|$

- For small λ , $f(A) \approx I$ and $||E_m|| \approx ||x_{\lambda} \overline{x}_{\lambda}||$. The method is not able to improve the accuracy of the solution of the initial system.
- For large λ we have that $x_{\lambda} \approx \overline{x}_{\lambda}$, but even assuming that $\|f(A)(x_{\lambda} \overline{x}_{\lambda})\| \approx 0$, we have another lower bound due the ill-conditioning of f(A).

<u>ATP</u>. The error is given by

 $E_m := f(Q)x_\lambda - p_{m-1}(Q)\overline{x}_\lambda$

 $(p_{m-1} \text{ interpolates } f \text{ at the eigenvalues of } Q)$ where $(A^TA + \lambda H^TH)x_{\lambda} = A^Tb$ and $(A^TA + \lambda H^TH)\overline{x}_{\lambda} = A^T\widetilde{b}$. As before

 $||E_m|| \le ||f(Q)\overline{x}_{\lambda} - p_{m-1}(Q)\overline{x}_{\lambda}|| + ||f(Q)(x_{\lambda} - \overline{x}_{\lambda})||.$

Theoretically we may choose λ very large, thus oversmoothing, in order to reduce the effect of noise. Unfortunately, as before, f(Q) may be ill-conditioned for λ large. A compromise for the selection of λ is necessary but contrary to the ASP method, it is difficult to design a theoretical strategy.

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- They are as fast and accurate as the most effective iterative solvers
- W.r.t. classical preconditioned iterative solvers, only one linear system with the preconditioner has to be solved
- They do not require an accurate parameter-choice strategy for λ

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