

How are Cloud structured?

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- Clients talk to Clouds using web browsers or the web services standards
 - But this only gets us to the outer "skin" of the Cloud data center, not the interior
 - Consider Amazon: it can host entire company web sites (like Target.com or Netflix.com), data (AC3), servers (EC2) and even user-provided virtual machines!

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Many styles of system

- Near the edge of the Cloud (tier one), focus is on vast numbers of clients and rapid response
- Inside we find high volume services that operate in a <u>pipelined</u> manner, <u>asynchronously</u>
- Deep inside the cloud (tier three) we see a world of <u>virtual computer clusters</u> that are scheduled to share resources and on which applications like MapReduce (Hadoop) are very popular

Asynchronous pipeline model

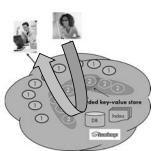
 The outside layer of the Cloud absorbs read accesses
 But writes (updates) are tricky

□ Popular model: guess at the

update outcome and respond

using that, but send the "real"

update asynchronously to a back-end server / database



□ Later, correct inconsistencies

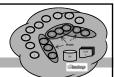
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In the outer tiers replication is key

- □ We need to replicate
 - Processing: each client has what seems to be a private, dedicated server (for a little while)
 - Data: as much as possible, that server has copies of the data it needs to respond to client requests without any delay at all
 - Control information: the entire structure is managed (orchestrated) in an agreed-upon way by a decentralized cloud management infrastructure

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Shared key-value store?



- The caching components running in tier two are central to the responsiveness of tier-one services
 - The idea is to always used cached data if possible, so that the inner services (here, a DB and a search index stored in a set of files) are shielded from "online" load
 - We need to replicate data within our cache to spread loads and provide fault tolerance
 - But not everything needs to be "fully" replicated
 Hence we often use "shards" with just a few replicas

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Sharding used in many ways

- - $\hfill\square$ Tier two could be any of a number of caching services:
 - Memcached: a sharable in-memory key-value store
 - Other kinds of Distributed Hash Tables that use key-value APIs
 - Dynamo: an Amazon service created as a scalable way to represent the shopping cart and similar data
 - BigTable: A very elaborate key-value store created by Google and used not just in tier two but throughout their "GooglePlex" for sharing information
 - □ The notion of sharding is cross-cutting
 - Most of these systems replicate data to some degree

Do we always need to shard data?

- Imagine a tier-one service running on 100k nodes
 Can it ever make sense to replicate data on the entire set?
- Yes, if some kinds of information might be so valuable that almost every external request touches it
 - Must think hard about patterns of data access and use
 - Some information needs to be heavily replicated to offer blindingly fast access on vast numbers of nodes
 - The principle is similar to the way Beehive operates
 - Even if we don't make a dynamic decision about the level of replication required, the principle is similar
 - We want the level of replication to match the level of load and the degree to which the data is needed on the critical path

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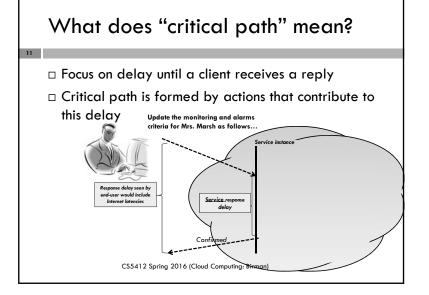
And it isn't just about updates

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- Should also be thinking about patterns that arise when doing reads ("queries")
 - Some can just be performed by a single representative of a service
 - But others might need the parallelism of having <u>several</u> (or even a huge number) of machines do parts of the work concurrently
- The term sharding is used for data, but here we might talk about "parallel computation on a shard"

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What if a request triggers updates?

If the updates are done "asynchronously" we might not experience much delay on the critical path

- Cloud systems often work this way
- Avoid waiting for slow services to process the updates but may force the tier-one service to "guess" the outcome
- For example, could optimistically apply update to value from a cache and just hope this was the right answer
- Many cloud systems use these sorts of "tricks" to speed up response time

First-tier parallelism

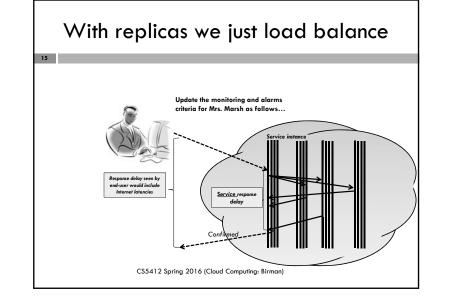
Parallelism is vital to speeding up tier-one services

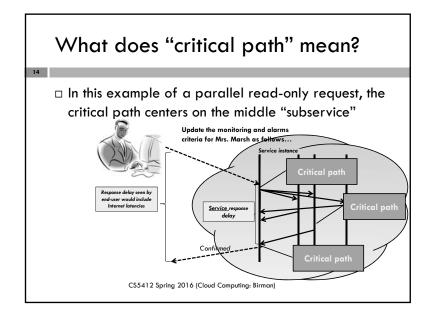
□ Key question:

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- Request has reached some service instance X
- Will it be faster...
- ... For X to just compute the response
- ... Or for X to subdivide the work by asking subservices to do parts of the job?
- □ Glimpse of an answer
 - Werner Vogels, CTO at Amazon, commented in one talk that many Amazon pages have content from 50 or more parallel subservices that run, in real-time, on your request!

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But when we add updates.... 16 Update the monitoring and alarms criteria for Mrs. Marsh as follows... Execution timeline for an individual first-tier replica С Soft-state first-tier service Send Response delay seen by end-user would also ... include Internet latencies 2-16 Now the delay associated with not measured in our wor waiting for the multicasts to finish could impact the critical path even in a single service 52 Confirmed CS5412 Spring 2016 (Cloud Computing:

What about updating w/o waiting?

Several issues now arise

- Are all the replicas applying updates in the same order?
- Might not matter unless the same data item is being changed
- But then clearly we do need some "agreement" on order
- What if the leader replies to the end user but then crashes and it turns out that the updates were lost in the network?
- Data center networks are surprisingly lossy at times
- Also, bursts of updates can queue up
- □ Such issues result in *inconsistency*

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Eric Brewer's CAP theorem

- □ In a famous 2000 keynote talk at ACM PODC, Eric Brewer proposed that "you can have just two from **Consistency, Availability and Partition Tolerance**"
 - He argues that data centers need very snappy response, hence availability is paramount
 - And they should be responsive even if a transient fault makes it hard to reach some service. So they should use cached data to respond faster even if the cached entry can't be validated and might be stale!
- □ Conclusion: weaken consistency for faster response

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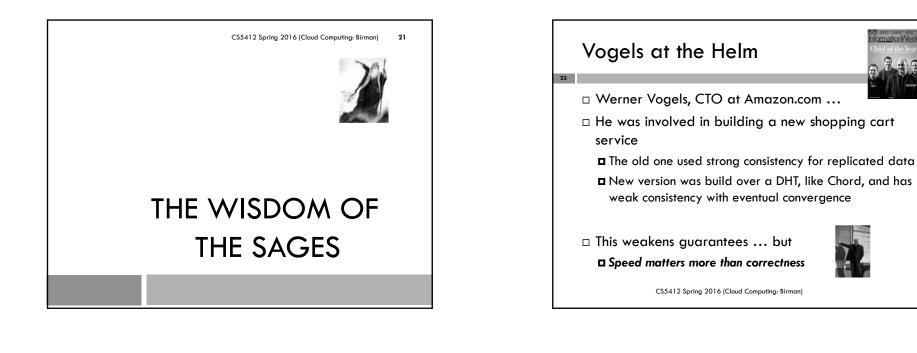
CAP theorem

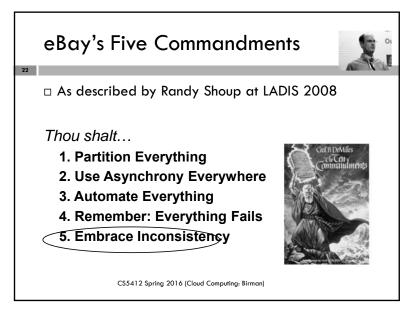
- □ A proof of CAP was later introduced by MIT's Seth Gilbert and Nancy Lynch
 - Suppose a data center service is active in two parts of the country with a wide-area Internet link between them
 - We temporarily cut the link ("partitioning" the network)
 - And present the service with conflicting requests
- □ The replicas can't talk to each other so can't sense the conflict
- □ If they respond at this point, inconsistency arises

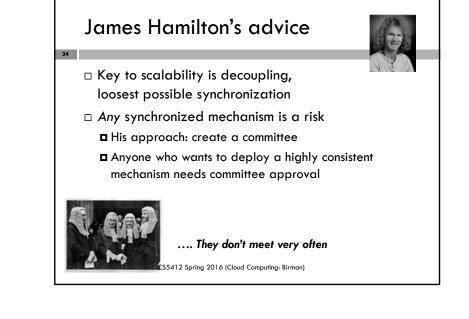
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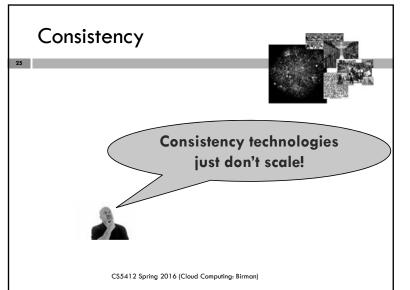
Is inconsistency a bad thing?

- □ How much consistency is really needed in tier one of the Cloud?
 - Think about YouTube videos: would consistency be an issue here?
 - What about the Amazon "number of units available" counters: will people notice if those are a bit off?
- □ Puzzle: can you come up with a general policy for knowing how much consistency a given thing needs?

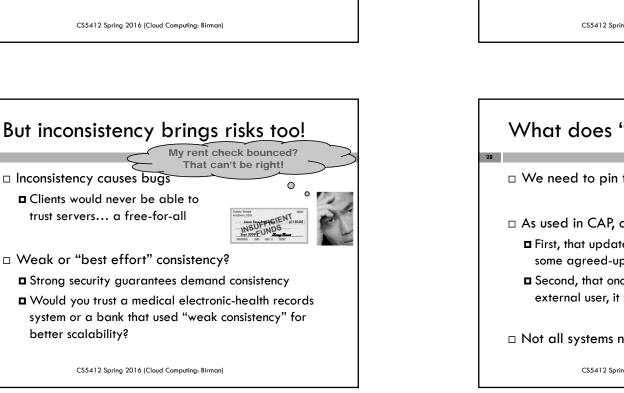


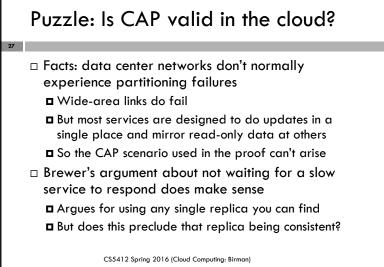






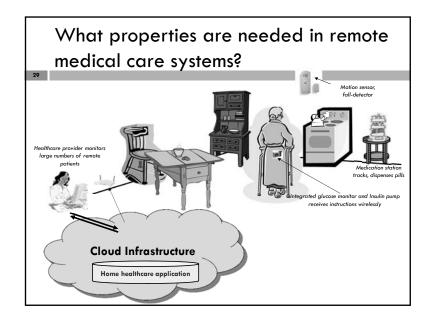
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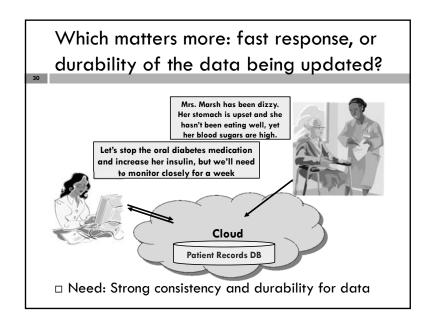


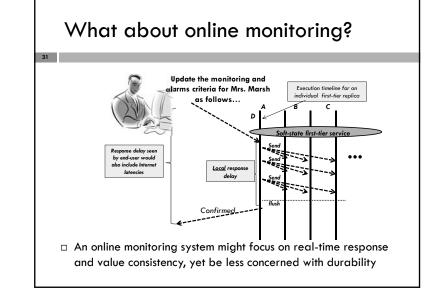


What does "consistency" mean?

- We need to pin this basic issue down!
 As used in CAP, consistency is about two things
 - First, that updates to the same data item are applied in some agreed-upon order
 - Second, that once an update is acknowledged to an external user, it won't be forgotten
- □ Not all systems need both properties







Why does monitoring have weaker needs?

When a monitoring system goes "offline" the device turns a red light or something on

- Later, on recovery, the monitoring policy may have changed and a node would need to reload it
- Moreover, with in-memory replication we may have a strong enough guarantee for most purposes
- Thus if durability costs enough to slow us down, we might opt for a weaker form of durability for better scalability and faster responses!

This illustrates a challenge!

- Cloud systems just can't be approached in a one-size fits all manner
- For performance-intensive scalability scenarios we need to look closely at tradeoffs
 Cost of stronger guarantee, versus
 Cost of being faster but offering weaker guarantee
- If systems builders blindly opt for strong properties when not needed, we just incur other costs!
 Amazon: Each 100ms delay reduces sales by 1%!

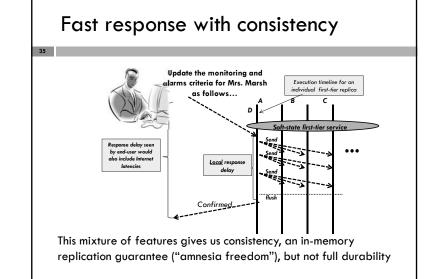
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Properties we might want

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- Consistency: Updates in an agreed order
- Durability: Once accepted, won't be forgotten
- □ Real-time <u>responsiveness</u>: Replies with bounded delay
- <u>Security</u>: Only permit authorized actions by authenticated parties
- □ <u>Privacy</u>: Won't disclose personal data
- <u>Resilience</u>: Failures can't prevent the system from providing desired services
- □ <u>Coordination</u>: Actions won't interfere with one another

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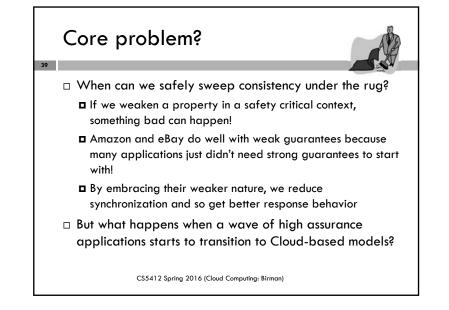
Does CAP apply deeper in the Cloud?

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- The principle of wanting speed and scalability certainly is universal
- But many Cloud services have strong consistency guarantees that we take for granted but depend on
- □ Marvin Theimer at Amazon explains:
 - Avoid costly guarantees that aren't even needed
 - But sometimes you just need to guarantee something
 - Then, be clever and engineer it to scale
 - And expect to revisit it each time you scale out 10x

Cloud services and their properties

Service	Properties it guarantees
Memcached	No special guarantees
Google's GFS	File is current if locking is used
BigTable	Shared key-value store with many consistency properties
Dynamo	Amazon's shopping cart: eventual consistency
Databases	Snapshot isolation with log-based mirroring (a fancy form of the ACID guarantees)
MapReduce	Uses a "functional" computing model within which offers very strong guarantees
Zookeeper	Yahoo! file system with sophisticated properties
PNUTS	Yahoo! database system, sharded data, spectrum of consistency options
Chubby	Locking service very strong guarantees
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Is there a conclusion to draw?

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- □ One thing to notice about those services...
 - Most of them cost 10's or 100's of millions to create!
 - Huge investment required to build strongly consistent and scalable and high performance solutions
 - Oracle's current parallel database: billions invested
- CAP isn't about telling Oracle how to build a database product...
 - CAP is a warning to <u>you</u> that strong properties can easily lead to slow services
 - But thinking in terms of weak properties is often a successful strategy that yields a good solution and requires less effort