University of Houston

COSC 3320: Algorithms and Data Structures Spring 2017

Homework 4

Due March 9, at the start of class

- 1. Given an array A[1, 2, ..., n] of n elements, a *majority* element of A is an element occurring at least $\lceil (n+1)/2 \rceil$ times. The elements cannot be ordered or sorted, but can be compared for equality. Design an efficient divide and conquer algorithm that returns a majority element of A (if any), and determine its complexity.
- 2. Consider the output A of the following algorithm.

```
w1 = 'blacksmith'
w2 = 'powdery'
A = new array of size 10
B = new array of size 7
for i = 0 to 9 do
        A[i] = (i+1)-th letter of w1
for i = 0 to 6 do
        B[i] = (i+1)-th letter of w2
        j = (i+1)-th digit of your 7-digit UH ID
        A[j] = B[i]
return A
```

You are told that $\sigma = A[0], A[1], \dots, A[9]$ is the sequence of nodes of a tree visited by some visit procedure.

- (a) Exhibit one tree T whose nodes are labeled with the letters in A and whose preorder visit would visit the nodes of T in the order given by σ .
- (b) Exhibit one tree T whose nodes are labeled with the letters in A and whose postorder visit would visit the nodes of T in the order given by σ .
- (c) Exhibit one binary tree T whose nodes are labeled with the letters in A and whose inorder visit would visit the nodes of T in the order given by σ .
- 3. Let T be a proper binary tree. Given a node $v \in T$, the *imbalance* of v, denoted imbalance(v), is defined as the difference, in absolute value, between the number of leaves of the left subtree of v and the number of leaves of the right subtree of v. (If v is a leaf, then imbalance(v) is defined to be 0.) Define imbalance(T) = $\max_{v \in T} \max_{v \in T} \max_{v \in T} (v)$.
 - (a) Provide an upper bound to the imbalance of a proper binary tree with n nodes, and exhibit a proper binary tree whose imbalance matches such an upper bound.
 - (b) Draw a proper binary tree T for which imbalance(T) = imbalance(v) and v is not the root of T.

- (c) Design an efficient algorithm to compute imbalance(T), and determine its complexity.
- 4. A proper d-ary tree, with $d \ge 1$, is a tree where every node has either 0 or d children.
 - (a) Draw two proper ternary (d = 3) trees T_1 and T_2 with 8 internal nodes each and such that T_1 has the maximum possible height and T_2 has the minimum possible height. How many leaves do T_1 and T_2 have?
 - (b) Let T be a proper d-ary tree with n nodes, m of which are leaves, and of height h. Which one of the following relations can be true for any T?
 - i. $m = (d-1)^h + 1$
 - ii. m = d(n m 1) 1
 - iii. m = (d 1)(n m) + 1
 - (c) Prove the relation guessed in (b).