On the multiple facets of synchronization

Runtimes for concurrency and distribution

Tullio Vardanega, tullio.vardanega@unipd.it

Academic year 2020/2021

Evaluating language features – 1

- The features of a programming language may be evaluated from two perspectives
- **Expressive power**
  - How far they help the user meet the application requirements
- **Usability**
  - How well those features interact (efficacy) versus how well they co-exist (coherence)

Evaluating language features – 2

- The synchronization constructs are an important ambit of evaluation
  - Toby Bloom, "Evaluating synchronisation mechanisms", Proc. 7th ACM Symposium on Operating System Principles (1979) https://doi.org/10.1145/800175.806566
- The cited work singles out 6 types of conditions on synchronization
  - Over and above exclusion synchronization

Conditions on synchronization – 1

1. Contingent on the synchronization state of the resource ☑
   - Number of current users or number of enqueued calls (`'Count`) in relation to the resource multiplicity
2. Contingent on the logical state of the resource ☑
   - No-write-on-full, no-read-from-empty
3. Contingent on the history of service ☑
   - For fairness, load balance, energy efficiency, …

Conditions on synchronization – 2

4. Contingent on the type of request ☑
    - Preferential treatment for some requests (e.g., writes over reads)
5. Contingent on the time of the request ☑
    - Reflected in queuing policies for calls and callers
6. **Contingent on the request parameters** ☒
    - Where serviceability depends on whether the resource can dispense as much as requested
        - E.g., as in paging or heap management

The resource allocation problem – 1

- Recurrent in any concurrent system
  - It involves all of the 6 dimensions
    - Our current model is unable to handle it properly
- **Example**
  a) A resource manager handles a statically fixed number N of resources $\{R_{j=1,\dots,N}\}$
  b) A number of concurrent clients $\{C_{i=1,\dots,M}\}$ may request any subset of such resources
  c) Accepted requests shall be satisfied fully
     i. Requests cannot return until satisfied
  d) Clients return resources after use

The resource allocation problem – 2

- Let us analyse the problem specification
  - The client interaction with the resource manager is synchronous
    - Wait-until-satisfied (cf. requirement c.i)
  - The volume of the request is specified as a parameter
    - This is the only plausible interface of the server
- What happens when the server finds itself unable to satisfy the request being examined
  - It cannot return to the caller prematurely
    - Hence it must keep that request on hold
  - How can it do that while continuing to serve others?
    - Serving others (cf. requirement d.) allows releases

The resource allocation problem – 3

- Do guards help?
  - They prevent synchronization when the required service cannot be delivered
  - But guards operate **before** synchronization takes effect, hence **before** the request itself is examined
  - Conclusions: guards as we know them, do not help
- Two alternatives are possible, both needing enhanced capabilities
  1. Allow guards expressions to access request parameters without this implying synchronization
  2. Transfer to another queue the request that presently cannot be satisfied
     - Beginning service but then holding it until further notice
     - Becoming able to serve other pending requests

Alternative 1

1 resource per request is the trivial case

```
protected Controller is
 entry Allocate (R : out Resource);
 procedure Release (R : Resource);
private
 Free : Natural := Full_Capacity;
 …
end Controller;
protected body Controller is
 entry Allocate (R : out Resource)
  when Free > 0 is
 begin
  Free := Free – 1;
  …
 end Allocate;
 procedure Release (R : Resource) is
 begin
  Free := Free + 1;
 end Release;
end Controller;
```

$1 \leq n_i \leq N$ resources per request is much harder

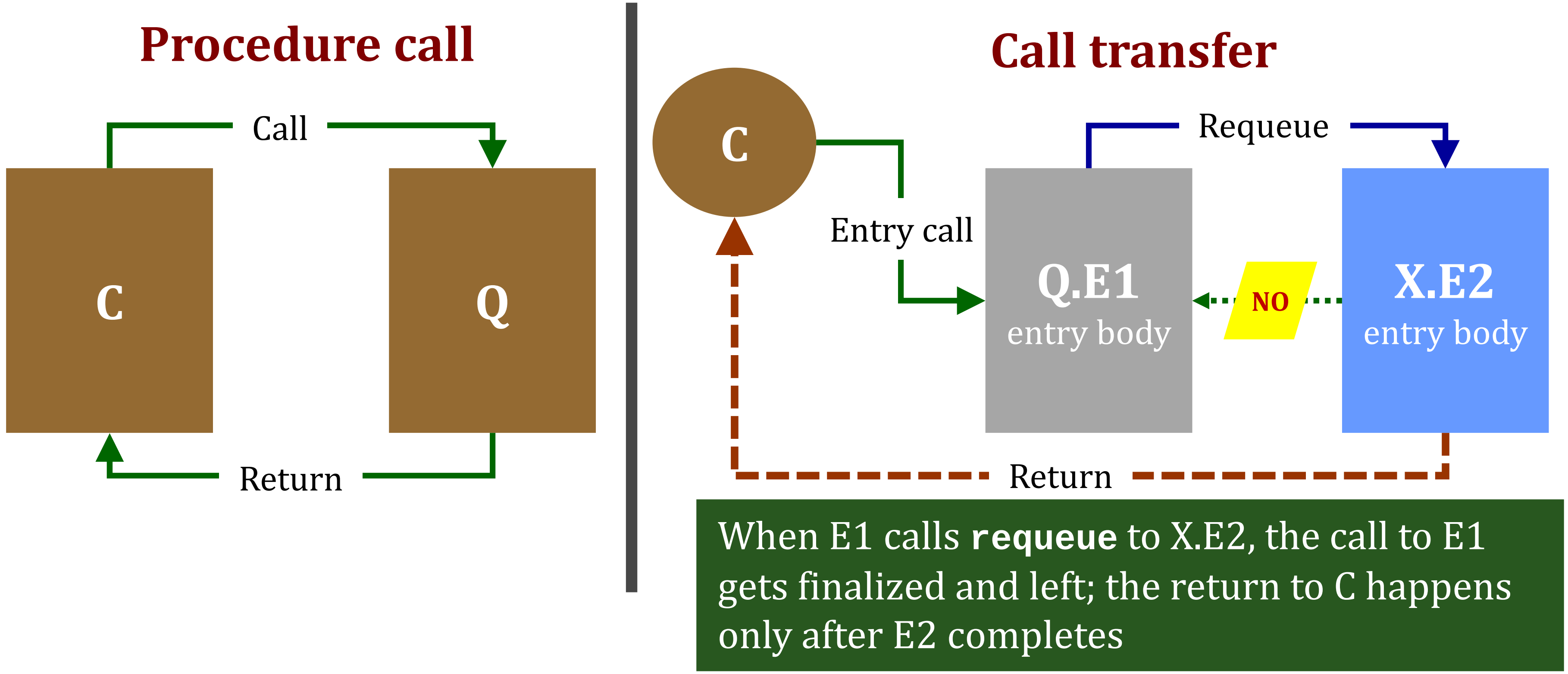
```
type Request is range 1..Max_Requests;
protected Controller is
 entry Allocate
  (R       :     out Resource;
   Amount : in      Request);
 procedure Release
  (R       : Resource;
   Amount : Request);
private
 Free : Request := Request'Last; …
end Controller;
protected body Controller is
 entry Allocate
  (R       :     out Resource;
   Amount : in      Request)
  when Amount <= Free is
 begin
  Free := Free – Amount;
 end Allocate;
 procedure Release (…) is …
end Controller;
```

Critique of alternative 1

- Requests that fail the guard are enqueued in the corresponding event queue (aka entry queue)
- Applying the eggshell model here would cause traversing the entire event queue every time a R/W access to the server state completes
  - Seeking any enqueued request that passes the guard
  - Untenably costly in the general case
- The problem is that each request has its own "state-change event"
  - The entry queue model that we know caters for a single-event queue and therefore does not really help

Alternative 2 – 1

- The call transfer (**`requeue`**) "reifies" the event queue: it is <u>not</u> a normal procedure call



Alternative 2 – 2

- A sophisticated feature, with challenging requirements on the runtime …
  - Transferring the call to another queue should **neither** suspend the server **nor** awake the client
  - Transfer should occur atomically, **without** evaluating the guard prefixed to the target queue
- This raises two "feature-interaction" questions
  1. Which entry queues can be allowable targets
  2. What happens to any time-out set on the call

Alternative 2 – 3

- **Question 1**
  - Any entry with a compatible interface, anywhere, even outside of the server, is an allowable target
    - The entry interface shall be either identical or with additional parameters all with default values, or with no parameters
- Implications
  - Transferring to a queue in the same server fits the eggshell model semantics nicely
    - In addition to yielding good functional cohesion
  - Transferring to an entry queue outside of the server requires releasing the R/W lock on it
    - Without the eggshell model semantics knowing what to do

Alternative 2 – 4

Question 2

- Two possible outcomes
  1) Call `B.E1` not accepted within `T1` is aborted
  2) Call `B.E1` accepted and then transferred to `B.E2`, is aborted if not accepted there within `T1`
- Outcome 2) incurs an ugly temporal distortion

What happens in this case?

```
-- client A
select
 B.E1;
or
 delay T1;
end select;
```

```
-- server B
select
 accept E1 do
  … -- T2 time units

  requeue E2 with abort;
 end E1;
or
 …
end select;
```

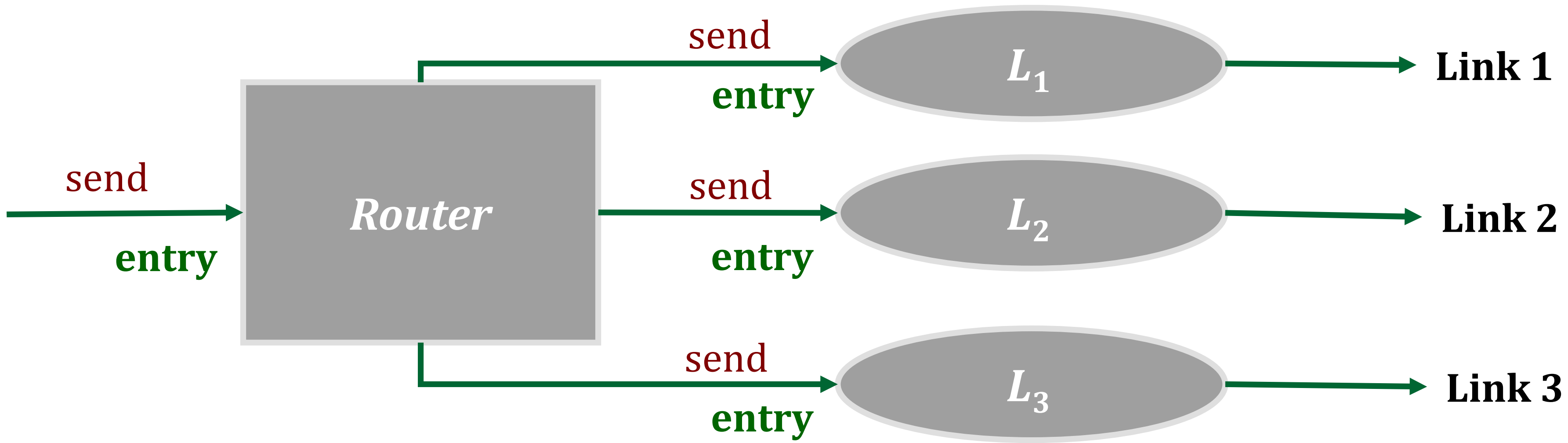
The **`with abort`** clause preserves the time-out effect upon call transfer

Use cases – 1

- Appropriate use of the `requeue` feature helps realize resource managers quite neatly
  - Check the implementation linked to today's lecture
- **Homework**
  - Try and improve the given solution in a manner that avoids useless call transfers
  - Try the same solution with a programming language of your choice

Use cases – 2

- A network router may forward inbound packets on to $N = 3$ outbound links
  - Link L_1 is the preferred choice, but the other links (first L_2 and then L_3) are used when L_1 risks overloading
- Likening packets to calls, and router and links to servers maps packet forwarding to a requeue



Flipped-class exercise

- Realize a circular-line metro service simulator
  - $M > 1$ train stations along a circular line
  - $N > M$ commuters who forever revolve around their duty cycle
    1. Go from home to the nearest train station
    2. Board the first possible train
    3. Get off the train and go to work (and work)
    4. Go from work to the nearest train station
    5. Board the first possible train
    6. Get off the train and go home (and rest)
  - 1 commuter train with capacity $C < N$ (no prebooking)